

# Node, Edge, Arc Routing and Turn Penalties: Multiple problems – One Neighborhood Extension

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This article explores a structural neighborhood decomposition for arc routing problems, in which the decisions about traversal orientations during services are made optimally as part of neighbor evaluation procedures. Using memory structures, bidirectional dynamic programming, and lower bounds, we show that a large neighborhood involving classical moves on the sequences of services along with optimal orientation decisions can be searched in amortized  $\mathcal{O}(1)$  time per move evaluation instead of  $\mathcal{O}(n)$  as in previous works. Because of its generality and now-reduced complexity, this approach can be efficiently applied to several variants of arc routing problems, extended into large neighborhoods such as ejection chains, and integrated into two classical metaheuristics. Our computational experiments lead to solutions of high quality on the main benchmark sets for the capacitated arc routing problem (CARP), the mixed capacitated general routing problem (MCGRP), the periodic CARP, the multi-depot CARP, and the min-max k-vehicles windy rural postman problem, with a total of 1528 instances. We also report sensitivity analyses for new MCGRP instances with turn penalties, which are uncommonly challenging yet critical for many practical applications.

*Key words:* Arc routing, General routing, Turn penalties, Service clusters, Heuristics, Structural problem decomposition, Local search, Large neighborhoods.

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## 1. Introduction

Arc routing problems arise in a wide variety of applications, including snow plowing, refuse collection, maintenance, and postal services. Practical issues introduce a variety of supplementary constraints, combined objectives, and decision sets, leading to a large family of problem variants (Corberán and Laporte 2014). Since exact solution methods are often impractical for industrial needs, many researchers have investigated heuristics. The literature on arc routing metaheuristics has thus seen a substantial increase in the last fifteen years, from a review of just four pages in Dror (2000) to more than one hundred articles now.

Arc routing is notably different from node routing, due to the necessity of considering precise network information and deciding on traversal orientations for the services on edges. In node routing algorithms, decisions related to the shortest paths between deliveries are usually concealed within the computation of the distance matrix. In arc routing applications, these path decisions are nontrivial, because they may be conditioned by the choices of service orientations, and they can also depend on other attributes such as delays at intersections and turn restrictions, which are essential features of urban applications. Recent capacitated arc routing (CARP) heuristics (e.g., Lacomme et al. 2004, Brandão and Eglese 2008, Usberti et al. 2013) thus deal with service orientations via additional neighborhoods. However, since the changes of orientation are considered only in specific move subsets, the resulting search spaces may be deceptive.

These distinct decision classes contributed to split the research effort on heuristics, between arc and node routing algorithms and concepts. This contrasts with the goal of several recent contributions on routing optimization, relating to unified metaheuristic strategies rather than problem-tailored developments (Irnich 2008a, Vidal et al. 2014a). Due to their purpose, these approaches could be extended to several variants of arc routing problems, but their success depends on our ability to deal efficiently and systematically with arc-routing-specific decision sets.

In this article, we consider a broad class of arc routing variants, with services on arcs, nodes, and edges as well as possible turn penalties, and formulate them as four main decision sets: ASSIGNMENT of services to routes, SEQUENCING, MODE CHOICES, and PATHS between services. The mode choices represent different ways of fulfilling a service, e.g., on either direction of a street or on a different lane. These choices arise naturally in the CARP, where two service orientations are possible for each edge, as well as in various advanced applications for snow plowing, street sweeping, salt spreading, meter reading, refuse collection, and courier delivery, among others (Corberán and Laporte 2014). Instead of designing separate neighborhoods for each decision class, we perform a heuristic search on a subset of the decisions while deducing exactly the others. A solution is formulated as sequences of services, without knowledge of the mode decisions. Alternative sequence and assignment decisions are heuristically enumerated by means of classical vehicle routing moves, while the optimal mode choices are derived via dynamic programming, using the approach of Beullens et al. (2003).

Such extended neighborhoods, known for a decade in the arc routing literature, have been scarcely used because of their higher computational complexity. In this article, we break this complexity barrier, demonstrating that each move can be evaluated in amortized  $\mathcal{O}(1)$  by means of efficient memory structures and bidirectional search. Moreover, by means of lower bounds on moves, we show that each evaluation requires roughly the same average number of elementary operations as those needed to evaluate distances and capacity constraints in a capacitated vehicle routing problem (CVRP). In other words, the mode decisions can be optimally addressed with little or no additional computational effort. The dynamic programming subproblem can also be used within polynomial ejection chains to search, in  $\mathcal{O}(n^2)$ , a exponential set of solutions obtained from chained service relocations and mode changes. Finally, we show that the approach is systematic: difficult decision sets for arc routing variants, such as turn penalties, can be fully concealed in the dynamic programming subproblems.

To investigate the improvement capabilities of these neighborhoods, we integrated them into two classical vehicle routing metaheuristics: the iterated local search (ILS) of Prins (2009) and the unified hybrid genetic search (UHGS) of Vidal et al. (2012, 2014a). We conducted extensive experiments on several arc routing variants: the CARP and the mixed capacitated general routing problem (MCGRP) – also called the node, edge, and arc routing problem (NEARP) – as well as problem extensions with turn penalties (CARP-TP and MCGRP-TP), multiple service periods (PCARP), multiple depots (MDCARP), and finally the min-max k-vehicles windy rural postman problem (MM-kWRPP). We considered a total of 1528 instances from 18 benchmark sets. With a single parameter setting, the proposed heuristics achieved results of very high quality for all the instance classes, matching or outperforming previous algorithms in the literature (some of which are fine-tuned for specific instances) while remaining conceptually simple and flexible. The key contributions of this article are thus the following:

- A full-fledged local search for the CARP, which evaluates moves on service sequences and jointly optimizes the service mode decisions in amortized  $\mathcal{O}(1)$  time per move evaluation;
- Lower bounds on move evaluations, which further reduce the computational effort;
- An extension of this methodology within ejection chains, exploring an exponential number of combinations of sequence and mode changes;
- An application to various arc routing problems via a broader interpretation of *service modes*;
- An integration of this neighborhood search into two state-of-the-art routing metaheuristics, leading to high-quality solutions for six important arc routing problem variants.

## 2. Problem statement

The CARP can be defined on a connected graph  $G = (V, E)$ , where  $V$  is the set of nodes and  $E$  the set of edges. A subset  $E_R \subseteq E$  of these edges must be serviced. Let  $n = |E_R|$  be the number of services. An edge  $(i, j) \in E$  can be traversed any number of times for a cost of  $c_{ij}$  each time, and a demand of  $q_{ij}$  is associated with any edge  $(i, j) \in E_R$ . The CARP aims to find a set of vehicle trips

with minimum cost, such that each trip starts and ends at a depot node  $v_0 \in V$ , each required edge is serviced by a single trip, and the total demand for any vehicle does not exceed a capacity  $Q$ .

The literature on CARP heuristics has grown rapidly in recent years. Most classical metaheuristic frameworks have been tested multiple times. Various tabu searches have been proposed (Hertz et al. 2000, Brandão and Eglese 2008), as well as a variable neighborhood descent (Hertz and Mittaz 2001), a guided local search (Beullens et al. 2003), and a GRASP (Usberti et al. 2013). Neighborhood-centered methods are often hybridized with population-based methods such as scatter search (Greistorfer 2003), path relinking (Usberti et al. 2013), genetic algorithms (Lacomme et al. 2001, 2004, Wang et al. 2015), ant colony optimization (Santos et al. 2010), and evolution strategies (Mei et al. 2013, 2014). Almost all these methods combine an aggressive optimization via local searches with diversification techniques such as crossovers, memories for guidance, or restarts from promising initial solutions.

Several variants of arc routing problems have been formulated to meet the needs of practical applications (see, e.g., Corberán and Laporte 2014, chapters 13–16). The MCGRP, in particular, can be defined on a connected multi-graph  $G = (V, E, A)$ , where  $V$  is the set of nodes,  $E$  the set of edges, and  $A$  the set of arcs. Services are required for a subset of nodes  $V_R \subseteq V$ , edges  $E_R \subseteq E$ , and arcs  $A_R \subseteq A$  such that  $n = |V_R| + |E_R| + |A_R|$ . Early research on this problem concerned constructive heuristics. Later, several metaheuristics have been introduced: a giant-tour-based hybrid GA (Prins and Bouchenoua 2005), a simulated annealing (Kokubugata et al. 2007), the *Spider* solver based on ILS and VNS (Hasle et al. 2012), a large neighborhood search and integer programming hybrid (Bosco et al. 2014), and an adaptive ILS (Dell’Amico et al. 2016). This last method generates high-quality results for a wide range of instances. All these metaheuristics use adaptations of classical neighborhoods for the CARP with eventual inversions of service orientations.

Another recurrent challenge of urban networks comes from possible turn restrictions and delays at intersections. These aspects are indispensable for commercial routing solvers, but more scarcely discussed in the academic literature. In arc routing problems with turn penalties, turn costs  $c_{ijk}$  are considered for each pair of connected edges and/or arcs  $(i, j)$  and  $(j, k)$ . Early research on this topic was primarily focused on ad-hoc constructive procedures and real applications. Later, heuristic and exact methods were proposed for single-vehicle settings (Benavent and Soler 1999, Corberán et al. 2002), based on problem reductions to generalized or asymmetric traveling salesman problems. Dussault (2012) and Dussault et al. (2014) consider the optimization of snow-plowing operations, leading to problem variants with multiple vehicles, a min-max objective, precedence constraints for deadheading, and turn penalties. Solution approaches based on integer programming, tour splitting, and cycle permutations were proposed. A few other articles discussed extensions of CARP with turn penalties (Lacomme et al. 2004, Bautista and Pereira 2004). The authors suggest specifying the service orientations in the solution representation, using a shortest path algorithm with turn penalties in the line graph. This approach works for extensions of the CARP but not in the presence of additional services on nodes as in the MCGRP-TP. For this last problem, we are aware of a single heuristic (Bräysy et al. 2011) based on a problem transformation into a generalized CVRP.

To extend the scope of the approach, we finally consider asymmetric travel costs and differentiate the travel (deadheading) cost  $c_{ij}$  and the service cost  $s_{ij}$  for any arc  $(i, j)$ . This allows us to solve *asymmetric* and *windy* arc routing problems (Benavent et al. 2009). Experimental analyses will also be reported for problems with multiple depots and visit periods. For the sake of brevity, these experiments are documented in the electronic companion of this paper.

### 3. A Question of Search Space

We consider a combinatorial optimization problem of the form  $\min_{x \in X} c(x)$ , where  $X$  is the solution space, and  $c$  is a cost function. A neighborhood is defined as a mapping  $\mathcal{N} : X \rightarrow 2^X$  that associates to each solution  $x$  a set of neighbors  $\mathcal{N}(x) \subset X$ . In the routing literature, neighborhoods are not directly specified as subsets  $\mathcal{N}(x)$ , but rather as the byproduct of a move definition. A move  $\phi$  is a local modification that can be applied to a given solution  $x$  to generate a neighbor  $\phi(x) \in \mathcal{N}(x)$ .

Examples of classical moves include the RELOCATE move, which changes the position of a service in the solution, and the SWAP move, which exchanges two services. A local search (LS) progresses iteratively from a solution to an improving neighbor until a local optimum is reached.

In the above definitions, the choice of solution space  $X$  appears to be transparent. This is, however, generally not the case, especially when a solution can be characterized by different decision sets. In production scheduling problems, for example, the solutions are often represented as sequences of activities for each machine, without precise information about their starting dates. The missing information is derived during solution evaluations, a task which is straightforward when the objective is *regular* (Giffler and Thompson 1960) but leads to a *timing* optimization problem in other cases (Vidal et al. 2015). For the CARP, a complete solution includes four decision subsets:

- ASSIGNMENT of services to routes,
- SEQUENCING of services for each route,
- MODE CHOICE for each service,
- PATHS between successive services,

and each of these four decision subsets leads to an exponential number of solutions. Now, when some of these decision sets are known (e.g., ASSIGNMENT and SEQUENCING), then the optimal choices for the other sets can be derived via dynamic programming. An incomplete solution representation can thus serve as a structural problem decomposition, dividing the search into the heuristic optimization of some decision sets and a simultaneous exact resolution of the others. By reducing the scope of the heuristic search, we aim to reduce its inherent error and to relegate a larger set of combinatorial decisions in the dynamic programming procedure. The main solution representations and neighborhoods used for the CARP are presented in Table 1.

R1	Explicit representation of ASSIGNMENT, SEQUENCING, MODE, and PATH decision subsets. In the figure, shortest paths between successive services (deadheading arcs) are represented with dashes.	
R2A	Each solution is represented as a set of routes. Each route is represented as a sequence of services with their modes (service orientations in the CARP). Thus, the ASSIGNMENT, SEQUENCING, and MODE decision subsets are explicitly represented.	
R2B	Each solution is represented as a giant tour, specifying a sequence of services with their modes. As such, only the SEQUENCING and MODE decision subsets are explicitly given.	In the representation R2A, each route is represented as illustrated above
R3A	Each solution is represented as a set of routes. Each route is represented as a sequence of services. As such, only the ASSIGNMENT and SEQUENCING decision subsets are explicit. The same representation is classically used for the CVRP.	
R3B	Each solution is represented as a giant tour containing a sequence of services. Hence, only the SEQUENCING decision subset is explicit. The same representation is classically used for the TSP.	In the representation R3A, each route is represented as illustrated above

**Table 1** Alternative solution representations for the CARP

First, in the complete solution representation (R1 in Table 1), all edges of the solution are specified, for both the services and the paths between services, thus giving explicitly the ASSIGNMENT,

SEQUENCING, MODE, and PATH decision subsets. Such a representation contains up to  $\mathcal{O}(n|V|)$  edges. This representation has been used in some seminal algorithms (e.g., Hertz et al. 2000), but it is now scarcely used in recent metaheuristics, since a preliminary computation of all-pairs distances allows to avoid the heuristic optimization of deadhauling paths. In mathematical programming methods, still, the sparsity of the underlying street network can be advantageously used (Letchford and Oukil 2009, Bode and Irnich 2012).

Representation R2A specifies the sequence of services to edges with their modes. Shortest-paths between edge endpoints have been processed prior to routing resolution. It requires  $\mathcal{O}(n)$  space per solution and  $\mathcal{O}(|V|^2)$  space for the shortest-path information. This representation, which allows move evaluations in  $\mathcal{O}(1)$  time, is used in all recent CARP metaheuristics during the local searches. However, as the SEQUENCING and MODE decisions are tightly correlated, some moves with joint modifications of these two decision sets are needed to progress toward high-quality solutions.

Representation R2B, without trip delimiters, has been frequently used during crossover operations in population metaheuristics (Lacomme et al. 2001, 2004). The  $\mathcal{O}(n^2)$  computational effort of the Split algorithm, used to decode the solutions, is affordable when used once before the LS, but not when systematically used during all move evaluations. For this reason, the above metaheuristics revert to R2A during the LS phases.

Solution representations without explicit MODE CHOICES (R3) have been rarely used until now. Beullens et al. (2003) perform mode optimization as a stand-alone procedure, and Irnich (2008b) combines mode optimization with the large neighborhood of Balas and Simonetti (2001), leading to promising results for mail delivery applications. Muyldermans et al. (2005) evaluate 2-OPT moves on R3A in  $\mathcal{O}(n)$  time per move. SWAP moves have been used on R3B by Wøhlk (2003, 2004), in a simulated annealing algorithm. However, the dynamic programming decoder, referred to as *Split with flips* in Prins et al. (2009), requires  $\mathcal{O}(n^2)$  operations per move evaluation. Ramdane-Cherif (2002) also suggest using R3B during crossover operations and reverting to R2A during LS. Overall, R3A and R3B are desirable since they conceal the MODE CHOICES in the dynamic programming subproblems, and thus guarantee the optimality of these decisions. However, the price to pay was, until now, a higher computational complexity for the move evaluations.

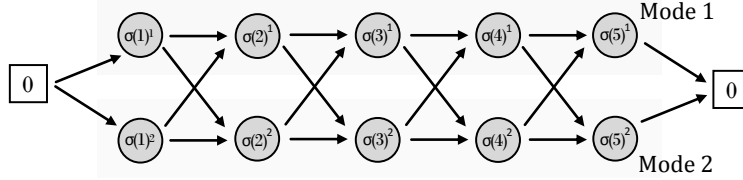
In the next section, we explain how to efficiently evaluate the solutions R3A in  $\mathcal{O}(1)$  instead of  $\mathcal{O}(n)$  in the context of a local search. As such, the heuristic search can be fully focused on the ASSIGNMENT and SEQUENCING decision subsets, as in CVRP metaheuristics, while the MODE CHOICE and PATH decisions are implicitly and optimally determined with little or no additional computational effort.

## 4. Neighborhood Search with Optimal Mode Decisions

In the following, a solution  $x \in X$  will be represented as a set of routes  $\mathcal{R}(x)$ . Each route  $\sigma \in \mathcal{R}(x)$  corresponds to a sequence of services  $\sigma = (\sigma(1), \dots, \sigma(|\sigma|))$ , starting and ending at the depot (counted as a dummy service, such that  $\sigma(1) = 0$  and  $\sigma(|\sigma|) = 0$ ). Any service  $i$  can be performed using one of a subset of modes  $M_i$ , with a service cost  $s_i^k$  for mode  $k \in M_i$ . We assume that the number of possible modes per service is bounded by a constant. Two modes are used for the CARP, one for each admissible service orientation. By convention,  $|M_0| = 1$  and  $s_0^1 = 0$  for the depot. The cost of traveling between services  $i$  and  $j$  in modes  $k \in M_i$  and  $l \in M_j$  is given by  $c_{ij}^{kl}$ . For the CARP, this corresponds to the distance of the shortest path between the exit node of edge  $i$  in mode  $k$ , and the entry node of edge  $j$  in mode  $l$ .

### 4.1. Optimization of mode choices

For any given sequence of services  $\sigma$ , an optimal selection of modes can be efficiently determined by dynamic programming as in Beullens et al. (2003). This procedure corresponds to the search for a shortest path in an acyclic auxiliary graph  $H_\sigma = (V_\sigma, A_\sigma)$  containing  $\sum_{i=1}^{|\sigma|} |M_{\sigma(i)}|$  nodes, one for each mode of each service. For each consecutive pair of services  $\sigma(i)$  and  $\sigma(i+1)$ , between each node pair associated with modes  $k \in M_{\sigma(i)}$  and  $l \in M_{\sigma(i+1)}$ , an arc is added in  $A_\sigma$  with cost  $c_{\sigma(i)\sigma(i+1)}^{kl} + s_{\sigma(i)}^k$ .



**Figure 1** Shortest path subproblem for the evaluation of a route  $\sigma$  with optimal mode choices.

Figure 1 illustrates the graph  $H_\sigma$  for the CARP. The shortest path problem can be solved in  $\mathcal{O}(n)$ , using the Bellman algorithm in the presence of a topological order. However, even linear-time move evaluations would be a significant computational bottleneck, such that further improvements are needed to efficiently employ this structure.

#### 4.2. Move evaluations by concatenation

To reduce this computational complexity, we propose incremental move evaluations in which the shortest-path subproblems are not solved independently. Instead, all-pairs shortest paths are preprocessed in the incumbent solution and then used for move evaluations. Thus, for each route  $\sigma$  in the incumbent solution, and each subsequence of consecutive services  $\bar{\sigma} \subset \sigma$ , the method keeps track of the cost of the partial shortest path  $C(\bar{\sigma})[k, l]$  in  $H_{\bar{\sigma}}$  between the first and last service in the sequence, for any combination of modes  $k$  and  $l$ , as well as the sum of the service demands  $Q(\bar{\sigma})$ .

The values  $C(\bar{\sigma})[k, l]$  are computed by induction on the operation of concatenation ( $\oplus$ ) of two sequences. For a sequence  $\bar{\sigma} = (v_i)$  containing a single service,  $C(\bar{\sigma})[k, l] = s_i^k$  for  $k = l$ , else  $C(\bar{\sigma})[k, l] = +\infty$ . Moreover, the following equations allow us to derive  $C$  and  $Q$  for any sequence  $\sigma_1 \oplus \sigma_2$  resulting from a concatenation of two sequences  $\sigma_1$  and  $\sigma_2$ , the first service of  $\sigma_2$  being done immediately after the last service of  $\sigma_1$ :

$$C(\sigma_1 \oplus \sigma_2)[k, l] = \min_{x \in M_{\sigma_1}(\sigma_1)} \left\{ \min_{y \in M_{\sigma_2}(1)} \left\{ C(\sigma_1)[k, x] + c_{\sigma_1(\sigma_1)\sigma_2(1)}^{xy} + C(\sigma_2)[y, l] \right\} \right\} \quad (1)$$

$$Q(\sigma_1 \oplus \sigma_2) = Q(\sigma_1) + Q(\sigma_2) \quad (2)$$

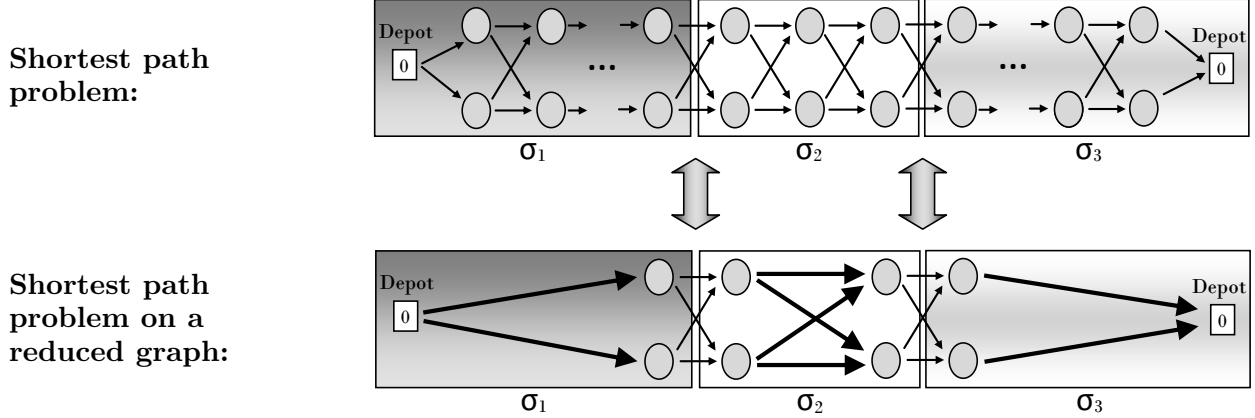
First, Equations (1–2) are used to preprocess the shortest paths in the incumbent solution in  $\mathcal{O}(n^2)$ : the services are enumerated according to their visit order, and the subsequences are evaluated. Then, the same equations are used for move evaluations. Any route  $\sigma$  obtained from a classical LS move corresponds to a concatenation of  $K$  service sequences from the incumbent solution:  $\sigma = \sigma_1 \oplus \dots \oplus \sigma_K$ . Evaluating the cost of  $\sigma$  involves evaluating the cost of a shortest path in  $H_\sigma$ , which is obtained via  $K - 1$  applications of Equation (1). As illustrated in Figure 2, this is equivalent to contracting known shortest paths in each sequence  $\sigma_i$  of  $H_\sigma$ , and then evaluating the shortest path on a reduced graph using the Floyd–Warshall algorithm.

**PROPOSITION 1.** *Using the proposed approach, the complexity of move evaluations for classical neighborhoods such as RELOCATE, SWAP, 2-OPT, and 2-OPT\* is in  $\mathcal{O}(1)$ .*

Indeed, any classical move on sequences can be assimilated to a recombination of a bounded number of sequences of services (Vidal et al. 2014a). Thus, each move evaluation is performed with  $\mathcal{O}(1)$  calls to Equations (1–2), requiring  $\mathcal{O}(1)$  elementary arithmetic operations since the number of alternative modes is bounded.

#### 4.3. Lower bounds on move evaluations

Even with an  $\mathcal{O}(1)$  complexity, the use of the *min* operator and the enumeration of modes leads to a constant but larger computational effort for move evaluations ( $\approx 4\times$  for the CARP when compared to the CVRP). To further reduce the CPU time, we propose a two-step evaluation procedure using



**Figure 2** Using preprocessed information on sequences of services to reduce the auxiliary graph  $H_\sigma$

lower bounds on moves. Consider two routes  $\sigma_1$  and  $\sigma_2$ , subject to a move  $\Pi$  that produces two new routes  $\sigma'_1$  and  $\sigma'_2$ . The move  $\Pi$  is improving if and only if  $\Delta_\Pi = C(\sigma'_1) + C(\sigma'_2) - C(\sigma_1) - C(\sigma_2) < 0$ . Let  $C_{\text{LB}}(\sigma')$  be a lower bound on the cost of a route  $\sigma'$ ; then any improving move fulfills the following condition, which can be used as a filter during a preliminary evaluation:

$$\Delta_\Pi^{\text{LB}} = C_{\text{LB}}(\sigma'_1) + C_{\text{LB}}(\sigma'_2) - C(\sigma_1) - C(\sigma_2) < 0. \quad (3)$$

Let  $C_{\text{MIN}}(\sigma) = \min_{k \in M_{\sigma(1)}} \{ \min_{l \in M_{\sigma(|\sigma|)}} \{ C(\sigma)[k, l] \} \}$  be the minimum distance of a shortest path in  $H_\sigma$  between any pair of modes for the first and last service of  $\sigma$ . Let  $c_{ij}^{\text{MIN}} = \min_{k \in M_i} \{ \min_{l \in M_j} \{ c_{ij}^{kl} \} \}$  be the minimum cost of a shortest path in  $G$  between services  $i$  and  $j$ , for any mode pair  $k \in M_i$  and  $l \in M_j$ . The following equation provides a lower bound on the cost of a route  $\sigma = \sigma_1 \oplus \dots \oplus \sigma_K$  composed of a concatenation of  $K$  sequences:

$$C_{\text{LB}}(\sigma_1 \oplus \dots \oplus \sigma_K) = \sum_{j=1}^K C_{\text{MIN}}(\sigma_j) + \sum_{j=1}^{K-1} c_{\sigma_j(|\sigma_j|)\sigma_{j+1}(1)}^{\text{MIN}}. \quad (4)$$

The values  $C_{\text{MIN}}(\sigma)$  can be preprocessed simultaneously with  $C(\sigma)[k, l]$ , and the shortest paths  $c_{ij}^{\text{MIN}}$  are directly evaluated from the distances  $c_{ij}^{kl}$ , prior to the routing optimization. Then, the evaluation of this lower bound takes exactly the same number of elementary operations as a classical move evaluation for the CVRP. Furthermore, most moves for routing problems are far from profitable, even when restricting moves to close services as in Toth and Vigo (2003). In our experiments (Section 5.2), we observed that this simple lower bound helped to discard 90% of the moves on average, and the exact evaluation of the remaining moves is no longer a computational bottleneck.

#### 4.4. Local search

The proposed techniques can serve as a building block for different local search algorithms, depending on the nature of the moves, the move acceptance policy, possible neighborhood restrictions, and the use of infeasible solutions, among other factors. The local search used in our computational experiments is summarized in Algorithm 1 and discussed in details in the rest of this section.

We rely on 2-OPT, 2-OPT\*, as well as RELOCATE and SWAP moves of up to  $k = 2$  consecutive services with possible reversals (Laporte et al. 2014). All together, these moves define the neighborhood  $\mathcal{N}(x^t)$  of an incumbent solution  $x^t$ . These moves consider service exchanges and relocations within the same route or between different routes, thus allowing an optimization of both SEQUENCING and ASSIGNMENT decision subsets. To speed up the search, the moves are only attempted between service pairs  $(i, j)$  where  $j$  belongs to a set  $\Gamma(i)$  of  $|\Gamma|$  closest services (Toth and Vigo 2003, Vidal et al. 2014a), the distance between services being defined as the shortest distance

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Input: An initial solution  $x^0$ 
Set  $t = 0$ 
repeat
  Update the auxiliary data structures for the routes of  $x^t$ . //  $\mathcal{O}(n^2)$  operations
  for each move  $\phi$  in random order, such that  $\phi(x) \in \mathcal{N}(x^t)$  //  $\mathcal{O}(n^2)$  possible moves
  do
    The move  $\phi$  can modify two routes of  $x^t$ , at most. Let  $z_{\text{BEFORE}}$  be the sum of the costs of these two
    routes in  $x^t$ , and let  $(\sigma_1, \dots, \sigma_K)$  and  $(\sigma'_1, \dots, \sigma'_L)$  be the sequences of services which form the
    two new routes in  $\phi(x)$ .
    Evaluate a lower bound on the cost of the new routes: //  $\mathcal{O}(1)$  operations
     $z_{\text{LB}} = C_{\text{LB}}(\sigma_1 \oplus \dots \oplus \sigma_K) + C_{\text{LB}}(\sigma'_1 \oplus \dots \oplus \sigma'_L)$ .
    if  $z_{\text{LB}} \geq z_{\text{BEFORE}}$  then continue (to the next move evaluation).
    Evaluate the cost of the new routes with optimal mode choices, //  $\mathcal{O}(1)$  operations
    using dynamic programming and the known auxiliary data structures:
     $z_{\text{AFTER}} = C(\sigma_1 \oplus \dots \oplus \sigma_K) + C(\sigma'_1 \oplus \dots \oplus \sigma'_L)$ .
    if  $z_{\text{AFTER}} \geq z_{\text{BEFORE}}$  then continue (to the next move evaluation).

    At this stage,  $\phi$  is known to be an improving move:
    Set  $x^{t+1} = \phi(x)$ ,  $t = t + 1$ , and break
  end
until a local minimum is attained.
Output: A local minimum  $x^t$ 

```

**Algorithm 1:** Local search using preprocessing, concatenations and lower bounds

between any two edge extremities. Furthermore, a “first improvement” policy is used: the moves are enumerated in random order and any improvement is directly applied.

**Capacity constraints.** As in previous works, we allow a controlled exploration of penalized infeasible solutions that violate the capacity constraints. Hence, the penalized cost of a route is evaluated as  $C_p(\sigma) = C(\sigma) + \omega \max\{0, Q(\sigma) - Q\}$ , where  $\omega$  is a penalty factor. This factor remains constant during one LS, and is self-adapted in the metaheuristic to attain a desired ratio of feasible solutions (Vidal et al. 2014a). The same penalty term is used in the evaluations of the lower bounds. Note that, in a LS variant that does not accept infeasible solutions, a load feasibility check should be included before the lower-bound evaluations.

**Preprocessing limitations.** Finally, the preprocessing phase represents only a small part of the overall computational effort for most benchmark instances. Still, additional limitation strategies can be used to prevent a larger time consumption in cases with few vehicles and long routes. In particular, Irnich (2008a) proposes limiting the preprocessing to a hierarchy of  $\mathcal{O}(n^{4/3})$  or  $\mathcal{O}(n^{8/7})$  sequences, at the cost of slightly more expensive but still constant-time move evaluations. To make this even simpler, we limited the preprocessing to sequences that start or end at the depot, or of a size smaller than ten. This still allows to evaluate inter-route moves in constant time (since these moves lead to routes of the form  $\sigma = \sigma_1 \oplus \bar{\sigma} \oplus \sigma_2$ , where  $\sigma_1(1) = 0$ ,  $\sigma_2(|\sigma_2|) = 0$ , and  $|\bar{\sigma}| \leq 2$ ), and the remaining intra-route moves, far less numerous, can be evaluated via a concatenation of a small (but possibly linear) number of sequences  $\bar{\sigma}$  such that  $|\bar{\sigma}| \leq 10$ .

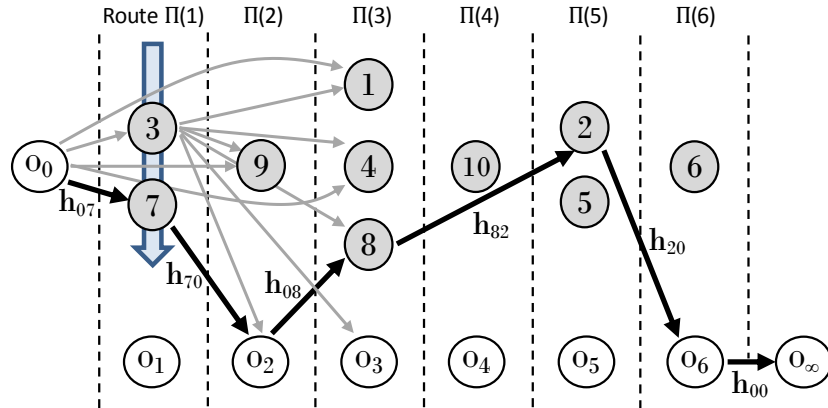
#### 4.5. Generalizations – Polynomial ejection chains

The previous sections have introduced a local search which evaluates an exponential number of solutions obtained from combined mode choices and simple service relocations and exchanges. We now generalize this technique into a variant of ejection chains (Glover and Rego 2006) that identifies, in  $\mathcal{O}(n^2)$  elementary operations, an improved solution obtained from multiple chained service



relocations among routes as well as combined optimal choices of service modes. The procedure works as follows:

- 1) Pick a random permutation  $\pi$  of the routes in the current solution.
- 2) Define the ejection graph  $H_\pi = (V_\pi, A_\pi)$  where
  - $V_\pi$  includes one node per service request as well as one *null* node per route, the purpose of which will be described below. Let  $R(i)$  be the corresponding route of a node  $i$ .
  - For any pair of nodes  $i \in V_\pi$  and  $j \in V_\pi$ ,  $A_\pi$  contains an arc  $(i, j)$  if and only if the route of  $i$  precedes the route of  $j$  in the permutation  $\pi$ . Note that  $i, j$ , or both can be *null* nodes.
  - The cost  $h_{ij}$  of an arc  $(i, j)$  corresponds to the difference in the cost of  $R(j)$  when *removing service  $j$  and inserting service  $i$  in its place*. A null node  $i$  stands for *inserting nothing*, while a null node  $j$  stands for *removing nothing*. In the case where a node is inserted and *nothing* is removed, the best insertion position is used to define the cost. All route costs are evaluated for sequences of services with an optimal choice of service modes. Note that the costs  $h_{ij}$  can be negative.
  - Complete the graph with a null node standing for a source, connected to all the service nodes, and a null node standing for a sink to which all the null nodes are connected.
- 3) Find a shortest path in this directed acyclic graph between the source and the sink. This can be done in  $\mathcal{O}(|V_\pi| + |A_\pi|) = \mathcal{O}(n^2)$  elementary operations, using Bellman's algorithm in topological order. If the path has a negative cost, apply the corresponding sequence of service removals and insertions to the current solution.



**Figure 3** Ejection graph  $H_\pi$  and a possible ejection chain in boldface. Grey and white colors are used for services and null nodes, respectively.

The ejection graph  $H_\pi$  is illustrated in Figure 3 for a problem with ten services in six different routes. The resulting shortest path relocates service 7 from route R1 to R2, relocates service 8 from R3 to R5, and relocates service 2 from R5 to R6.

#### 4.6. Generalizations – Other mode decisions

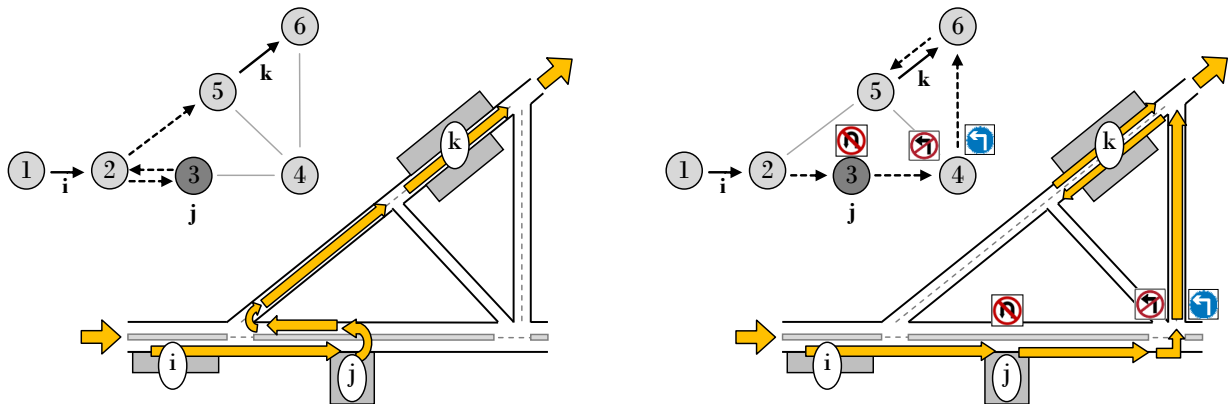
This section discusses extensions of the proposed methodology to several problem variants, including node, edge, and arc routing, turn penalties, and service clusters.

**4.6.1. Services on nodes, edges, and arcs.** Most existing metaheuristics for the MCGRP (Prins and Bouchenoua 2005, Kokubugata et al. 2007, Hasle et al. 2012) translate the variety of service types into a larger number of move classes, dealing with multiple cases of service location and orientation changes. Alternatively, Irnich et al. (2015) suggest to transform each service on a node into a service on a loop. In the proposed methodology, we can associate one or two modes for each service depending on its nature:

NODE	$ M_i  = 1$	One mode for service;
ARC	$ M_i  = 1$	One mode for the only feasible service orientation;
EDGE	$ M_i  = 2$	Two modes, one for each service orientation.

Move evaluations are then done exactly as before, the subproblem being even simpler since many services possess a single mode.

**4.6.2. Turn restrictions and delays at complex intersections.** Turns and delays at intersections can account for 30% of the total transit time in urban networks (Nielsen et al. 1998). As such, considering these features during optimization is a necessity for multiple applications. Shortest path problems with turn restrictions and penalties can be solved efficiently via tailored labeling approaches (Gutiérrez and Medaglia 2008) or graph transformations (line graph or node splitting – Vanhove and Fack 2012). However, representing an MCGRP solution as a permutation of services and using the shortest path with turn restrictions between each pair of services does not necessarily lead to a feasible overall solution. For example, the left part of Figure 4 shows a route with three consecutive services: to an edge  $i$ , a node  $j$ , and an edge  $k$ . The shortest paths between services  $i$  and  $j$  and between  $j$  and  $k$  do not contain any forbidden turns. However, the arrival direction at  $j$  is the opposite of the starting direction of the next path, leading to an infeasible U-turn.



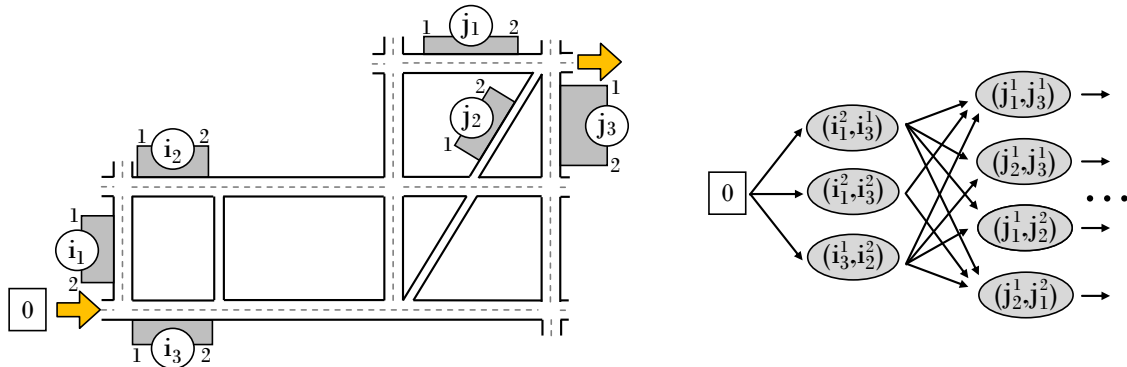
**Figure 4** Shortest path for a given sequence of three services. On the left, without turn penalties or constraints; and on the right, in the presence of one forbidden U-turn and other turn limitations.

This inconsistency is due to a lack of characterization of the arrival edge when servicing the node  $j$ . The proposed methodology provides a simple solution, by including the arrival edge and driving direction as part of the definition of the mode for a node:

NODE	$ M_i  = p_i$	$p_i$ modes to specify the arrival direction, where $p_i$ is the in-degree of $v_i$ ;
ARC	$ M_i  = 1$	One mode for the only feasible service orientation;
EDGE	$ M_i  = 2$	Two modes, one for each service orientation.

With this definition, the orientation of a service on an edge and the direction of arrival when servicing a node are implied by the mode. The travel times are computed, prior to the optimization, between any pair of (service, mode) couples, i.e., any pairs of arcs in the graph. This is done via an all-pairs shortest path algorithm on the line graph. The other steps are unchanged. The resulting algorithm produces the alternative turn-feasible solution illustrated on the right of Figure 4.

**4.6.3. Service clusters.** Problem aggregation is a natural way to deal with large-scale applications that contain multiple drop points. A natural aggregation occurs, in arc routing problems, when assimilating drops on the same street to a single edge service. This aggregation implies that these drops are done consecutively, hence forbidding split deliveries. Further aggregations can also be relevant, e.g., considering a geographically delimited group of visits as a service cluster, leading to a *clustered* routing problem, which aims to find routes in which all the services of the same cluster are done consecutively (Battarra et al. 2014).



**Figure 5** On the left, an example of a street network with two clusters of deliveries:  $\{i_1, i_2, i_3\}$  and  $\{j_1, j_2, j_3\}$ . Each delivery on an edge is represented by a gray rectangle, with two possible orientations (1 and 2). On the right, the associated shortest path problem for mode optimization.

This generalization can be handled via a broader definition of the service mode. A service now represents a cluster of visits, and the service mode of a cluster corresponds to one (entry, exit) direction pair for the cluster, as illustrated in Figure 5. The set of possible (entry, exit) pairs may be exhaustive, limited by practical constraints, or restricted to nondominated choices. The associated service costs are obtained by solving a variant of rural postman problem in the cluster where we impose to start and end at prescribed nodes. This can be done before starting the routing optimization, either exactly in the presence of small clusters, or via efficient heuristics. The other components of the approach are unchanged.

Finally, note that any CARP instance can be reduced to a clustered vehicle routing instance with two nodes in each cluster, one for each end of the associated edge. The visit-order choices in the cluster, i.e., the modes, are then equivalent to the service orientation in the original problem. This immediate reduction highlights the close connections between these problem classes.

**4.6.4. Other characteristics.** Service-mode choices arise in many other problem applications. The generalized VRP, for example, considers for any service  $i$  a set of  $g_i$  delivery locations, and one location per group must be visited. The proposed neighborhoods can be applied to this setting by assimilating a service to a group of locations and considering  $g_i$  modes per service. The generalized VRP may be viewed as the archetype of a routing problem with mode choices. Each delivery location is an alternative mode with a distinct cost. In this context, problem reductions from arc routing to generalized VRP arise naturally (Baldacci et al. 2009, Micó and Soler 2011).

Some other problems involve generalized concepts of the service mode. These include decisions about visiting an intermediate facility (Polacek et al. 2008), returning to a depot (Vidal et al. 2014b), selecting customers (Vidal et al. 2016), moderating speed (Kramer et al. 2015), or doing a preliminary stop at a charging station (Schneider et al. 2014). In some cases, the mode choices may involve other resources such as time, capacity, and service levels, as objectives or constraints. The methods of this paper can be further generalized to these applications, but leading to resource constrained shortest path subproblems (RCSP) in cases where several resources are involved. The RCSP is NP-hard in general, but many specific variants can be efficiently addressed (Irnich and Desaulniers 2005, Lozano and Medaglia 2013).

## 5. Computational Experiments

We conducted extensive experimental analyses to investigate the effectiveness of the proposed neighborhoods. This section describes the metaheuristic frameworks used for these tests, the benchmark instances, a comparison of the resulting solutions with previous state-of-the-art algorithms, an analysis of method scalability, and finally new benchmark instances and analyses for CARP and MCGRP with turn penalties.

### 5.1. An extension of two classical metaheuristics

The extended neighborhoods were included in two classical metaheuristic frameworks for vehicle routing: the multi-start iterated local search (ILS) of Prins (2009), and the unified hybrid genetic search (UHGS) of Vidal et al. (2012, 2014a). These metaheuristics were then applied to the CARP without any other problem-specific adaptation or tuning.

ILS starts from a random initial solution. Subsequently,  $n_C$  solutions are iteratively generated by applying a shaking operator and a local search, the best solution being kept as the new incumbent solution. The shaking operator applies  $k = 2 + \lfloor n/200 \rfloor$  random exchanges of services on a representation of the solution without trip delimiters, and a Split algorithm is used to reinsert the delimiters. The search is restarted  $n_P$  times, each run ending when  $n_I$  consecutive iterations have been performed without improvement of the best solution or when a time  $T_{\text{MAX}}$  is reached. The best overall solution is returned.

UHGS follows the methodology of Vidal et al. (2012) with the original binary tournament selection, crossover, and advanced population-diversity management operators. The population is managed to contain between  $\mu^{\text{MIN}}$  and  $\mu^{\text{MIN}} + \mu^{\text{GEN}}$  solutions, and each new individual is generated by an OX crossover followed by the local search procedure. The method terminates when  $I_{\text{MAX}}$  individual generations have been performed without improvement, or a time limit  $T_{\text{MAX}}$  is reached.

Both methods rely on the local search described in Section 4.4 for solution improvement, considering possible infeasible solutions, followed by one iteration of ejection chains (construction of the ejection graph and resolution). Mode decisions are optimally taken within route evaluations, such that the classical Split algorithm is also implicitly extended into a variant that optimizes the mode choices in  $\mathcal{O}(n^2)$ , as in Ramdane-Cherif (2002) and Wøhlk (2004). Each metaheuristic is used with its original parameter setting:  $(n_I, n_C) = (100, 50)$  for ILS and  $(\mu^{\text{EL}}, \mu^{\text{MIN}}, \mu^{\text{GEN}}) = (12, 25, 40)$  for UHGS. The termination criterion is set to  $(n_P = 5, T_{\text{MAX}} = 1\text{h})$  for ILS and  $(I_{\text{MAX}} = 20, 000, T_{\text{MAX}} = 1\text{h})$  for UHGS to produce results in a time comparable to that of other algorithms.

### 5.2. Comparison with previous literature

The CARP and MCGRP literature contains several well-known sets of benchmark instances, and some of the largest, such as EGL, DI-NEARP, and CBMix, are still challenging for heuristics. We use these instances to assess the performance of the new metaheuristics with extended neighborhoods in comparison with previous algorithms. Table 2 gives the instance characteristics; the files can be found at <http://logistik.bwl.uni-mainz.de/benchmarks.php>, <http://www.sintef.no/Projectweb/TOP/nearp/>.

We ran the ILS and UHGS with extended neighborhoods ten times for each instance with different random seeds, on a single Xeon 3.07 GHz CPU with 16 GB of RAM. Note that some previous CARP algorithms were calibrated for limited groups of instances and frequently tested with a single run, leading to possible *over-tuning* effects. To avoid this, we 1) report the average results from several runs, 2) consider all available nontrivial sets of instances, 3) use new seeds for the final experiments, and 4) disregard any benchmark-set-specific parameter calibration. Finally, some previous algorithms terminate when a known optimal solution is reached. This approach is dependent on exogenous information, leading to inaccurate CPU time comparisons. Hence, we do not trigger the termination of ILS and UHGS when reaching a known optimal solution.

Table 3 gives a summary of the results. It lists for each set of instances the results of the best three methods (selected independently for each set) in terms of the average percentage gap,

**Table 2** Benchmark instances

#	Reference	$ N_R $	$ E_R $	$ A_R $	$n$	Specificities
<b>CARP</b>						
GDB	(23) Golden et al. (1983)	0	[11,55]	0	[11,55]	Random graphs; Only required edges
VAL	(34) Benavent et al. (1992)	0	[39,97]	0	[39,97]	Random graphs; Only required edges
BMCV	(100) Beullens et al. (2003)	0	[28,121]	0	[28,121]	Intercity road network in Flanders
EGL	(24) Li and Eglese (1996)	0	[51,190]	0	[51,190]	Winter-gritting application in Lancashire
EGL-L	(10) Brandão and E. (2008)	0	[347,375]	0	[347,375]	Larger winter-gritting application
<b>MCGRP</b>						
MGGDB	(138) Bosco et al. (2012)	[3,16]	[1,9]	[4,31]	[8,48]	From CARP instances GDB
MGVAL	(210) Bosco et al. (2012)	[7,46]	[6,33]	[12,79]	[36,129]	From CARP instances VAL
CBMix	(23) Prins and B. (2005)	[0,93]	[0,94]	[0,149]	[20,212]	Randomly generated planar networks
BHW	(20) Bach et al. (2013)	[4,50]	[0,51]	[7,380]	[20,410]	From CARP instances GDB, VAL, & EGL
DI-NEARP	(24) Bach et al. (2013)	[120,347]	[120,486]	0	[240,833]	Newspaper and media product distribution

computed as  $100(z - z_{\text{BKS}})/z_{\text{BKS}}$ , where  $z$  is the solution value obtained by the method and  $z_{\text{BKS}}$  is the best known solution (BKS) value for the instance. This measure is an estimate of the solution quality of a single run. The table also indicates, where available, the gap associated with the best solution of several runs, in which case the number of runs is indicated; the average CPU time “T” from the start of the method to the end (in minutes); the average time “T\*” to attain the final solution of the run; and the type of processor used. The new methods are highlighted on a grey background, and the best method, in terms of average gap, is highlighted in bold. The acronyms for the algorithms are listed in Table 4. Finally, the detailed results for each instance are given in Tables EC.1 to EC.15, available in the electronic companion of this article, and at <https://w1.cirreлт.ca/~vidalt/en/VRP-resources.html>.

Table 3 gives an overview of the good performance of the two proposed methods. The new neighborhoods lead to better solutions, attained by either ILS or UHGS, that have not been achieved until now with simpler neighborhoods. The inclusion of more advanced diversification principles, as in UHGS, helps to improve the solution quality even further: UHGS outperforms the current approaches by 0.523% and 1.142% on the largest instances (CBMix and EGL-L), and it outperforms ILS by up to 0.239%. Some of the BKSs for large instances (see Table EC.6) have been improved by up to 2.275%, which is a large difference given the research effort devoted to these problems. The proposed methods also produce solutions of consistent quality, with the average standard deviation for each set ranging from 0.000% to 0.243%. Overall, the impact of the extended neighborhoods and that of using UHGS rather than ILS are greater for larger instances.

We conducted paired-sample Wilcoxon tests to compare the performance of the proposed UHGS to other algorithms in terms of the gap for each instance set. The results of this analysis are illustrated in the boxplots in Figures 6 and 7, which give the gaps for the larger sets, as well as the p-values and significant effects detected by the tests. UHGS performs significantly better ( $p < 0.003$  in all cases) than the other methods on the large test sets (EGL, EGL-L, CbMix, BHW, and DI-NEARP). From these experiments, we also observed that the smaller sets have lost their discriminating power, since recent algorithms find the best solutions for the vast majority of instances, leading to many inconclusive observations. We thus encourage future contributors to investigate the performance of heuristics on the larger instances.

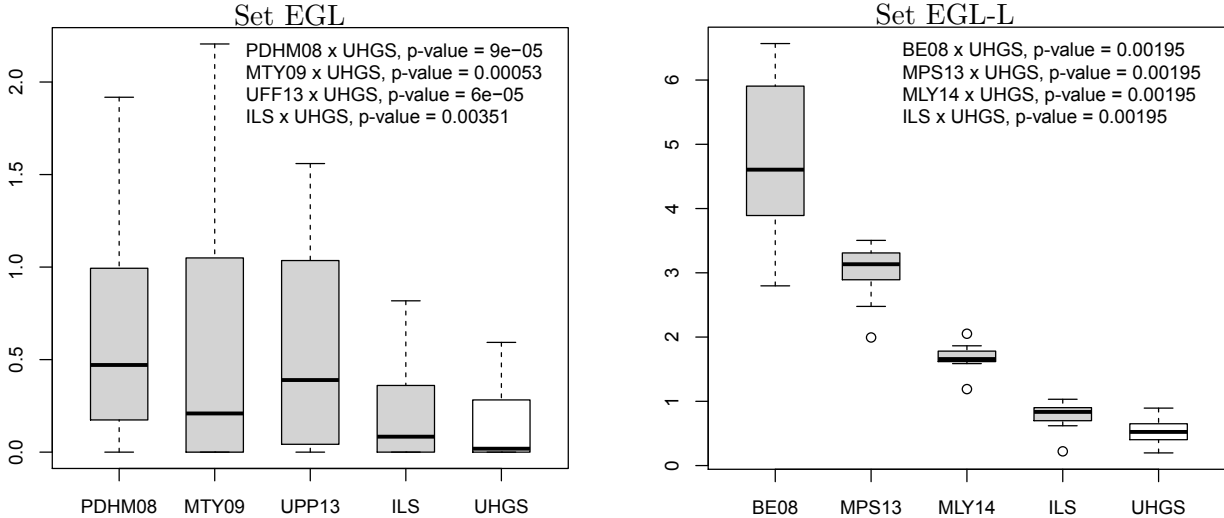
The CPU time of the proposed algorithms ranges from a fraction of a second on the smaller instances to 60 minutes for problems with 833 services. In the MCGRP experiments, this CPU time is smaller than other methods, which were tested on processors of a similar generation. For the CARP, some comparisons involve older processors. Since different CPU speed conversion techniques can give very different results, we opted to present the raw time data and processor information, letting the readers choose their preferred approach. Figure 8 displays the computational effort of

**Table 3** Results for CARP and MCGRP benchmark instances.

Variant	Bench.	$n$	Author	Runs	Avg.	Best	T	T*	CPU
CARP	GDB	[11,55]	TMY09	30	0.009%	0.000%	0.11	—	Xe 2.0G
			<b>BMCV03</b>	1	<b>0.000%</b>	—	—	0.03	P-II 500M
			<b>MTY09</b>	1	<b>0.000%</b>	—	—	0.01	Xe 2.0G
			<b>ILS</b>	10	<b>0.000%</b>	<b>0.000%</b>	0.30	< 0.01	Xe 3.07G
			<b>UHGS</b>	10	<b>0.000%</b>	<b>0.000%</b>	0.31	< 0.01	Xe 3.07G
	VAL	[39,97]	MTY09	1	0.142%	—	—	0.11	Xe 2.0G
			LPR01	1	0.126%	—	2.00	—	P-III 500M
			BMCV03	1	0.060%	—	—	1.36	P-II 500M
			<b>ILS</b>	10	0.044%	0.021%	1.10	0.07	Xe 3.07G
			<b>UHGS</b>	10	<b>0.041%</b>	<b>0.013%</b>	1.13	0.07	Xe 3.07G
	BMCV	[28,121]	BE08	1	0.158%	—	—	1.08	P-M 1.4G
			MTY09	1	0.075%	—	—	0.35	Xe 2.0G
			BMCV03	1	0.038%	—	2.57	—	P-II 450M
			<b>ILS</b>	10	0.020%	0.000%	1.26	0.10	Xe 3.07G
			<b>UHGS</b>	10	<b>0.013%</b>	<b>0.003%</b>	1.19	0.09	Xe 3.07G
	EGL	[51,190]	PDHM08	10	0.625%	—	30.0	8.39	P-IV 3.6G
			UFF13	15	0.562%	0.207%	13.3	—	I4 3.0G
			MTY09	1	0.555%	—	—	2.10	Xe 2.0G
			<b>ILS</b>	10	0.209%	0.088%	3.45	1.47	Xe 3.07G
			<b>UHGS</b>	10	<b>0.141%</b>	<b>0.049%</b>	5.59	3.46	Xe 3.07G
EGL-L	[347,375]	BE08	1	4.749%	—	—	17.0	P-M 1.4G	
		MPS13	10	3.018%	2.591%	20.7	—	I5 3.2G	
		MLY14	30	1.671%	0.962%	33.4	—	I7 3.4G	
		<b>ILS</b>	10	0.768%	0.468%	28.4	15.8	Xe 3.07G	
		<b>UHGS</b>	10	<b>0.529%</b>	<b>0.206%</b>	46.1	36.5	Xe 3.07G	
MCGRP	MGGDB	[8,48]	BLMV14	1	1.342%	—	0.31	—	Xe 3.0G
			DHDI14	1	0.018%	—	60.0	0.86	CPU 3G
			<b>ILS</b>	10	<b>0.006%</b>	<b>0.000%</b>	0.22	0.01	Xe 3.07G
			<b>UHGS</b>	10	<b>0.006%</b>	<b>0.000%</b>	0.27	0.01	Xe 3.07G
	MGVAL	[36,129]	BLMV14	1	2.621%	—	16.7	—	Xe 3.0G
			DHDI14	1	0.072%	—	60.0	3.69	CPU 3G
			<b>ILS</b>	10	0.047%	0.015%	1.41	0.15	Xe 3.07G
			<b>UHGS</b>	10	<b>0.047%</b>	<b>0.013%</b>	1.65	0.20	Xe 3.07G
	CBMix	[20,212]	HKSG12	2	—	3.083%	120	56.9	CPU 3G
			BLMV14	1	2.705%	—	44.7	—	Xe 3.0G
			DHDI14	1	0.891%	—	60.0	19.6	CPU 3G
			<b>ILS</b>	10	0.574%	0.287%	3.13	1.54	Xe 3.07G
			<b>UHGS</b>	10	<b>0.361%</b>	<b>0.088%</b>	5.09	3.09	Xe 3.07G
	BHW	[20,410]	HKSG12	2	—	1.976%	120	60.1	CPU 3G
			DHDI14	1	0.581%	—	60.0	21.4	CPU 3G
			<b>ILS</b>	10	0.338%	0.126%	6.91	4.00	Xe 3.07G
			<b>UHGS</b>	10	<b>0.230%</b>	<b>0.084%</b>	9.95	7.12	Xe 3.07G
	DI-NEARP	[240,833]	HKSG12	2	—	1.640%	120	93.0	CPU 3G
			DHDI14	1	0.537%	—	60.0	36.3	CPU 3G
			<b>ILS</b>	10	0.159%	0.074%	58.2	23.6	Xe 3.07G
<b>UHGS</b>			10	<b>0.106%</b>	<b>0.041%</b>	53.5	30.8	Xe 3.07G	

**Table 4** Current state-of-the-art methods for the considered benchmark instances, and their acronyms

BCS10	Benavent et al. (2010)	HKSG12	Hasle et al. (2012)	MPY11	Mei et al. (2011)
BE08	Brandão and Eglese (2008)	KY10	Kansou and Yassine (2010)	MTY09	Mei et al. (2009)
BLMV14	Bosco et al. (2014)	LPR01	Lacomme et al. (2001)	PDHM08	Polacek et al. (2008)
BMCV03	Beullens et al. (2003)	LPR05	Lacomme et al. (2005)	TMY09	Tang et al. (2009)
CLP06	Chu et al. (2006)	MLY14	Mei et al. (2014)	UFF13	Usberti et al. (2013)
DHDI14	Dell’Amico et al. (2016)	MPS13	Martinelli et al. (2013)		

**Figure 6** Boxplots of the percentage gap of recent algorithms on the CARP sets EGL and EGL-L. Paired-sample Wilcoxon tests between UHGS and the other methods are also reported. The boxplot associated with a method X is shaded in grey when the solutions of UHGS are of significantly better quality than those of X.

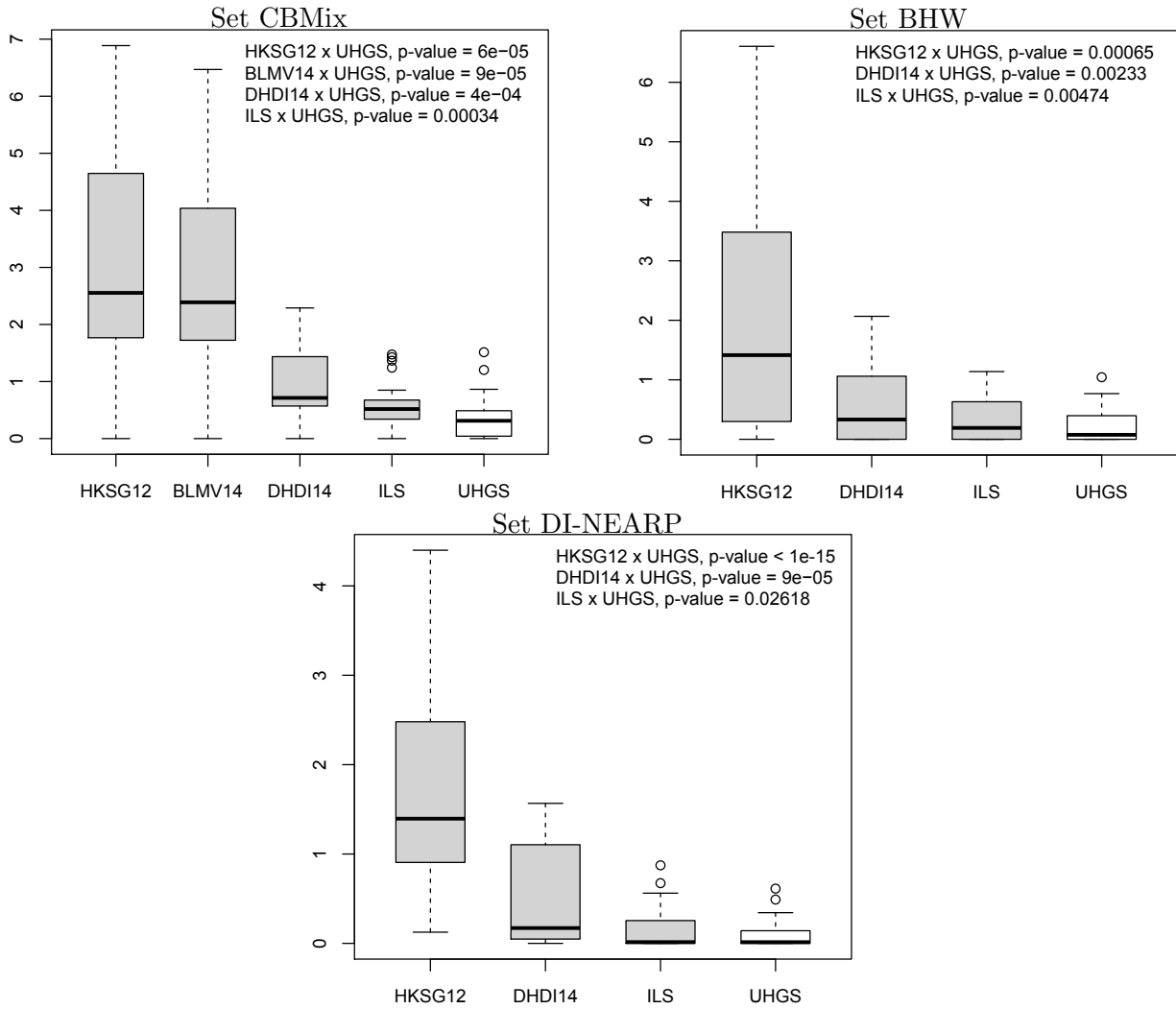
UHGS, in log-log scale, as a function of the number of customer services. The CPU time appears to rise in  $\mathcal{O}(n^2)$ . Measuring the respective share of computational effort on the EGL instances, we observed that the local search and ejection chains use 60% and 13% of the CPU time of UHGS, respectively. Within the LS, the preprocessing phases use 13% of the time, while the lower bounds and exact move evaluations consume 34% and 6%, respectively.

On the CARP benchmark sets, 187 of 191 BKS have been matched or improved, and 18 of these solutions have been strictly improved. Moreover, 153 solutions out of the 155 known optimal solutions were found. For the last two instances, val9D and val10D, the results are only one unit of distance away from the optimum. For the MCGRP, 408 of 409 BKS have been matched or improved, including 80 new BKS. All 217 known optimal solutions were found. The detailed results are reported in Tables EC.1 to EC.17.

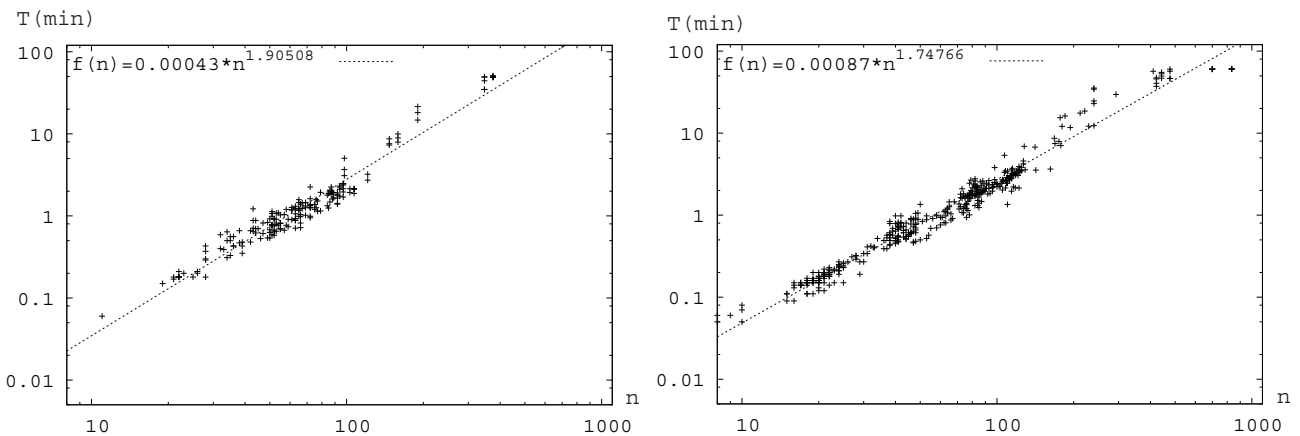
Finally, to complement the experimental comparisons, we considered three additional problems variants: with multiple delivery periods, multiple depots, and the min-max k-vehicles windy rural postman problem. The UHGS with extended neighborhoods also produced solutions of high quality for these problems. All BKS were either retrieved or improved, and for some PCARP instances, the improvements over previous BKS reached up to 10%. These additional experiments, summarized in Table 5, are thoroughly documented in Tables EC.18 to EC.25 of the E-companion.

### 5.3. Direct approach or problem transformation

The previous sections have investigated whether methods using extended neighborhoods – with optimal choices of service orientations – can outperform methods based on more traditional neighborhoods. This section analyzes whether relying on a problem reduction from CARP to CVRP (Baldacci and Maniezzo 2006) with a classical routing metaheuristic can be profitable. The reduction increases the number of services by a factor of two. Half of the edges of a CVRP solution, with a large fixed negative cost, directly determine the service orientations in the associated CARP solution.



**Figure 7** Boxplots of the percentage gap of recent algorithms on the MCGRP sets CBMix, BHW, and DI-NEARP. Same conventions as in Figure 6.



**Figure 8** Growth of the CPU time of UHGS as a function of the number of services, for the CARP instances (left figure) and MCGRP instances (right figure). Log-log scale. A linear fit, with a least square regression, has been performed on the sample after logarithmic transformation.



**Table 5** Results of UHGS on other CARP variants

Variant	Bench.	$n$	Author	Runs	Avg.	Best	T	T*	CPU
PCARP	PGDB	[65,165]	LPR05	1	9.448%	—	12.5	—	P-IV 1.4G
			CLP06	1	7.741%	—	1.86	—	P-IV 2.4G
			MPY11	30	3.900%	1.951%	0.20	—	Xe 2.0G
			UHGS <sup>†</sup>	10	0.730%	0.217%	0.14	0.09	Xe 3.07G
	<b>UHGS</b>	10	<b>0.256%</b>	<b>0.071%</b>	0.91	0.41	Xe 3.07G		
	PVAL	[94,300]	CLP06	1	16.494%	—	7.38	—	P-IV 2.4G
			MPY11	30	8.691%	6.317%	0.87	—	Xe 2.0G
			UHGS <sup>†</sup>	10	1.614%	0.721%	0.82	0.61	Xe 3.07G
<b>UHGS</b>			10	<b>0.636%</b>	<b>0.161%</b>	4.91	3.15	Xe 3.07G	
MDCARP	GDB	[8,48]	KY10	1	2.041%	—	0.02	—	P-IV 1.4G
			UHGS <sup>†</sup>	10	0.296%	0.104%	0.01	0.01	Xe 3.07G
			<b>UHGS</b>	10	<b>0.017%</b>	<b>0.000%</b>	0.37	0.04	Xe 3.07G
MM-kWRPP	2V	[7,78]	BCS10	1	0.103%	—	0.94	—	I2 2.4G
			<b>UHGS</b>	10	<b>0.008%</b>	<b>0.002%</b>	0.18	0.07	Xe 3.07G
	3V	[7,78]	BCS10	1	0.230%	—	0.41	—	I2 2.4G
			<b>UHGS</b>	10	<b>0.008%</b>	<b>0.000%</b>	0.18	0.07	Xe 3.07G
	4V	[7,78]	BCS10	1	0.303%	—	0.29	—	I2 2.4G
			<b>UHGS</b>	10	<b>0.014%</b>	<b>0.000%</b>	0.18	0.06	Xe 3.07G
	5V	[7,78]	BCS10	1	0.392%	—	0.24	—	I2 2.4G
			<b>UHGS</b>	10	<b>0.021%</b>	<b>0.000%</b>	0.19	0.07	Xe 3.07G

†: A shorter termination criteria has been used (Section EC.3).

We thus applied the same ILS and UHGS to the transformed instances, now using a classical move evaluation for the CVRP. Table 6 compares the results of these approaches, ILS<sub>CVRP</sub> and UHGS<sub>CVRP</sub>, with those of the proposed direct methods for the CARP.

**Table 6** Results of ILS<sub>CVRP</sub> and UHGS<sub>CVRP</sub> on the transformed CVRP. Comparison with the direct approaches.

	Gap(%)		T(min)		Gap(%)		T(min)		
	ILS	ILS <sub>CVRP</sub>	ILS	ILS <sub>CVRP</sub>	UHGS	UHGS <sub>CVRP</sub>	UHGS	UHGS <sub>CVRP</sub>	
GDB	0.000%	0.000%	0.30	0.59	GDB	0.000%	0.000%	0.31	0.72
VAL	0.044%	0.061%	1.10	2.39	VAL	0.041%	0.048%	1.13	2.98
BMCV	0.020%	0.047%	1.26	2.79	BMCV	0.013%	0.017%	1.19	3.02
EGL	0.209%	0.346%	3.45	8.50	EGL	0.141%	0.202%	5.59	12.65
EGL-L	0.768%	1.478%	28.4	60.0	EGL-L	0.529%	1.067%	46.1	59.7

Table 6 shows that a transformation into CVRP leads to solutions of lower quality, for a twofold increase in CPU time. The impact is stronger for larger instances, for which significant gap differences can be observed, e.g., 1.067% versus 0.529% for UHGS<sub>CVRP</sub> and 1.478% versus 0.768% for ILS<sub>CVRP</sub> on set EGL-L. The increased time is a natural consequence of the larger number of services resulting from the transformation. Hence, a direct approach with the extended neighborhoods appears to be a much better alternative than a problem transformation and CVRP resolution.

#### 5.4. Problem settings with turn penalties

Finally, a significant contribution of the proposed methodology is its ability to efficiently handle CARP and MCGRP with turn penalties. These problem classes have been little studied in the

operations research literature, and no benchmark instances for the direct problem formulation subsist nowadays. To fill this gap, we generated two benchmark sets.

The first set, CMMS, extends the 21 mixed rural postman problem instances with turn penalties from Table 3 of Corberán et al. (2002). No services on nodes are required, leading to instances of CARP with turn penalties. The demand quantities were selected with uniform probability in  $[5, 50]$ , and the vehicle capacities were scaled so as to satisfy 10 or 20 services on average, leading to a total of 42 instances based on random graphs. U-turns are not allowed.

To progress toward more realistic instances, we also extended the DI-NEARP instances of Bach et al. (2013) by adding turn penalties, leading to a set of 28 instances named DI-TP. These instances include services on nodes and edges. They are larger, and they correspond to an application of newspaper and media product delivery in the Nordic countries. Although the graph structure and the edge distances were given, this information was insufficient to identify left or right turns. Thus, we had to reconstruct a plausible planar layout for each instance. This was done with the FM<sup>3</sup> algorithm of Hachul and Jünger (2005).

The turn penalties were produced using the following rules. Let  $T_j$  be the sets of turns at a node  $j$ , ordered by increasing polar angle,  $T_j = \{(i, j, k) \mid i \neq k, (i, j) \in E \text{ and } (j, k) \in E\}$ .

- If  $|T_j| = 1$ , then the unique turn  $x \in T_j$  is not penalized:  $p_x = 0$ .
- Otherwise, if  $|T_j| \geq 2$ ,
  1. The turn  $\bar{x}$  with angle closest to  $180^\circ$  is the *straight* direction. It receives a penalty  $p_{\bar{x}} = \gamma$ .
  2. Any turn  $x$  with smaller angle than  $\bar{x}$  is a *right turn*, and receives a penalty  $p_x = 0$ .
  3. Any turn  $x$  with greater angle than  $\bar{x}$  is a *left turn*, and receives a penalty  $p_x = 3\gamma$ .
- The penalty of any U-turn  $x = (i, j, k)$  such that  $i = k$  is set to  $p_x = 5\gamma$ .

The coefficient  $\gamma$  was set, for each instance, to 10% of the average edge length. With this value, the turn costs amount to approximately 25% of the total objective function, a realistic estimate in light of the study of Nielsen et al. (1998).

We investigated the performance of the proposed methods on these sets of instances, also diminishing the population size by two to allow for a faster convergence. Since no solutions are available in the literature, we compare the results obtained with a termination criterion of ( $I_{\text{MAX}} = 10,000$ ;  $T_{\text{MAX}} = 1\text{h}$ ) and ( $I_{\text{MAX}} = 20,000$ ;  $T_{\text{MAX}} = 2\text{h}$ ), to observe to what extent the solution quality could be improved with additional computational effort. A small standard deviation in solution quality and a small gap with respect to the best solution of several runs are usually good indications of performance.

Tables EC.16 and EC.17 give the main characteristics of the new instances and the results of these experiments. The gap and standard deviation over 10 runs remain moderate: 0.526% for UHGS (with  $\sigma = 0.227\%$ ), and 0.626% for ILS (with  $\sigma = 0.249\%$ ), over the whole instance set. The best solutions of the longer runs are, on average, 0.184% better on set CMMS and 0.130% better on set DI-TP. This gap is larger than that for MCGRP instances without turn penalties.

To study further the impact of the turn penalties on the algorithm performance and the structure of the solutions, we conducted additional test runs on the 14 DI-TP instances with vehicle capacities  $Q \in \{2000, 4000\}$ , multiplying the penalties by a factor  $f_{\text{TP}} \in \{0, 0.25, 0.5, 1, 2, 5, 10\}$ . The average results for each level of  $\gamma$  are given in Table 7, while Figure 9 displays boxplots of the percentage gap of UHGS as a function of  $f_{\text{TP}}$  as well as the number of U-turns, left turns, and right turns observed in the solutions.

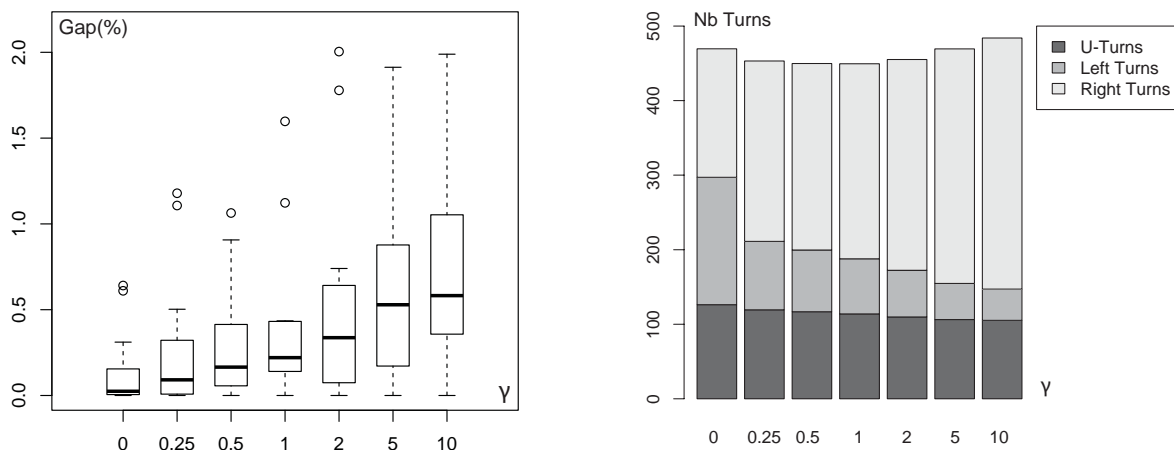
We observe a correlation between the weight of the turn penalties and the difficulty of the resulting instances, as reflected by a slightly larger gap (up to 0.752%) and standard deviation (up to  $\sigma = 0.366\%$ ) for higher values of  $f_{\text{TP}}$ . High penalties lead to more deceptive search spaces, possibly because the quality of a sequence of services now depends more strongly on the route context, i.e., the previous services and their impact on the best choices of arrival direction.

Higher turn penalty values lead to significantly fewer left turns in the solutions: from 170.85 turns on average for  $f_{\text{TP}} = 0$  to 73.91 for  $f_{\text{TP}} = 1$  and 42.01 for  $f_{\text{TP}} = 10$ . The number of U-turns also decreases until it reaches a plateau, due to some nodes with an out-degree of 1, where a U-turn may

**Table 7** Impact of the scale of the turn penalties on algorithm performance and solution characteristics.

$\gamma$	Gap (%)	T	Cost	Distance	No. Turns			
					U-turns	Left	Right	All
0	0.141%	50.68	25076.61	25076.61	126.24	170.85	172.35	469.44
0.25	0.280%	51.32	27500.70	25164.44	119.40	91.72	241.98	453.10
0.5	0.281%	51.65	29806.22	25250.74	116.79	82.77	250.17	449.73
1	0.373%	51.74	34339.29	25451.40	113.87	73.91	261.63	449.41
2	0.511%	51.77	43103.49	25986.19	109.84	62.54	282.69	455.06
5	0.607%	51.90	68258.91	27243.48	106.31	48.52	314.51	469.34
10	0.752%	51.92	109011.41	28534.13	105.23	42.01	336.76	484.00

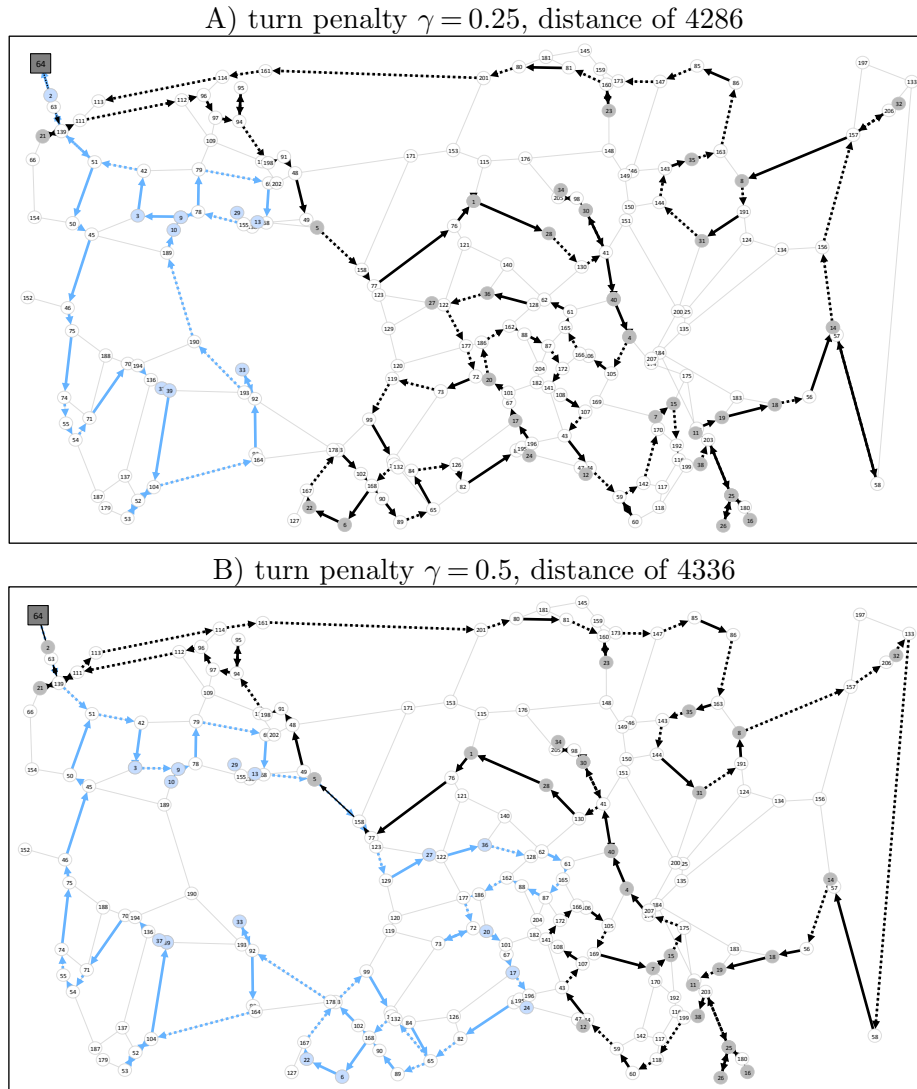
be unavoidable. Decreasing the number of turns naturally comes at the expense of a greater total distance. As Table 7 shows, moderate penalty values below  $\gamma = 1$  significantly reduce left turns and U-turns and lead to an increase of just 2% in the distance. This compromise is likely to be acceptable in practice. Finally, as illustrated in Figure 10, significant changes in the structure of the solutions occur as the turn penalties grow. The goal of minimizing left turns leads to clockwise cycles in some areas of the solutions. Despite these changes in solution structure, the proposed metaheuristic with extended neighborhoods still produces high-quality solutions.

**Figure 9** On the left, percentage gap of UHGS for different values of the turn penalties. On the right, average number of U-turns, left turns, and right turns in the solutions as a function of the turn penalties.

## 6. Conclusions and Future Research

In this article, we have investigated a family of extended neighborhoods for arc routing problems, which involve service relocations and exchanges with combined optimal mode choices on the routes. We have shown that a complete exploration of this extended neighborhood can be done in  $\mathcal{O}(n^2)$ , the same complexity as a classical neighborhood search without mode optimization. Moreover, as the new approach relegates service mode choices in a dynamic programming subproblem, it guarantees their optimality and helps to concentrate the heuristic search on a more limited decision set. We have generalized this methodology to multiple arc routing variants, extended it into large-scale neighborhood searches such as ejection chains, and integrated it into two classical metaheuristic frameworks. Our experiments led to solutions of high quality on 18 benchmark sets for these problems, with a total of 1528 instances. We also conducted experiments on the CARP-TP and MCGRP-TP, which are uncommonly challenging and critical for practical applications.

This work will help to build more connections between the arc routing and vehicle routing communities and take research into heuristics in new directions. Current research is often excessively



**Figure 10** Comparison of solutions for the smallest DI-TP instance, n80-Q4k, when doubling the turn penalties.

concerned with the design of *better* high-level metaheuristic frameworks rather than the study of alternative neighborhood structures. Significant future breakthroughs may arise from a better understanding of problem structure and *neighborhood search* – as attempted here – rather than brute-force *neighborhood enumeration*.

Future research into arc routing could also extend these neighborhoods to an even wider class of problems with different resources, such as time, duration, workload, and energy, as well as possible time-dependent resource consumptions (e.g., congestion aspects in urban environments). In the presence of additional resources and constraints, move evaluations would be formulated as resource-constrained shortest paths problems, and additional move lower bounds might be needed. Finally, similar neighborhood decompositions and dynamic-programming-based move evaluations should be investigated for other difficult combinatorial optimization problems, such as scheduling, allocation, and packing problems.

### Revision Note

An article of Chen et al. (2016) appeared online during our final revision steps. It presents an hybrid GA with tabu search, merging moves, and infeasible descents as an intensification procedure. This approach also

produces solutions of high quality, albeit in higher CPU time, via more dedicated metaheuristic strategies rather than extended neighborhoods. This leads to two state-of-the-art methods, which are both interesting in their own right since they bring significant improvements to two symbiotic aspects of the search: the low-level neighborhoods (this paper), and the higher level metaheuristic strategies (the mentioned article). Future research should gather key elements from these methods to progress on a new generation of simple and efficient heuristics.

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Thibaut Vidal is professor at the department of computer science of the Pontifical Catholic University of Rio de Janeiro, Brazil. His research interests include combinatorial optimization, metaheuristics, integer and dynamic programming, with applications to logistics and distribution systems design, vehicle routing and scheduling, resource allocation and signal processing.

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# Node, Edge, Arc Routing and Turn Penalties: Multiple Problems – One Neighborhood Extension

## *Electronic Companion*

This E-Companion provides the detailed results, on individual instances, for all the benchmark sets considered in this paper.

### EC.1. Conventions on fleet size limits

In the usual CARP definition, no limit is imposed on the number of vehicles used (see, e.g., Corberán and Laporte 2014). We use this convention when testing the algorithms. Note that all the solutions found by ILS and UHGS for CARP instances use the minimum number of vehicles, so these solutions remain valid in the presence of the fleet size constraint.

The situation is different for problem extensions with services on nodes, edges, and arcs. The NEARP, as defined by Prins and Bouchenoua (2005), does not include a limit on the number of vehicles; whereas the MCGRP (the same problem), as defined by Bosco et al. (2012), includes a maximum fleet size limit. We respected the fleet size conventions specified by the authors for each benchmark set. Note that for mggdb-0.xx-13 and mgval-0.50-1C, slightly better solutions may be obtained if the fleet size is relaxed. Finally, the solution of mgval-0.50-1C from Dell’Amico et al. (2016) appears to include an extra vehicle. For this reason, it was not considered a BKS.

### EC.2. Format of the tables

Tables EC.1 to EC.17 present the detailed results for ILS and UHGS with extended neighborhoods as well as other state-of-the-art methods for the CARP and MCGRP benchmark instances, with possible turn penalties. The first group of columns presents the instance identifiers and the number of services on nodes  $|N_R|$ , edges  $|E_R|$ , and arcs  $|A_R|$ . The next group of columns presents the results. When available, we report both average and best results over several runs (the number of runs is specified in the table headings). The best solution among the methods is highlighted in bold. In addition, the two last columns indicate the previous and new BKS. The new BKS were collected from all the runs, including the preliminary tests. New BKS are underlined. Known optimal solutions from the literature, gathered at <http://logistik.bwl.uni-mainz.de/benchmarks.php> and <http://www.sintef.no/Projectweb/TOP/nearp/>, are indicated with an asterisk. For three benchmark instances (BMCV-E9, mgval-0.25-4D, and mgval-0.25-10D) our new BKS matches the best known lower bound from Bode and Irnich (2015) and Irnich et al. (2015), so these three solutions are also optimal. Finally, the last lines of the tables report average measures over the sets of instances: the CPU time for each method, the time to reach the final solution, the average percentage gap relative to the new BKS, and the processor.

### EC.3. Additional experiments on CARP variants

To complement these experimental analyses and provide more elements of comparison with current and future state-of-the-art methods for arc routing problems, we considered three additional problem settings, which are the arc-routing counterparts of important problem variants addressed by the hybrid genetic search of Vidal et al. (2012), used in this paper. As such, we can apply the same metaheuristic with the generalized move evaluations.

- In the **periodic capacitated arc routing problem (PCARP)**, the deliveries are performed in several time periods. The objective is to choose, for each service, a combination of periods in a set of allowed patterns and to design the routing plan for each period. The hybrid genetic search of Vidal et al. (2012) already contains an additional local search phase (called pattern improvement – PI) and a crossover to explore alternative customer-pattern assignment choices. Our only adaptation concerns the evaluations of the routes, which now involve an optimal choice of service orientations and serve as a building block for all other search components. Moreover, the classical objective for the PCARP minimizes the fleet size, and then the travel distance. This is done by iteratively decrementing the fleet size and performing shorter runs of UHGS ( $I_{\text{MAX}} = 5,000$ ) until no solution is found. Then, a final run is performed ( $I_{\text{MAX}} = 10,000$ ) with the smallest feasible fleet size. Tables EC.18 and EC.19 report the performance of UHGS on the benchmark instances of Lacomme et al. (2005) and Chu et al. (2006), in comparison with the current state-of-the-art methods for the PCARP: the improved periodic memetic algorithm of Lacomme et al. (2005) (LPR05), the scatter search of Chu et al. (2006) (CLP06), and the memetic algorithm with route merging of Mei et al. (2011) (MPY11). To allow a better comparison with MPY11, which produces results in very short CPU time, we report additional computational results of UHGS with a short termination criterion,  $I_{\text{MAX}} = 1,000$ .
- In the **multi-depot capacitated arc routing problem (MDCARP)**, the vehicles are stationed at several depot locations. The objective is to optimize the depot-to-service assignments and the routing plans. Despite the importance of the multi-depot attribute in the vehicle routing literature, few articles have investigated it in the arc routing domain. In Krushinsky and Van Woensel (2015), all services are requested on (oriented) arcs. This problem, however, can be reduced to an asymmetric MDVRP. Xing et al. (2010) consider an extension of MDCARP with prohibited turns and duration constraints. Finally, Kansou and Yassine (2010) address the MDCARP, reporting computational results for extensions of the GDB benchmark, obtained by considering depot locations at the vertices 1 and  $|V|$ . We followed this latter approach, also generalizing the VAL and EGL instances with  $d = 3$  and  $d = 4$  depots, respectively, located at the vertices 1 and  $x \lfloor \frac{|V|}{d-1} \rfloor$  for  $x \in \{1, \dots, d-1\}$ . The results for these three sets of instances are reported in Tables EC.20 to EC.22. To allow a better comparison with the algorithm of Kansou and Yassine (2010), on the set GDB, we report additional computational results with  $I_{\text{MAX}} = 500$ .
- Finally, in the **min-max k-vehicles windy rural postman problem (MM-kWRPP)**, the travel and service costs are asymmetric, and the objective is to minimize the distance of the longest route. We addressed this objective with UHGS by setting a distance constraint  $D$  for the routes, and iteratively setting  $D = D_S - 1$  when a feasible solution  $S$  is found, where  $D_S$  represents the maximum distance of a route in  $S$ . The termination criterion was set to  $I_{\text{MAX}} = 2,000$  at each iteration. Experiments were conducted on the benchmark of Benavent et al. (2009), which includes 24 classes of 6 instances, with 2 to 6 vehicles. The UHGS results are compared with those of the multi-start ILS of Benavent et al. (2010) (detailed solution values were obtained by contacting the authors), and with the BKS from subsequent works on exact methods, available at <http://www.uv.es/corberan/instancias.htm>. The total number of instances (720) is too large for a display of all the individual results. Hence, Tables EC.23 to EC.25 report detailed results for the larger instances of classes C20 to C24, as well as C16 to C19 for problems with six vehicles. For the remaining instances, the optimal solutions are known, and these solutions were found by UHGS on all runs. The average results, at the bottom of Table EC.25, are for the whole set of instances. Note that, for the 76 currently open instances, 71 upper bounds have been improved. This will be helpful for future research into exact methods.

## EC.4. Open source code library

Finally, to help reproducing and extending this work, we provide at <https://github.com/vidalthi/HGS-CARP> the source code of the proposed methods under the GPL 3.0 license as well as the data sets considered in this article.

**Table EC.1** Results for the CARP – GDB instances

Inst	$ E_R $	BMCV03	TMY09		MTY09	ILS				UHGS				BKS
		Single	Avg-30	Best-30	Single	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Opt
gdb1	22	<b>316</b>	<b>316.0</b>	<b>316</b>	<b>316</b>	<b>316.0</b>	<b>316</b>	0.14	0.00	<b>316.0</b>	<b>316</b>	0.18	0.00	316*
gdb2	26	<b>339</b>	<b>339.0</b>	<b>339</b>	<b>339</b>	<b>339.0</b>	<b>339</b>	0.19	0.00	<b>339.0</b>	<b>339</b>	0.21	0.00	339*
gdb3	22	<b>275</b>	<b>275.0</b>	<b>275</b>	<b>275</b>	<b>275.0</b>	<b>275</b>	0.15	0.00	<b>275.0</b>	<b>275</b>	0.18	0.00	275*
gdb4	19	<b>287</b>	<b>287.0</b>	<b>287</b>	<b>287</b>	<b>287.0</b>	<b>287</b>	0.12	0.00	<b>287.0</b>	<b>287</b>	0.15	0.00	287*
gdb5	26	<b>377</b>	<b>377.0</b>	<b>377</b>	<b>377</b>	<b>377.0</b>	<b>377</b>	0.18	0.00	<b>377.0</b>	<b>377</b>	0.20	0.00	377*
gdb6	22	<b>298</b>	<b>298.0</b>	<b>298</b>	<b>298</b>	<b>298.0</b>	<b>298</b>	0.14	0.00	<b>298.0</b>	<b>298</b>	0.18	0.00	298*
gdb7	22	<b>325</b>	<b>325.0</b>	<b>325</b>	<b>325</b>	<b>325.0</b>	<b>325</b>	0.14	0.00	<b>325.0</b>	<b>325</b>	0.18	0.00	325*
gdb8	46	<b>348</b>	348.7	<b>348</b>	<b>348</b>	<b>348.0</b>	<b>348</b>	0.58	0.05	<b>348.0</b>	<b>348</b>	0.53	0.01	348*
gdb9	51	<b>303</b>	<b>303.0</b>	<b>303</b>	<b>303</b>	<b>303.0</b>	<b>303</b>	0.59	0.02	<b>303.0</b>	<b>303</b>	0.59	0.02	303*
gdb10	25	<b>275</b>	<b>275.0</b>	<b>275</b>	<b>275</b>	<b>275.0</b>	<b>275</b>	0.20	0.00	<b>275.0</b>	<b>275</b>	0.18	0.00	275*
gdb11	45	<b>395</b>	<b>395.0</b>	<b>395</b>	<b>395</b>	<b>395.0</b>	<b>395</b>	0.62	0.00	<b>395.0</b>	<b>395</b>	0.70	0.00	395*
gdb12	23	<b>458</b>	<b>458.0</b>	<b>458</b>	<b>458</b>	<b>458.0</b>	<b>458</b>	0.16	0.00	<b>458.0</b>	<b>458</b>	0.20	0.00	458*
gdb13	28	<b>536</b>	<b>536.0</b>	<b>536</b>	<b>536</b>	<b>536.0</b>	<b>536</b>	0.28	0.01	<b>536.0</b>	<b>536</b>	0.30	0.01	536*
gdb14	21	<b>100</b>	<b>100.0</b>	<b>100</b>	<b>100</b>	<b>100.0</b>	<b>100</b>	0.14	0.00	<b>100.0</b>	<b>100</b>	0.18	0.00	100*
gdb15	21	<b>58</b>	<b>58.0</b>	<b>58</b>	<b>58</b>	<b>58.0</b>	<b>58</b>	0.16	0.00	<b>58.0</b>	<b>58</b>	0.17	0.00	58*
gdb16	28	<b>127</b>	<b>127.0</b>	<b>127</b>	<b>127</b>	<b>127.0</b>	<b>127</b>	0.28	0.00	<b>127.0</b>	<b>127</b>	0.29	0.00	127*
gdb17	28	<b>91</b>	<b>91.0</b>	<b>91</b>	<b>91</b>	<b>91.0</b>	<b>91</b>	0.20	0.00	<b>91.0</b>	<b>91</b>	0.18	0.00	91*
gdb18	36	<b>164</b>	<b>164.0</b>	<b>164</b>	<b>164</b>	<b>164.0</b>	<b>164</b>	0.43	0.00	<b>164.0</b>	<b>164</b>	0.43	0.00	164*
gdb19	11	<b>55</b>	<b>55.0</b>	<b>55</b>	<b>55</b>	<b>55.0</b>	<b>55</b>	0.05	0.00	<b>55.0</b>	<b>55</b>	0.06	0.00	55*
gdb20	22	<b>121</b>	<b>121.0</b>	<b>121</b>	<b>121</b>	<b>121.0</b>	<b>121</b>	0.23	0.00	<b>121.0</b>	<b>121</b>	0.21	0.00	121*
gdb21	33	<b>156</b>	<b>156.0</b>	<b>156</b>	<b>156</b>	<b>156.0</b>	<b>156</b>	0.38	0.00	<b>156.0</b>	<b>156</b>	0.39	0.00	156*
gdb22	44	<b>200</b>	<b>200.0</b>	<b>200</b>	<b>200</b>	<b>200.0</b>	<b>200</b>	0.62	0.01	<b>200.0</b>	<b>200</b>	0.62	0.01	200*
gdb23	55	<b>233</b>	<b>233.0</b>	<b>233</b>	<b>233</b>	<b>233.0</b>	<b>233</b>	0.81	0.02	<b>233.0</b>	<b>233</b>	0.81	0.02	233*
Gap (%)		0.000%	0.009%	0.000%	0.000%	0.000%	0.000%			0.000%	0.000%			
T(min)		—	0.11		—			0.30				0.31		
T*(min)		0.03	—		0.01			<0.01				<0.01		
CPU		P-II 500M	Xe 2.0G		Xe 2.0G			Xe 3.07G				Xe 3.07G		

Table EC.2 Results for the CARP – VAL instances

Inst	$ E_R $	LPR01	BMCV03	MTY09	ILS				UHGS				BKS
		Single	Single	Single	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Opt
1A	39	<b>173</b>	<b>173</b>	<b>173</b>	<b>173.0</b>	<b>173</b>	0.43	0.00	<b>173.0</b>	<b>173</b>	0.35	0.00	173*
1B	39	<b>173</b>	<b>173</b>	<b>173</b>	<b>173.0</b>	<b>173</b>	0.62	0.01	<b>173.0</b>	<b>173</b>	0.48	0.01	173*
1C	39	<b>245</b>	<b>245</b>	<b>245</b>	<b>245.0</b>	<b>245</b>	0.38	0.00	<b>245.0</b>	<b>245</b>	0.43	0.00	245*
2A	34	<b>227</b>	<b>227</b>	<b>227</b>	<b>227.0</b>	<b>227</b>	0.55	0.00	<b>227.0</b>	<b>227</b>	0.64	0.00	227*
2B	34	<b>259</b>	<b>259</b>	<b>259</b>	<b>259.0</b>	<b>259</b>	0.44	0.00	<b>259.0</b>	<b>259</b>	0.50	0.00	259*
2C	34	<b>457</b>	<b>457</b>	<b>457</b>	<b>457.0</b>	<b>457</b>	0.33	0.01	<b>457.0</b>	<b>457</b>	0.31	0.01	457*
3A	35	<b>81</b>	<b>81</b>	<b>81</b>	<b>81.0</b>	<b>81</b>	0.53	0.00	<b>81.0</b>	<b>81</b>	0.56	0.00	81*
3B	35	<b>87</b>	<b>87</b>	<b>87</b>	<b>87.0</b>	<b>87</b>	0.46	0.00	<b>87.0</b>	<b>87</b>	0.50	0.00	87*
3C	35	<b>138</b>	<b>138</b>	<b>138</b>	<b>138.0</b>	<b>138</b>	0.30	0.00	<b>138.0</b>	<b>138</b>	0.33	0.01	138*
4A	69	<b>400</b>	<b>400</b>	<b>400</b>	<b>400.0</b>	<b>400</b>	1.36	0.01	<b>400.0</b>	<b>400</b>	1.46	0.01	400*
4B	69	<b>412</b>	<b>412</b>	<b>412</b>	<b>412.0</b>	<b>412</b>	1.19	0.01	<b>412.0</b>	<b>412</b>	1.32	0.01	412*
4C	69	<b>428</b>	<b>428</b>	<b>428</b>	<b>428.0</b>	<b>428</b>	1.28	0.03	<b>428.0</b>	<b>428</b>	1.26	0.02	428*
4D	69	530	530	530	530.1	<b>528</b>	1.12	0.26	529.8	<b>528</b>	1.02	0.06	528*
5A	65	<b>423</b>	<b>423</b>	<b>423</b>	<b>423.0</b>	<b>423</b>	1.33	0.01	<b>423.0</b>	<b>423</b>	1.46	0.01	423*
5B	65	<b>446</b>	<b>446</b>	<b>446</b>	<b>446.0</b>	<b>446</b>	1.11	0.01	<b>446.0</b>	<b>446</b>	1.24	0.01	446*
5C	65	<b>474</b>	<b>474</b>	<b>474</b>	<b>474.0</b>	<b>474</b>	1.06	0.01	<b>474.0</b>	<b>474</b>	1.17	0.01	474*
5D	65	581	579	583	575.8	<b>575</b>	1.14	0.40	575.6	<b>575</b>	1.58	0.64	575*
6A	50	<b>223</b>	<b>223</b>	<b>223</b>	<b>223.0</b>	<b>223</b>	0.74	0.00	<b>223.0</b>	<b>223</b>	0.63	0.00	223*
6B	50	<b>233</b>	<b>233</b>	<b>233</b>	<b>233.0</b>	<b>233</b>	0.65	0.01	<b>233.0</b>	<b>233</b>	0.68	0.00	233*
6C	50	<b>317</b>	<b>317</b>	<b>317</b>	<b>317.0</b>	<b>317</b>	0.54	0.01	<b>317.0</b>	<b>317</b>	0.54	0.01	317*
7A	66	<b>279</b>	<b>279</b>	<b>279</b>	<b>279.0</b>	<b>279</b>	0.84	0.01	<b>279.0</b>	<b>279</b>	0.72	0.00	279*
7B	66	<b>283</b>	<b>283</b>	<b>283</b>	<b>283.0</b>	<b>283</b>	1.01	0.00	<b>283.0</b>	<b>283</b>	1.04	0.00	283*
7C	66	<b>334</b>	<b>334</b>	<b>334</b>	<b>334.0</b>	<b>334</b>	0.86	0.02	<b>334.0</b>	<b>334</b>	0.82	0.02	334*
8A	63	<b>386</b>	<b>386</b>	<b>386</b>	<b>386.0</b>	<b>386</b>	0.96	0.01	<b>386.0</b>	<b>386</b>	0.88	0.00	386*
8B	63	<b>395</b>	<b>395</b>	<b>395</b>	<b>395.0</b>	<b>395</b>	1.15	0.01	<b>395.0</b>	<b>395</b>	1.23	0.01	395*
8C	63	527	<b>521</b>	524	<b>521.0</b>	<b>521</b>	1.16	0.36	<b>521.0</b>	<b>521</b>	1.06	0.19	521*
9A	92	<b>323</b>	<b>323</b>	<b>323</b>	<b>323.0</b>	<b>323</b>	1.65	0.02	<b>323.0</b>	<b>323</b>	1.39	0.05	323*
9B	92	<b>326</b>	<b>326</b>	<b>326</b>	<b>326.0</b>	<b>326</b>	1.87	0.02	<b>326.0</b>	<b>326</b>	1.78	0.03	326*
9C	92	<b>332</b>	<b>332</b>	<b>332</b>	<b>332.0</b>	<b>332</b>	1.76	0.01	<b>332.0</b>	<b>332</b>	1.91	0.04	332*
9D	92	391	391	391	390.8	390	1.43	0.29	390.6	<b>389</b>	1.65	0.19	388*
10A	97	<b>428</b>	<b>428</b>	<b>428</b>	<b>428.0</b>	<b>428</b>	2.45	0.06	<b>428.0</b>	<b>428</b>	2.45	0.03	428*
10B	97	<b>436</b>	<b>436</b>	<b>436</b>	<b>436.0</b>	<b>436</b>	2.36	0.03	<b>436.0</b>	<b>436</b>	2.38	0.08	436*
10C	97	<b>446</b>	<b>446</b>	<b>446</b>	<b>446.0</b>	<b>446</b>	2.14	0.07	<b>446.0</b>	<b>446</b>	2.14	0.06	446*
10D	97	530	526	534	526.3	526	2.03	0.83	526.4	<b>526</b>	2.48	0.82	525*
Gap (%)		0.126%	0.060%	0.142%	0.044%	0.021%			0.041%	0.013%			
T(min)		2.00	—	—			1.10				1.13		
T*(min)		—	1.36	0.11				0.07				0.07	
CPU		P-III 500M	P-II 500M	Xe 2.0G			Xe 3.07G				Xe 3.07G		

Table EC.3 Results for the CARP – BMCV instances

Inst	$ E_R $	BMCV03	BE08	MTY09	ILS				UHGS				BKS	
		Single	Single	Single	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Old	New
C01	79	4150	4150	4150	4154.5	4150	1.33	0.33	4150.0	4150	1.15	0.08	4150*	4150
C02	53	3135	3135	3135	3135.0	3135	0.65	0.01	3135.0	3135	0.65	0.01	3135*	3135
C03	51	2575	2575	2575	2575.0	2575	0.85	0.03	2575.0	2575	0.78	0.09	2575*	2575
C04	72	3510	3510	3510	3510.0	3510	1.04	0.01	3510.0	3510	0.99	0.01	3510*	3510
C05	65	5370	5365	5365	5365.0	5365	0.99	0.08	5365.0	5365	0.99	0.18	5365*	5365
C06	51	2535	2535	2535	2535.0	2535	0.78	0.02	2535.0	2535	0.76	0.02	2535*	2535
C07	52	4075	4075	4075	4075.0	4075	0.61	0.01	4075.0	4075	0.58	0.01	4075*	4075
C08	63	4090	4090	4090	4090.0	4090	0.78	0.02	4090.0	4090	0.71	0.03	4090*	4090
C09	97	5265	5270	5270	5279.0	5260	1.41	0.28	5260.0	5260	1.70	0.43	5260	5260
C10	55	4720	4700	4700	4700.0	4700	0.81	0.05	4700.0	4700	0.69	0.08	4700*	4700
C11	94	4640	4640	4630	4634.0	4630	1.94	0.74	4636.0	4630	1.74	0.36	4630	4630
C12	72	4240	4240	4240	4240.0	4240	0.98	0.02	4240.0	4240	0.95	0.02	4240*	4240
C13	52	2955	2955	2955	2955.0	2955	0.77	0.02	2955.0	2955	0.67	0.02	2955*	2955
C14	57	4030	4030	4030	4030.0	4030	0.80	0.03	4030.0	4030	0.67	0.02	4030*	4030
C15	107	4940	4945	4940	4940.5	4940	2.32	0.57	4940.0	4940	1.88	0.36	4940	4940
C16	32	1475	1475	1475	1475.0	1475	0.45	0.00	1475.0	1475	0.40	0.00	1475*	1475
C17	42	3555	3555	3555	3555.0	3555	0.53	0.02	3555.0	3555	0.48	0.02	3555*	3555
C18	121	5645	5650	5660	5621.0	5605	2.54	0.80	5625.5	5620	3.22	1.37	5620	5605
C19	61	3115	3120	3115	3115.0	3115	0.97	0.02	3115.0	3115	0.90	0.07	3115*	3115
C20	53	2120	2120	2120	2120.0	2120	0.89	0.01	2120.0	2120	0.79	0.01	2120*	2120
C21	76	3970	3970	3970	3970.0	3970	1.27	0.02	3970.0	3970	1.26	0.02	3970*	3970
C22	43	2245	2245	2245	2245.0	2245	0.70	0.00	2245.0	2245	0.71	0.00	2245*	2245
C23	92	4085	4095	4095	4085.0	4085	1.79	0.18	4085.0	4085	1.42	0.14	4085	4085
C24	84	3400	3400	3400	3400.0	3400	1.51	0.10	3400.0	3400	1.33	0.05	3400*	3400
C25	38	2310	2310	2310	2310.0	2310	0.46	0.01	2310.0	2310	0.47	0.01	2310*	2310
D01	79	3215	3230	3230	3215.0	3215	1.95	0.32	3218.0	3215	1.94	0.44	3215*	3215
D02	53	2520	2520	2520	2520.0	2520	0.84	0.00	2520.0	2520	0.91	0.00	2520*	2520
D03	51	2065	2065	2065	2065.0	2065	1.03	0.00	2065.0	2065	0.99	0.00	2065*	2065
D04	72	2785	2785	2785	2785.0	2785	1.44	0.01	2785.0	2785	1.38	0.00	2785*	2785
D05	65	3935	3935	3935	3935.0	3935	1.14	0.01	3935.0	3935	1.21	0.01	3935*	3935
D06	51	2125	2125	2125	2125.0	2125	1.15	0.01	2125.0	2125	1.07	0.00	2125*	2125
D07	52	3115	3115	3115	3115.0	3115	0.92	0.03	3115.0	3115	0.77	0.02	3115*	3115
D08	63	3045	3045	3045	3045.0	3045	1.05	0.02	3045.0	3045	0.89	0.02	3045*	3045
D09	97	4120	4120	4120	4120.0	4120	1.87	0.02	4120.0	4120	1.96	0.03	4120*	4120
D10	55	3340	3340	3340	3340.0	3340	0.81	0.00	3340.0	3340	0.82	0.00	3340*	3340
D11	94	3755	3785	3755	3745.0	3745	2.72	0.59	3745.0	3745	2.24	0.24	3745*	3745
D12	72	3310	3310	3310	3310.0	3310	1.29	0.02	3310.0	3310	1.28	0.01	3310*	3310
D13	52	2535	2540	2535	2535.0	2535	1.09	0.05	2535.0	2535	0.92	0.03	2535*	2535
D14	57	3280	3290	3280	3280.0	3280	1.09	0.01	3280.0	3280	1.04	0.01	3280*	3280
D15	107	3990	4030	4000	3990.0	3990	2.65	0.22	3990.0	3990	2.16	0.17	3990*	3990
D16	32	1060	1060	1060	1060.0	1060	0.55	0.00	1060.0	1060	0.59	0.00	1060*	1060
D17	42	2620	2620	2620	2620.0	2620	0.65	0.00	2620.0	2620	0.66	0.00	2620*	2620
D18	121	4165	4165	4185	4165.0	4165	3.00	0.09	4165.0	4165	2.73	0.10	4165*	4165
D19	61	2400	2410	2400	2400.0	2400	1.33	0.00	2400.0	2400	1.27	0.00	2400*	2400
D20	53	1870	1870	1870	1870.0	1870	1.07	0.00	1870.0	1870	1.08	0.00	1870*	1870
D21	76	3050	3070	3055	3050.0	3050	1.95	0.31	3050.0	3050	1.58	0.07	3050	3050
D22	43	1865	1865	1865	1865.0	1865	1.17	0.00	1865.0	1865	1.22	0.00	1865*	1865
D23	92	3130	3130	3130	3130.0	3130	2.01	0.06	3130.0	3130	1.81	0.04	3130	3130
D24	84	2710	2710	2710	2710.0	2710	1.92	0.02	2710.0	2710	1.79	0.02	2710*	2710
D25	38	1815	1815	1815	1815.0	1815	0.63	0.00	1815.0	1815	0.66	0.00	1815*	1815

Table EC.4 Results for the CARP – BMCV instances (continued)

Inst	$ E_R $	BMCV03	BE08	MTY09	ILS				UHGS				BKS	
		Single	Single	Single	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Old	New
E01	85	4915	<b>4910</b>	<b>4910</b>	<b>4910.0</b>	<b>4910</b>	1.46	0.13	<b>4910.0</b>	<b>4910</b>	1.25	0.11	4910	4910
E02	58	<b>3990</b>	<b>3990</b>	<b>3990</b>	<b>3990.0</b>	<b>3990</b>	0.87	0.03	<b>3990.0</b>	<b>3990</b>	0.77	0.02	3990*	3990
E03	47	<b>2015</b>	<b>2015</b>	<b>2015</b>	<b>2015.0</b>	<b>2015</b>	0.69	0.01	<b>2015.0</b>	<b>2015</b>	0.61	0.01	2015*	2015
E04	77	<b>4155</b>	4160	<b>4155</b>	<b>4155.0</b>	<b>4155</b>	1.44	0.07	<b>4155.0</b>	<b>4155</b>	1.14	0.08	4155*	4155
E05	61	4595	<b>4585</b>	4610	4592.5	<b>4585</b>	1.20	0.56	<b>4585.0</b>	<b>4585</b>	1.19	0.37	4585*	4585
E06	43	<b>2055</b>	<b>2055</b>	<b>2055</b>	<b>2055.0</b>	<b>2055</b>	0.62	0.00	<b>2055.0</b>	<b>2055</b>	0.65	0.00	2055*	2055
E07	50	<b>4155</b>	<b>4155</b>	<b>4155</b>	<b>4155.0</b>	<b>4155</b>	0.65	0.02	<b>4155.0</b>	<b>4155</b>	0.63	0.00	4155*	4155
E08	59	<b>4710</b>	4715	<b>4710</b>	<b>4710.0</b>	<b>4710</b>	0.87	0.05	<b>4710.0</b>	<b>4710</b>	0.74	0.02	4710*	4710
E09	103	5835	5885	5870	5840.0	<b>5810</b>	2.14	0.60	<b>5810.0</b>	<b>5810</b>	1.93	0.41	5820	<u>5810</u>
E10	49	<b>3605</b>	<b>3605</b>	<b>3605</b>	<b>3605.0</b>	<b>3605</b>	0.64	0.01	<b>3605.0</b>	<b>3605</b>	0.54	0.01	3605*	3605
E11	94	4670	4675	4670	4658.0	<b>4650</b>	2.06	0.65	4655.0	<b>4650</b>	2.31	0.89	4650*	4650
E12	67	4195	4215	4200	4185.5	<b>4180</b>	1.23	0.52	4190.0	<b>4180</b>	1.44	0.53	4180*	4180
E13	52	<b>3345</b>	<b>3345</b>	<b>3345</b>	<b>3345.0</b>	<b>3345</b>	0.68	0.02	<b>3345.0</b>	<b>3345</b>	0.64	0.01	3345*	3345
E14	55	<b>4115</b>	<b>4115</b>	<b>4115</b>	<b>4115.0</b>	<b>4115</b>	0.80	0.02	<b>4115.0</b>	<b>4115</b>	0.70	0.02	4115*	4115
E15	107	4225	4225	4225	4214.0	<b>4205</b>	1.96	0.63	4219.0	<b>4205</b>	2.12	0.57	4205*	4205
E16	54	<b>3775</b>	<b>3775</b>	<b>3775</b>	<b>3775.0</b>	<b>3775</b>	0.93	0.04	<b>3775.0</b>	<b>3775</b>	0.77	0.04	3775*	3775
E17	36	<b>2740</b>	<b>2740</b>	<b>2740</b>	<b>2740.0</b>	<b>2740</b>	0.51	0.02	<b>2740.0</b>	<b>2740</b>	0.43	0.01	2740*	2740
E18	88	<b>3835</b>	<b>3835</b>	<b>3835</b>	<b>3835.0</b>	<b>3835</b>	1.59	0.03	<b>3835.0</b>	<b>3835</b>	1.61	0.07	3835	3835
E19	66	<b>3235</b>	<b>3235</b>	<b>3235</b>	<b>3235.0</b>	<b>3235</b>	1.11	0.03	<b>3235.0</b>	<b>3235</b>	0.98	0.02	3235*	3235
E20	63	<b>2825</b>	<b>2825</b>	<b>2825</b>	<b>2825.0</b>	<b>2825</b>	0.96	0.01	<b>2825.0</b>	<b>2825</b>	0.93	0.00	2825*	2825
E21	72	<b>3730</b>	<b>3730</b>	<b>3730</b>	<b>3730.0</b>	<b>3730</b>	1.57	0.15	<b>3730.0</b>	<b>3730</b>	1.26	0.15	3730*	3730
E22	44	<b>2470</b>	<b>2470</b>	<b>2470</b>	<b>2470.0</b>	<b>2470</b>	0.73	0.01	<b>2470.0</b>	<b>2470</b>	0.70	0.01	2470*	2470
E23	89	<b>3710</b>	3725	<b>3710</b>	<b>3710.0</b>	<b>3710</b>	2.01	0.20	3713.0	<b>3710</b>	1.92	0.67	3710	3710
E24	86	<b>4020</b>	<b>4020</b>	<b>4020</b>	<b>4020.0</b>	<b>4020</b>	1.48	0.05	<b>4020.0</b>	<b>4020</b>	1.48	0.06	4020*	4020
E25	28	<b>1615</b>	<b>1615</b>	<b>1615</b>	<b>1615.0</b>	<b>1615</b>	0.34	0.00	<b>1615.0</b>	<b>1615</b>	0.37	0.00	1615*	1615
F01	85	<b>4040</b>	4060	<b>4040</b>	<b>4040.0</b>	<b>4040</b>	2.05	0.07	<b>4040.0</b>	<b>4040</b>	1.88	0.06	4040*	4040
F02	58	<b>3300</b>	<b>3300</b>	<b>3300</b>	<b>3300.0</b>	<b>3300</b>	1.22	0.01	<b>3300.0</b>	<b>3300</b>	1.20	0.01	3300*	3300
F03	47	<b>1665</b>	<b>1665</b>	<b>1665</b>	<b>1665.0</b>	<b>1665</b>	0.84	0.00	<b>1665.0</b>	<b>1665</b>	0.82	0.00	1665*	1665
F04	77	<b>3485</b>	3505	3495	<b>3485.0</b>	<b>3485</b>	1.77	0.33	<b>3485.0</b>	<b>3485</b>	1.52	0.05	3485*	3485
F05	61	<b>3605</b>	<b>3605</b>	<b>3605</b>	<b>3605.0</b>	<b>3605</b>	1.25	0.02	<b>3605.0</b>	<b>3605</b>	1.18	0.01	3605*	3605
F06	43	<b>1875</b>	<b>1875</b>	<b>1875</b>	<b>1875.0</b>	<b>1875</b>	0.82	0.00	<b>1875.0</b>	<b>1875</b>	0.88	0.00	1875*	1875
F07	50	<b>3335</b>	<b>3335</b>	<b>3335</b>	<b>3335.0</b>	<b>3335</b>	1.00	0.02	<b>3335.0</b>	<b>3335</b>	0.92	0.01	3335*	3335
F08	59	<b>3705</b>	<b>3705</b>	<b>3705</b>	<b>3705.0</b>	<b>3705</b>	0.97	0.01	<b>3705.0</b>	<b>3705</b>	0.91	0.00	3705*	3705
F09	103	<b>4730</b>	4755	<b>4730</b>	<b>4730.0</b>	<b>4730</b>	2.33	0.11	<b>4730.0</b>	<b>4730</b>	2.15	0.06	4730*	4730
F10	49	<b>2925</b>	<b>2925</b>	<b>2925</b>	<b>2925.0</b>	<b>2925</b>	0.84	0.00	<b>2925.0</b>	<b>2925</b>	0.83	0.00	2925*	2925
F11	94	<b>3835</b>	<b>3835</b>	<b>3835</b>	<b>3835.0</b>	<b>3835</b>	2.22	0.05	<b>3835.0</b>	<b>3835</b>	2.11	0.07	3835*	3835
F12	67	<b>3395</b>	<b>3395</b>	<b>3395</b>	<b>3395.0</b>	<b>3395</b>	1.24	0.03	<b>3395.0</b>	<b>3395</b>	1.23	0.03	3395*	3395
F13	52	<b>2855</b>	<b>2855</b>	<b>2855</b>	<b>2855.0</b>	<b>2855</b>	0.93	0.00	<b>2855.0</b>	<b>2855</b>	0.91	0.00	2855*	2855
F14	55	<b>3330</b>	3340	3340	<b>3330.0</b>	<b>3330</b>	1.14	0.02	<b>3330.0</b>	<b>3330</b>	1.03	0.02	3330*	3330
F15	107	<b>3560</b>	3605	<b>3560</b>	<b>3560.0</b>	<b>3560</b>	2.53	0.06	<b>3560.0</b>	<b>3560</b>	2.12	0.05	3560*	3560
F16	54	<b>2725</b>	<b>2725</b>	<b>2725</b>	<b>2725.0</b>	<b>2725</b>	1.14	0.00	<b>2725.0</b>	<b>2725</b>	1.09	0.00	2725*	2725
F17	36	<b>2055</b>	2080	<b>2055</b>	<b>2055.0</b>	<b>2055</b>	0.52	0.00	<b>2055.0</b>	<b>2055</b>	0.56	0.00	2055*	2055
F18	88	<b>3075</b>	<b>3075</b>	<b>3075</b>	<b>3075.0</b>	<b>3075</b>	2.04	0.02	<b>3075.0</b>	<b>3075</b>	1.94	0.03	3075	3075
F19	66	<b>2525</b>	2540	<b>2525</b>	<b>2525.0</b>	<b>2525</b>	1.51	0.01	<b>2525.0</b>	<b>2525</b>	1.39	0.01	2525	2525
F20	63	<b>2445</b>	<b>2445</b>	<b>2445</b>	<b>2445.0</b>	<b>2445</b>	1.42	0.04	<b>2445.0</b>	<b>2445</b>	1.31	0.02	2445*	2445
F21	72	<b>2930</b>	<b>2930</b>	<b>2930</b>	<b>2930.0</b>	<b>2930</b>	1.70	0.01	<b>2930.0</b>	<b>2930</b>	1.62	0.01	2930*	2930
F22	44	<b>2075</b>	<b>2075</b>	<b>2075</b>	<b>2075.0</b>	<b>2075</b>	0.91	0.00	<b>2075.0</b>	<b>2075</b>	0.88	0.00	2075*	2075
F23	89	<b>3005</b>	3010	<b>3005</b>	<b>3005.0</b>	<b>3005</b>	2.11	0.22	<b>3005.0</b>	<b>3005</b>	1.66	0.06	3005	3005
F24	86	<b>3210</b>	3245	3240	<b>3210.0</b>	<b>3210</b>	2.24	0.05	<b>3210.0</b>	<b>3210</b>	2.06	0.07	3210*	3210
F25	28	<b>1390</b>	<b>1390</b>	<b>1390</b>	<b>1390.0</b>	<b>1390</b>	0.43	0.00	<b>1390.0</b>	<b>1390</b>	0.43	0.00	1390*	1390
Gap (%)		0.038%	0.158%	0.075%	0.020%	0.000%			0.013%	0.003%				
T(min)		2.57	—	—			1.26				1.19			
T*(min)		—	1.08	0.35				0.10				0.09		
CPU		P-II 450M	P-M 1.4G	Xe 2.0G			Xe 3.07G				Xe 3.07G			

Table EC.5 Results for the CARP – EGL instances

Inst	$ E_R $	PDHM08	MTY09	UFF13		ILS				UHGS				BKS	
		Avg-10	Single	Avg-15	Best-15	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Old	New
egl-e1-A	51	<b>3548.0</b>	<b>3548</b>	<b>3548.0</b>	<b>3548</b>	<b>3548.0</b>	<b>3548</b>	0.75	0.00	<b>3548.0</b>	<b>3548</b>	0.76	0.00	3548*	3548
egl-e1-B	51	4522.2	<b>4498</b>	4508.6	<b>4498</b>	<b>4498.0</b>	<b>4498</b>	1.07	0.18	<b>4498.0</b>	<b>4498</b>	1.13	0.26	4498*	4498
egl-e1-C	51	5608.0	<b>5595</b>	5615.3	<b>5595</b>	<b>5595.0</b>	<b>5595</b>	0.77	0.03	<b>5595.0</b>	<b>5595</b>	0.78	0.02	5595*	5595
egl-e2-A	72	5023.8	<b>5018</b>	<b>5018.0</b>	<b>5018</b>	<b>5018.0</b>	<b>5018</b>	1.13	0.02	<b>5018.0</b>	<b>5018</b>	1.26	0.02	5018*	5018
egl-e2-B	72	6335.4	<b>6317</b>	6330.7	<b>6317</b>	6317.2	<b>6317</b>	1.65	0.55	6321.2	<b>6317</b>	2.25	0.98	6317	6317
egl-e2-C	72	8355.9	<b>8335</b>	8335.8	<b>8335</b>	<b>8335.0</b>	<b>8335</b>	1.20	0.09	<b>8335.0</b>	<b>8335</b>	1.34	0.18	8335*	8335
egl-e3-A	87	<b>5898.0</b>	<b>5898</b>	<b>5898.0</b>	<b>5898</b>	<b>5898.0</b>	<b>5898</b>	1.79	0.03	<b>5898.0</b>	<b>5898</b>	1.99	0.10	5898*	5898
egl-e3-B	87	7806.4	7787	7787.3	7777	7775.6	<b>7775</b>	1.96	0.61	7776.4	<b>7775</b>	1.83	0.44	7775	7775
egl-e3-C	87	10322.3	10305	10296.5	<b>10292</b>	<b>10292.0</b>	<b>10292</b>	1.84	0.78	<b>10292.0</b>	<b>10292</b>	2.25	0.77	10292	10292
egl-e4-A	98	6459.4	6461	6461.1	<b>6444</b>	6458.0	6446	2.22	0.17	<b>6444.0</b>	<b>6444</b>	3.68	1.80	6444	6444
egl-e4-B	98	9016.3	9026	9037.1	9002	8996.8	8987	2.25	1.11	8985.3	<b>8961</b>	3.10	1.47	8961	8961
egl-e4-C	98	11750.1	11598	11670.0	11626	11563.7	<b>11529</b>	2.66	1.33	11562.8	<b>11529</b>	5.04	3.35	11561	<u>11529</u>
egl-s1-A	75	<b>5018.0</b>	<b>5018</b>	5038.9	<b>5018</b>	<b>5018.0</b>	<b>5018</b>	1.30	0.02	<b>5018.0</b>	<b>5018</b>	1.38	0.02	5018*	5018
egl-s1-B	75	<b>6388.0</b>	6394	6388.4	<b>6388</b>	<b>6388.0</b>	<b>6388</b>	1.49	0.27	<b>6388.0</b>	<b>6388</b>	1.38	0.14	6388*	6388
egl-s1-C	75	8518.2	<b>8518</b>	8521.5	<b>8518</b>	<b>8518.0</b>	<b>8518</b>	1.14	0.06	<b>8518.0</b>	<b>8518</b>	1.20	0.08	8518*	8518
egl-s2-A	147	9997.9	9970	9980.5	9903	9893.2	<b>9875</b>	4.86	2.16	9886.5	<b>9875</b>	8.70	5.65	9884	<u>9875</u>
egl-s2-B	147	13176.0	13345	13240.6	13169	13125.6	13095	4.92	2.25	13101.9	<b>13081</b>	7.59	4.54	13100	<u>13057</u>
egl-s2-C	147	16551.6	16600	16539.9	16442	16451.2	<b>16425</b>	4.59	3.00	16440.2	<b>16425</b>	7.31	4.57	16425*	16425
egl-s3-A	159	10291.2	10284	10276.1	<b>10221</b>	10243.2	<b>10221</b>	5.21	2.32	10240.0	<b>10221</b>	7.92	4.39	10220	<u>10201</u>
egl-s3-B	159	13829.2	13857	13860.7	13694	13714.4	<b>13682</b>	5.25	1.76	13693.5	<b>13682</b>	9.98	6.63	13682	13682
egl-s3-C	159	17327.9	17316	17277.7	17221	17243.2	17196	5.65	2.38	17191.3	<b>17188</b>	8.93	5.85	17188*	17188
egl-s4-A	190	12440.4	12348	12406.5	12297	12313.3	<b>12257</b>	6.52	3.10	12287.9	12273	14.72	10.46	12268	<u>12216</u>
egl-s4-B	190	16410.3	16442	16432.0	16333	16303.3	16265	7.71	4.75	16283.9	<b>16230</b>	18.18	14.21	16321	<u>16214</u>
egl-s4-C	190	20731.5	20821	20660.5	20563	20637.4	20577	14.76	8.20	20591.4	<b>20500</b>	21.51	17.24	20481	<u>20461</u>
Gap (%)		0.625%	0.555%	0.562%	0.207%	0.209%	0.088%			0.141%	0.049%				
T(min)		30.00	—	13.31	—			3.45				5.59			
T*(min)		8.39	2.10	—	—				1.47				3.46		
CPU		P-IV 3.6G	Xe 2G	I4 3.0G			Xe 3.07G				Xe 3.07G				

Table EC.6 Results for the CARP – EGL-L instances

Inst	$ E_R $	BE08		MPS13		MLY14		ILS			UHGS			BKS		
		Single	Best-10	Avg-10	Best-10	Avg-30	Best-30	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Old
egl-g1-A	347	1049708	1010937.4	1004864	1007619	998777	993374.8	<b>992045</b>	19.84	11.99	993127.4	992227	34.80	24.61	998777	991176
egl-g1-B	347	1140692	1137141.5	1129937	1122863	1118030	1117402.7	1114565	19.97	7.77	1116617.0	<b>1112149</b>	34.76	23.26	1118030	<u>1109656</u>
egl-g1-C	347	1282270	1266576.8	1262888	1250174	1243403	1241243.4	1238534	27.18	16.92	1236062.0	<b>1232501</b>	44.39	36.00	1243403	<u>1230155</u>
egl-g1-D	347	1420126	1406929.0	1398958	1386120	1373389	1374848.2	1372867	28.11	16.88	1370963.0	<b>1365393</b>	49.15	40.78	1373389	<u>1361862</u>
egl-g1-E	347	1583133	1554220.2	1543804	1525629	1517424	1514409.0	1507131	35.71	23.38	1511572.0	<b>1503467</b>	49.52	41.30	1517424	<u>1501801</u>
egl-g2-A	375	1129229	1118363.0	1115339	1104944	1097578	1093665.3	1089698	28.08	16.21	1090396.0	<b>1087353</b>	49.00	39.17	1097578	<u>1086932</u>
egl-g2-B	375	1255907	1233720.5	1226645	1221429	1209694	1206808.9	1201860	25.73	14.50	1202901.0	<b>1198633</b>	48.76	38.12	1209694	<u>1196873</u>
egl-g2-C	375	1418145	1374479.7	1371004	1355548	1342637	1342004.7	1338873	28.27	15.08	1336104.0	<b>1333430</b>	49.43	38.49	1342637	<u>1330744</u>
egl-g2-D	375	1516103	1515119.3	1509990	1492063	1483558	1479210.4	1473039	32.11	16.02	1476285.0	<b>1471783</b>	50.36	40.53	1483558	<u>1468310</u>
egl-g2-E	375	1701681	1658378.1	1659217	1629002	1620692	1618758.7	1613410	38.60	19.54	1616556.0	<b>1610919</b>	51.06	42.46	1620692	<u>1602229</u>
Gap (%)		4.749%	3.018%	2.591%	1.671%	0.962%	0.768%	0.468%			0.529%	0.206%				
T(min)		—	20.66		33.43		28.36				46.12					
T*(min)		17.02	—		—		15.83									
CPU		P-M 1.4G	I5 3.2G		I7 3.4G		Xe 3.07G				Xe 3.07G					



**Table EC.7** Results for the MCGRP – MGGDB instances

Inst	$ N_R $	$ E_R $	$ A_R $	BLMV14	DHDI14	ILS				UHGS				BKS	
				Single	Single	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Old	New
mggdb-0.25-1	6	3	12	<b>280</b>	<b>280</b>	<b>280.0</b>	<b>280</b>	0.13	0.00	<b>280.0</b>	<b>280</b>	0.19	0.00	280*	280
mggdb-0.25-2	6	4	15	359	<b>349</b>	<b>349.0</b>	<b>349</b>	0.24	0.02	<b>349.0</b>	<b>349</b>	0.24	0.01	349*	349
mggdb-0.25-3	7	3	12	286	<b>278</b>	<b>278.0</b>	<b>278</b>	0.14	0.00	<b>278.0</b>	<b>278</b>	0.20	0.00	278*	278
mggdb-0.25-4	4	3	11	<b>289</b>	<b>289</b>	<b>289.0</b>	<b>289</b>	0.11	0.00	<b>289.0</b>	<b>289</b>	0.15	0.00	289*	289
mggdb-0.25-5	5	4	15	410	<b>394</b>	<b>394.0</b>	<b>394</b>	0.20	0.00	<b>394.0</b>	<b>394</b>	0.21	0.00	394*	394
mggdb-0.25-6	6	3	12	295	<b>292</b>	<b>292.0</b>	<b>292</b>	0.15	0.00	<b>292.0</b>	<b>292</b>	0.19	0.00	292*	292
mggdb-0.25-7	5	3	12	302	<b>290</b>	<b>290.0</b>	<b>290</b>	0.12	0.00	<b>290.0</b>	<b>290</b>	0.17	0.00	290*	290
mggdb-0.25-8	11	8	26	351	<b>336</b>	<b>336.0</b>	<b>336</b>	0.56	0.06	<b>336.0</b>	<b>336</b>	0.57	0.02	336	336
mggdb-0.25-9	9	9	29	316	<b>309</b>	<b>309.0</b>	<b>309</b>	0.62	0.06	<b>309.0</b>	<b>309</b>	0.63	0.04	309	309
mggdb-0.25-10	4	4	14	<b>265</b>	<b>265</b>	<b>265.0</b>	<b>265</b>	0.18	0.00	<b>265.0</b>	<b>265</b>	0.19	0.00	265*	265
mggdb-0.25-11	8	8	25	369	<b>356</b>	<b>356.0</b>	<b>356</b>	0.55	0.00	<b>356.0</b>	<b>356</b>	0.70	0.00	356*	356
mggdb-0.25-12	6	3	13	465	<b>459</b>	<b>459.0</b>	<b>459</b>	0.14	0.00	<b>459.0</b>	<b>459</b>	0.20	0.00	459*	459
mggdb-0.25-13	6	5	15	392	<b>388</b>	388.8	<b>388</b>	0.25	0.08	<b>388.0</b>	<b>388</b>	0.27	0.02	388*	388
mggdb-0.25-14	5	3	12	<b>107</b>	<b>107</b>	<b>107.0</b>	<b>107</b>	0.12	0.00	<b>107.0</b>	<b>107</b>	0.15	0.00	107*	107
mggdb-0.25-15	5	3	12	<b>55</b>	<b>55</b>	<b>55.0</b>	<b>55</b>	0.09	0.00	<b>55.0</b>	<b>55</b>	0.12	0.00	55*	55
mggdb-0.25-16	5	5	15	<b>98</b>	<b>98</b>	<b>98.0</b>	<b>98</b>	0.22	0.00	<b>98.0</b>	<b>98</b>	0.23	0.00	98*	98
mggdb-0.25-17	5	5	15	<b>71</b>	<b>71</b>	<b>71.0</b>	<b>71</b>	0.13	0.00	<b>71.0</b>	<b>71</b>	0.15	0.00	71*	71
mggdb-0.25-18	6	6	20	<b>144</b>	<b>144</b>	<b>144.0</b>	<b>144</b>	0.33	0.00	<b>144.0</b>	<b>144</b>	0.42	0.00	144*	144
mggdb-0.25-19	3	1	6	<b>53</b>	<b>53</b>	<b>53.0</b>	<b>53</b>	0.04	0.00	<b>53.0</b>	<b>53</b>	0.08	0.00	53*	53
mggdb-0.25-20	5	3	12	117	<b>116</b>	<b>116.0</b>	<b>116</b>	0.14	0.00	<b>116.0</b>	<b>116</b>	0.20	0.00	116*	116
mggdb-0.25-21	7	6	18	<b>146</b>	<b>146</b>	<b>146.0</b>	<b>146</b>	0.33	0.00	<b>146.0</b>	<b>146</b>	0.41	0.00	146*	146
mggdb-0.25-22	6	8	24	168	<b>160</b>	<b>160.0</b>	<b>160</b>	0.54	0.02	<b>160.0</b>	<b>160</b>	0.61	0.01	160*	160
mggdb-0.25-23	8	9	31	186	<b>181</b>	<b>181.0</b>	<b>181</b>	0.72	0.04	<b>181.0</b>	<b>181</b>	0.77	0.05	181*	181
mggdb-0.30-1	7	3	11	276	<b>273</b>	<b>273.0</b>	<b>273</b>	0.13	0.00	<b>273.0</b>	<b>273</b>	0.18	0.00	273*	273
mggdb-0.30-2	6	4	14	314	<b>301</b>	<b>301.0</b>	<b>301</b>	0.15	0.00	<b>301.0</b>	<b>301</b>	0.19	0.00	301*	301
mggdb-0.30-3	5	3	11	278	<b>270</b>	<b>270.0</b>	<b>270</b>	0.12	0.00	<b>270.0</b>	<b>270</b>	0.16	0.00	270*	270
mggdb-0.30-4	6	2	10	<b>260</b>	<b>260</b>	<b>260.0</b>	<b>260</b>	0.10	0.00	<b>260.0</b>	<b>260</b>	0.15	0.00	260*	260
mggdb-0.30-5	7	4	14	399	<b>388</b>	<b>388.0</b>	<b>388</b>	0.18	0.00	<b>388.0</b>	<b>388</b>	0.23	0.00	388*	388
mggdb-0.30-6	8	3	11	<b>276</b>	<b>276</b>	<b>276.0</b>	<b>276</b>	0.13	0.00	<b>276.0</b>	<b>276</b>	0.19	0.00	276*	276
mggdb-0.30-7	6	3	11	277	<b>273</b>	<b>273.0</b>	<b>273</b>	0.10	0.00	<b>273.0</b>	<b>273</b>	0.16	0.00	273*	273
mggdb-0.30-8	15	7	24	338	<b>331</b>	<b>331.0</b>	<b>331</b>	0.50	0.02	<b>331.0</b>	<b>331</b>	0.60	0.02	331	331
mggdb-0.30-9	11	8	27	284	<b>281</b>	<b>281.0</b>	<b>281</b>	0.63	0.06	<b>281.0</b>	<b>281</b>	0.62	0.06	281*	281
mggdb-0.30-10	5	4	13	<b>242</b>	<b>242</b>	<b>242.0</b>	<b>242</b>	0.16	0.00	<b>242.0</b>	<b>242</b>	0.21	0.00	242*	242
mggdb-0.30-11	13	7	23	399	<b>387</b>	<b>387.0</b>	<b>387</b>	0.62	0.01	<b>387.0</b>	<b>387</b>	0.73	0.01	387*	387
mggdb-0.30-12	6	3	12	472	<b>467</b>	<b>467.0</b>	<b>467</b>	0.13	0.00	<b>467.0</b>	<b>467</b>	0.18	0.00	467*	467
mggdb-0.30-13	6	4	14	<b>483</b>	<b>483</b>	<b>483.0</b>	<b>483</b>	0.20	0.00	<b>483.0</b>	<b>483</b>	0.24	0.00	483*	483
mggdb-0.30-14	3	3	11	<b>101</b>	<b>101</b>	<b>101.0</b>	<b>101</b>	0.10	0.00	<b>101.0</b>	<b>101</b>	0.15	0.00	101*	101
mggdb-0.30-15	5	3	11	<b>44</b>	<b>44</b>	<b>44.0</b>	<b>44</b>	0.09	0.00	<b>44.0</b>	<b>44</b>	0.11	0.00	44*	44
mggdb-0.30-16	6	4	14	107	<b>105</b>	<b>105.0</b>	<b>105</b>	0.24	0.01	<b>105.0</b>	<b>105</b>	0.24	0.01	105*	105
mggdb-0.30-17	4	4	14	67	<b>65</b>	<b>65.0</b>	<b>65</b>	0.11	0.00	<b>65.0</b>	<b>65</b>	0.14	0.00	65*	65
mggdb-0.30-18	6	6	18	<b>144</b>	<b>144</b>	<b>144.0</b>	<b>144</b>	0.25	0.00	<b>144.0</b>	<b>144</b>	0.27	0.00	144*	144
mggdb-0.30-19	3	1	6	<b>51</b>	<b>51</b>	<b>51.0</b>	<b>51</b>	0.04	0.00	<b>51.0</b>	<b>51</b>	0.07	0.00	51*	51
mggdb-0.30-20	4	3	11	97	<b>94</b>	<b>94.0</b>	<b>94</b>	0.11	0.00	<b>94.0</b>	<b>94</b>	0.17	0.00	94*	94
mggdb-0.30-21	6	5	17	122	<b>121</b>	<b>121.0</b>	<b>121</b>	0.26	0.00	<b>121.0</b>	<b>121</b>	0.32	0.00	121*	121
mggdb-0.30-22	7	7	23	156	<b>153</b>	<b>153.0</b>	<b>153</b>	0.44	0.00	<b>153.0</b>	<b>153</b>	0.53	0.00	153*	153
mggdb-0.30-23	9	9	29	171	<b>167</b>	<b>167.0</b>	<b>167</b>	0.73	0.05	<b>167.0</b>	<b>167</b>	0.88	0.25	167*	167

**Table EC.8** Results for the MCGRP – MGGDB instances (continued)

Inst	$ N_R $	$ E_R $	$ A_R $	BLMV14	DHDI14	ILS				UHGS				BKS	
				Single	Single	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Old	New
mggdb-0.35-1	7	3	11	<b>252</b>	<b>252</b>	<b>252.0</b>	<b>252</b>	0.12	0.00	<b>252.0</b>	<b>252</b>	0.17	0.00	252*	252
mggdb-0.35-2	6	3	13	<b>284</b>	<b>284</b>	<b>284.0</b>	<b>284</b>	0.13	0.00	<b>284.0</b>	<b>284</b>	0.18	0.00	284*	284
mggdb-0.35-3	6	3	11	<b>243</b>	<b>243</b>	<b>243.0</b>	<b>243</b>	0.11	0.00	<b>243.0</b>	<b>243</b>	0.17	0.00	243*	243
mggdb-0.35-4	6	2	9	<b>242</b>	<b>242</b>	<b>242.0</b>	<b>242</b>	0.10	0.00	<b>242.0</b>	<b>242</b>	0.14	0.00	242*	242
mggdb-0.35-5	7	3	13	317	<b>309</b>	<b>309.0</b>	<b>309</b>	0.15	0.00	<b>309.0</b>	<b>309</b>	0.21	0.00	309*	309
mggdb-0.35-6	7	3	11	<b>262</b>	<b>262</b>	<b>262.0</b>	<b>262</b>	0.12	0.00	<b>262.0</b>	<b>262</b>	0.18	0.00	262*	262
mggdb-0.35-7	8	3	11	<b>272</b>	<b>272</b>	<b>272.0</b>	<b>272</b>	0.12	0.00	<b>272.0</b>	<b>272</b>	0.18	0.00	272*	272
mggdb-0.35-8	9	7	22	321	<b>316</b>	<b>316.0</b>	<b>316</b>	0.51	0.16	<b>316.0</b>	<b>316</b>	0.49	0.04	316	316
mggdb-0.35-9	13	7	25	274	<b>266</b>	<b>266.0</b>	<b>266</b>	0.55	0.02	<b>266.0</b>	<b>266</b>	0.58	0.03	266	266
mggdb-0.35-10	9	3	12	<b>268</b>	<b>268</b>	<b>268.0</b>	<b>268</b>	0.20	0.00	<b>268.0</b>	<b>268</b>	0.27	0.00	268*	268
mggdb-0.35-11	12	7	22	313	<b>303</b>	<b>303.0</b>	<b>303</b>	0.57	0.01	<b>303.0</b>	<b>303</b>	0.67	0.01	303*	303
mggdb-0.35-12	6	3	11	<b>461</b>	<b>461</b>	<b>461.0</b>	<b>461</b>	0.11	0.00	<b>461.0</b>	<b>461</b>	0.15	0.00	461*	461
mggdb-0.35-13	7	4	13	435	<b>417</b>	<b>417.0</b>	<b>417</b>	0.19	0.00	<b>417.0</b>	<b>417</b>	0.23	0.01	417*	417
mggdb-0.35-14	5	3	10	85	<b>84</b>	<b>84.0</b>	<b>84</b>	0.10	0.00	<b>84.0</b>	<b>84</b>	0.16	0.00	84*	84
mggdb-0.35-15	5	3	10	<b>44</b>	<b>44</b>	<b>44.0</b>	<b>44</b>	0.08	0.00	<b>44.0</b>	<b>44</b>	0.11	0.00	44*	44
mggdb-0.35-16	5	4	13	<b>75</b>	<b>75</b>	<b>75.0</b>	<b>75</b>	0.14	0.00	<b>75.0</b>	<b>75</b>	0.23	0.00	75*	75
mggdb-0.35-17	6	4	13	<b>62</b>	<b>62</b>	<b>62.0</b>	<b>62</b>	0.12	0.00	<b>62.0</b>	<b>62</b>	0.15	0.00	62*	62
mggdb-0.35-18	8	5	17	137	<b>135</b>	<b>135.0</b>	<b>135</b>	0.32	0.00	<b>135.0</b>	<b>135</b>	0.34	0.00	135*	135
mggdb-0.35-19	3	1	5	<b>51</b>	<b>51</b>	<b>51.0</b>	<b>51</b>	0.03	0.00	<b>51.0</b>	<b>51</b>	0.06	0.00	51*	51
mggdb-0.35-20	6	3	11	<b>96</b>	<b>96</b>	<b>96.0</b>	<b>96</b>	0.14	0.00	<b>96.0</b>	<b>96</b>	0.18	0.00	96*	96
mggdb-0.35-21	7	5	16	122	<b>120</b>	<b>120.0</b>	<b>120</b>	0.24	0.00	<b>120.0</b>	<b>120</b>	0.32	0.00	120*	120
mggdb-0.35-22	8	7	21	143	<b>139</b>	<b>139.0</b>	<b>139</b>	0.38	0.00	<b>139.0</b>	<b>139</b>	0.49	0.00	139*	139
mggdb-0.35-23	9	8	27	185	<b>179</b>	<b>179.0</b>	<b>179</b>	0.59	0.02	<b>179.0</b>	<b>179</b>	0.66	0.05	179*	179
mggdb-0.40-1	6	3	10	<b>279</b>	<b>279</b>	<b>279.0</b>	<b>279</b>	0.10	0.00	<b>279.0</b>	<b>279</b>	0.16	0.00	279*	279
mggdb-0.40-2	7	3	12	320	<b>308</b>	<b>308.0</b>	<b>308</b>	0.16	0.00	<b>308.0</b>	<b>308</b>	0.20	0.00	308*	308
mggdb-0.40-3	7	3	10	229	<b>225</b>	<b>225.0</b>	<b>225</b>	0.12	0.00	<b>225.0</b>	<b>225</b>	0.17	0.00	225*	225
mggdb-0.40-4	6	2	9	<b>238</b>	<b>238</b>	<b>238.0</b>	<b>238</b>	0.10	0.00	<b>238.0</b>	<b>238</b>	0.15	0.00	238*	238
mggdb-0.40-5	7	3	12	346	<b>344</b>	<b>344.0</b>	<b>344</b>	0.15	0.00	<b>344.0</b>	<b>344</b>	0.19	0.00	344*	344
mggdb-0.40-6	6	3	10	281	<b>270</b>	<b>270.0</b>	<b>270</b>	0.11	0.00	<b>270.0</b>	<b>270</b>	0.16	0.00	270*	270
mggdb-0.40-7	6	3	10	283	<b>282</b>	<b>282.0</b>	<b>282</b>	0.11	0.00	<b>282.0</b>	<b>282</b>	0.16	0.00	282*	282
mggdb-0.40-8	13	6	21	340	<b>331</b>	<b>331.0</b>	<b>331</b>	0.42	0.01	<b>331.0</b>	<b>331</b>	0.49	0.01	331	331
mggdb-0.40-9	15	7	23	285	<b>275</b>	<b>275.0</b>	<b>275</b>	0.50	0.01	<b>275.0</b>	<b>275</b>	0.55	0.01	275	275
mggdb-0.40-10	8	3	11	<b>191</b>	<b>191</b>	<b>191.0</b>	<b>191</b>	0.15	0.00	<b>191.0</b>	<b>191</b>	0.18	0.00	191*	191
mggdb-0.40-11	12	6	20	287	<b>283</b>	<b>283.0</b>	<b>283</b>	0.46	0.00	<b>283.0</b>	<b>283</b>	0.58	0.01	283	283
mggdb-0.40-12	6	3	10	<b>412</b>	<b>412</b>	<b>412.0</b>	<b>412</b>	0.11	0.00	<b>412.0</b>	<b>412</b>	0.16	0.00	412*	412
mggdb-0.40-13	7	4	12	<b>405</b>	406	<b>405.0</b>	<b>405</b>	0.17	0.01	<b>405.0</b>	<b>405</b>	0.22	0.00	405*	405
mggdb-0.40-14	6	3	9	<b>62</b>	<b>62</b>	<b>62.0</b>	<b>62</b>	0.10	0.00	<b>62.0</b>	<b>62</b>	0.15	0.00	62*	62
mggdb-0.40-15	6	3	9	<b>37</b>	<b>37</b>	<b>37.0</b>	<b>37</b>	0.08	0.00	<b>37.0</b>	<b>37</b>	0.11	0.00	37*	37
mggdb-0.40-16	5	4	12	<b>84</b>	<b>84</b>	<b>84.0</b>	<b>84</b>	0.16	0.00	<b>84.0</b>	<b>84</b>	0.22	0.00	84*	84
mggdb-0.40-17	5	4	12	<b>65</b>	<b>65</b>	<b>65.0</b>	<b>65</b>	0.12	0.00	<b>65.0</b>	<b>65</b>	0.15	0.00	65*	65
mggdb-0.40-18	6	5	16	122	<b>119</b>	<b>119.0</b>	<b>119</b>	0.24	0.00	<b>119.0</b>	<b>119</b>	0.31	0.00	119*	119
mggdb-0.40-19	4	1	5	<b>38</b>	<b>38</b>	<b>38.0</b>	<b>38</b>	0.04	0.00	<b>38.0</b>	<b>38</b>	0.05	0.00	38*	38
mggdb-0.40-20	6	3	10	<b>94</b>	<b>94</b>	<b>94.0</b>	<b>94</b>	0.13	0.00	<b>94.0</b>	<b>94</b>	0.20	0.00	94*	94
mggdb-0.40-21	9	4	15	106	<b>104</b>	<b>104.0</b>	<b>104</b>	0.23	0.00	<b>104.0</b>	<b>104</b>	0.29	0.01	104*	104
mggdb-0.40-22	8	6	19	132	<b>129</b>	<b>129.0</b>	<b>129</b>	0.33	0.01	<b>129.0</b>	<b>129</b>	0.40	0.00	129*	129
mggdb-0.40-23	10	7	25	165	<b>160</b>	160.3	<b>160</b>	0.63	0.17	160.4	<b>160</b>	0.80	0.24	160*	160

**Table EC.9** Results for the MCGRP – MGGDB instances (end)

Inst	$ N_R $	$ E_R $	$ A_R $	BLMV14	DHDI14	ILS				UHGS				BKS	
				Single	Single	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Old	New
mggdb-0.45-1	6	2	9	<b>259</b>	<b>259</b>	<b>259.0</b>	<b>259</b>	0.09	0.00	<b>259.0</b>	<b>259</b>	0.14	0.00	259*	259
mggdb-0.45-2	7	3	11	302	<b>298</b>	<b>298.0</b>	<b>298</b>	0.13	0.00	<b>298.0</b>	<b>298</b>	0.18	0.00	298*	298
mggdb-0.45-3	8	2	9	245	<b>237</b>	<b>237.0</b>	<b>237</b>	0.11	0.00	<b>237.0</b>	<b>237</b>	0.16	0.00	237*	237
mggdb-0.45-4	7	2	8	<b>228</b>	<b>228</b>	<b>228.0</b>	<b>228</b>	0.09	0.00	<b>228.0</b>	<b>228</b>	0.14	0.00	228*	228
mggdb-0.45-5	7	3	11	357	<b>350</b>	<b>350.0</b>	<b>350</b>	0.14	0.00	<b>350.0</b>	<b>350</b>	0.19	0.00	350*	350
mggdb-0.45-6	7	2	9	225	<b>218</b>	<b>218.0</b>	<b>218</b>	0.10	0.00	<b>218.0</b>	<b>218</b>	0.15	0.00	218*	218
mggdb-0.45-7	9	2	9	<b>243</b>	<b>243</b>	<b>243.0</b>	<b>243</b>	0.09	0.00	<b>243.0</b>	<b>243</b>	0.15	0.00	243*	243
mggdb-0.45-8	16	6	19	312	<b>296</b>	<b>296.0</b>	<b>296</b>	0.44	0.01	<b>296.0</b>	<b>296</b>	0.47	0.01	296*	296
mggdb-0.45-9	14	6	21	287	<b>277</b>	<b>277.0</b>	<b>277</b>	0.51	0.04	<b>277.0</b>	<b>277</b>	0.54	0.05	277*	277
mggdb-0.45-10	9	3	10	<b>214</b>	<b>214</b>	<b>214.0</b>	<b>214</b>	0.17	0.00	<b>214.0</b>	<b>214</b>	0.22	0.00	214*	214
mggdb-0.45-11	15	6	18	310	<b>297</b>	<b>297.0</b>	<b>297</b>	0.46	0.00	<b>297.0</b>	<b>297</b>	0.57	0.01	297	297
mggdb-0.45-12	10	2	9	406	<b>393</b>	<b>393.0</b>	<b>393</b>	0.13	0.00	<b>393.0</b>	<b>393</b>	0.19	0.00	393*	393
mggdb-0.45-13	7	3	11	<b>423</b>	<b>423</b>	<b>423.0</b>	<b>423</b>	0.22	0.01	<b>423.0</b>	<b>423</b>	0.20	0.00	423*	423
mggdb-0.45-14	6	2	8	67	<b>66</b>	<b>66.0</b>	<b>66</b>	0.08	0.00	<b>66.0</b>	<b>66</b>	0.13	0.00	66*	66
mggdb-0.45-15	6	2	8	36	<b>34</b>	<b>34.0</b>	<b>34</b>	0.07	0.00	<b>34.0</b>	<b>34</b>	0.09	0.00	34*	34
mggdb-0.45-16	6	3	11	<b>70</b>	<b>70</b>	<b>70.0</b>	<b>70</b>	0.12	0.00	<b>70.0</b>	<b>70</b>	0.17	0.00	70*	70
mggdb-0.45-17	7	3	11	<b>53</b>	<b>53</b>	<b>53.0</b>	<b>53</b>	0.09	0.00	<b>53.0</b>	<b>53</b>	0.12	0.00	53*	53
mggdb-0.45-18	7	4	14	<b>123</b>	<b>123</b>	<b>123.0</b>	<b>123</b>	0.20	0.00	<b>123.0</b>	<b>123</b>	0.26	0.00	123	123
mggdb-0.45-19	3	1	4	<b>48</b>	<b>48</b>	<b>48.0</b>	<b>48</b>	0.03	0.00	<b>48.0</b>	<b>48</b>	0.06	0.00	48*	48
mggdb-0.45-20	5	2	9	<b>78</b>	<b>78</b>	<b>78.0</b>	<b>78</b>	0.11	0.00	<b>78.0</b>	<b>78</b>	0.15	0.00	78*	78
mggdb-0.45-21	7	4	13	128	<b>122</b>	<b>122.0</b>	<b>122</b>	0.20	0.00	<b>122.0</b>	<b>122</b>	0.25	0.00	122*	122
mggdb-0.45-22	9	6	18	139	<b>136</b>	<b>136.0</b>	<b>136</b>	0.32	0.00	<b>136.0</b>	<b>136</b>	0.41	0.00	136*	136
mggdb-0.45-23	9	7	23	147	145	144.6	<b>144</b>	0.52	0.11	144.8	<b>144</b>	0.69	0.19	144*	144
mggdb-0.50-1	8	2	8	<b>214</b>	<b>214</b>	<b>214.0</b>	<b>214</b>	0.09	0.00	<b>214.0</b>	<b>214</b>	0.14	0.00	214*	214
mggdb-0.50-2	6	3	10	281	<b>269</b>	<b>269.0</b>	<b>269</b>	0.11	0.00	<b>269.0</b>	<b>269</b>	0.15	0.00	269*	269
mggdb-0.50-3	9	2	8	<b>218</b>	<b>218</b>	<b>218.0</b>	<b>218</b>	0.10	0.00	<b>218.0</b>	<b>218</b>	0.15	0.00	218*	218
mggdb-0.50-4	6	2	7	<b>219</b>	<b>219</b>	<b>219.0</b>	<b>219</b>	0.07	0.00	<b>219.0</b>	<b>219</b>	0.11	0.00	219*	219
mggdb-0.50-5	7	3	10	<b>292</b>	<b>292</b>	<b>292.0</b>	<b>292</b>	0.11	0.00	<b>292.0</b>	<b>292</b>	0.16	0.00	292*	292
mggdb-0.50-6	7	2	8	<b>276</b>	<b>276</b>	<b>276.0</b>	<b>276</b>	0.09	0.00	<b>276.0</b>	<b>276</b>	0.14	0.00	276*	276
mggdb-0.50-7	9	2	8	274	<b>265</b>	<b>265.0</b>	<b>265</b>	0.11	0.00	<b>265.0</b>	<b>265</b>	0.16	0.00	265*	265
mggdb-0.50-8	15	5	17	<b>310</b>	<b>310</b>	<b>310.0</b>	<b>310</b>	0.35	0.01	<b>310.0</b>	<b>310</b>	0.39	0.01	310	310
mggdb-0.50-9	16	6	19	270	<b>265</b>	<b>265.0</b>	<b>265</b>	0.42	0.00	<b>265.0</b>	<b>265</b>	0.48	0.01	265	265
mggdb-0.50-10	7	3	9	<b>194</b>	<b>194</b>	<b>194.0</b>	<b>194</b>	0.12	0.00	<b>194.0</b>	<b>194</b>	0.17	0.00	194*	194
mggdb-0.50-11	16	5	17	278	<b>275</b>	<b>275.0</b>	<b>275</b>	0.53	0.02	<b>275.0</b>	<b>275</b>	0.53	0.02	275	275
mggdb-0.50-12	8	2	9	<b>445</b>	<b>445</b>	<b>445.0</b>	<b>445</b>	0.11	0.00	<b>445.0</b>	<b>445</b>	0.16	0.00	445*	445
mggdb-0.50-13	8	3	10	<b>259</b>	261	<b>259.0</b>	<b>259</b>	0.16	0.00	<b>259.0</b>	<b>259</b>	0.20	0.00	259*	259
mggdb-0.50-14	6	2	8	76	<b>75</b>	<b>75.0</b>	<b>75</b>	0.09	0.00	<b>75.0</b>	<b>75</b>	0.14	0.00	75*	75
mggdb-0.50-15	5	2	8	<b>37</b>	<b>37</b>	<b>37.0</b>	<b>37</b>	0.06	0.00	<b>37.0</b>	<b>37</b>	0.09	0.00	37*	37
mggdb-0.50-16	6	3	10	<b>66</b>	<b>66</b>	<b>66.0</b>	<b>66</b>	0.11	0.00	<b>66.0</b>	<b>66</b>	0.16	0.00	66*	66
mggdb-0.50-17	7	3	10	<b>53</b>	<b>53</b>	<b>53.0</b>	<b>53</b>	0.11	0.00	<b>53.0</b>	<b>53</b>	0.13	0.00	53*	53
mggdb-0.50-18	8	4	13	122	<b>121</b>	<b>121.0</b>	<b>121</b>	0.21	0.00	<b>121.0</b>	<b>121</b>	0.26	0.00	121	121
mggdb-0.50-19	3	1	4	<b>44</b>	<b>44</b>	<b>44.0</b>	<b>44</b>	0.03	0.00	<b>44.0</b>	<b>44</b>	0.05	0.00	44*	44
mggdb-0.50-20	5	2	8	<b>81</b>	<b>81</b>	<b>81.0</b>	<b>81</b>	0.08	0.00	<b>81.0</b>	<b>81</b>	0.11	0.00	81*	81
mggdb-0.50-21	8	4	12	88	<b>86</b>	<b>86.0</b>	<b>86</b>	0.17	0.00	<b>86.0</b>	<b>86</b>	0.21	0.00	86*	86
mggdb-0.50-22	10	5	16	127	<b>123</b>	<b>123.0</b>	<b>123</b>	0.33	0.02	<b>123.0</b>	<b>123</b>	0.34	0.01	123*	123
mggdb-0.50-23	7	6	21	<b>125</b>	126	<b>125.0</b>	<b>125</b>	0.45	0.06	<b>125.0</b>	<b>125</b>	0.52	0.12	125*	125
Gap (%)				1.342%	0.018%	0.006%	0.000%			0.006%	0.000%				
T(min)				0.31	60.00			0.22				0.27			
T*(min)				—	0.86			0.01				0.01			
CPU				Xe 3.0G	CPU 3G	Xe 3.07G				Xe 3.07G					

Table EC.10 Results for the MCGRP – MGVAL instances

Inst	$ N_R $	$ E_R $	$ A_R $	BLMV14	DHDI14	ILS				UHGS				BKS	
				Single	Single	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Old	New
mgval-0.25-1A	13	15	26	<b>177</b>	<b>177</b>	<b>177.0</b>	<b>177</b>	0.57	0.01	<b>177.0</b>	<b>177</b>	0.57	0.01	177*	177
mgval-0.25-1B	10	9	28	<b>217</b>	<b>217</b>	<b>217.0</b>	<b>217</b>	0.85	0.02	<b>217.0</b>	<b>217</b>	0.72	0.02	217*	217
mgval-0.25-1C	12	12	27	335	<b>279</b>	279.9	<b>279</b>	0.76	0.26	<b>279.0</b>	<b>279</b>	0.84	0.19	279	279
mgval-0.25-2A	7	12	21	<b>259</b>	<b>259</b>	<b>259.0</b>	<b>259</b>	0.76	0.01	<b>259.0</b>	<b>259</b>	0.98	0.01	259*	259
mgval-0.25-2B	9	9	30	<b>336</b>	<b>336</b>	<b>336.0</b>	<b>336</b>	0.72	0.00	<b>336.0</b>	<b>336</b>	0.92	0.00	336*	336
mgval-0.25-2C	12	10	26	528	<b>480</b>	<b>480.0</b>	<b>480</b>	0.56	0.02	<b>480.0</b>	<b>480</b>	0.66	0.04	480	480
mgval-0.25-3A	9	11	24	<b>89</b>	<b>89</b>	<b>89.0</b>	<b>89</b>	0.55	0.00	<b>89.0</b>	<b>89</b>	0.55	0.01	89*	89
mgval-0.25-3B	8	12	21	<b>125</b>	<b>125</b>	<b>125.0</b>	<b>125</b>	0.57	0.00	<b>125.0</b>	<b>125</b>	0.78	0.00	125*	125
mgval-0.25-3C	10	13	18	161	<b>153</b>	<b>153.0</b>	<b>153</b>	0.42	0.00	<b>153.0</b>	<b>153</b>	0.50	0.01	153*	153
mgval-0.25-4A	19	19	51	<b>514</b>	<b>514</b>	<b>514.0</b>	<b>514</b>	1.85	0.05	<b>514.0</b>	<b>514</b>	2.40	0.05	514*	514
mgval-0.25-4B	20	14	62	541	<b>537</b>	<b>537.0</b>	<b>537</b>	1.74	0.02	<b>537.0</b>	<b>537</b>	2.36	0.07	537*	537
mgval-0.25-4C	24	15	61	549	<b>525</b>	<b>525.0</b>	<b>525</b>	1.90	0.04	<b>525.0</b>	<b>525</b>	2.38	0.07	525*	525
mgval-0.25-4D	19	15	62	724	683	<b>675.0</b>	<b>675</b>	1.81	0.25	<b>675.0</b>	<b>675</b>	2.17	0.57	683*	<u>675</u>
mgval-0.25-5A	21	16	55	<b>485</b>	<b>485</b>	<b>485.0</b>	<b>485</b>	1.54	0.03	<b>485.0</b>	<b>485</b>	1.81	0.06	485*	485
mgval-0.25-5B	18	26	42	500	<b>493</b>	<b>493.0</b>	<b>493</b>	1.65	0.05	<b>493.0</b>	<b>493</b>	2.04	0.08	493*	493
mgval-0.25-5C	21	12	60	599	<b>584</b>	<b>584.0</b>	<b>584</b>	1.60	0.03	<b>584.0</b>	<b>584</b>	2.11	0.07	584*	584
mgval-0.25-5D	17	21	47	681	644	644.0	644	1.39	0.39	643.8	<b>642</b>	1.80	0.33	644*	<u>642</u>
mgval-0.25-6A	16	16	35	<b>274</b>	<b>274</b>	<b>274.0</b>	<b>274</b>	0.93	0.02	<b>274.0</b>	<b>274</b>	0.91	0.03	274*	274
mgval-0.25-6B	14	16	33	<b>263</b>	<b>263</b>	<b>263.0</b>	<b>263</b>	0.96	0.02	<b>263.0</b>	<b>263</b>	1.12	0.03	263*	263
mgval-0.25-6C	16	17	33	337	<b>324</b>	<b>324.0</b>	<b>324</b>	0.76	0.03	<b>324.0</b>	<b>324</b>	0.87	0.03	324	324
mgval-0.25-7A	20	27	37	<b>297</b>	<b>297</b>	<b>297.0</b>	<b>297</b>	1.33	0.05	<b>297.0</b>	<b>297</b>	1.31	0.06	297*	297
mgval-0.25-7B	18	18	49	<b>355</b>	<b>355</b>	<b>355.0</b>	<b>355</b>	1.63	0.04	<b>355.0</b>	<b>355</b>	1.66	0.04	355*	355
mgval-0.25-7C	18	21	46	407	<b>378</b>	<b>378.0</b>	<b>378</b>	1.31	0.09	<b>378.0</b>	<b>378</b>	1.40	0.05	378	378
mgval-0.25-8A	16	15	57	<b>510</b>	<b>510</b>	<b>510.0</b>	<b>510</b>	2.07	0.07	<b>510.0</b>	<b>510</b>	1.90	0.07	510*	510
mgval-0.25-8B	16	20	48	<b>423</b>	<b>423</b>	<b>423.0</b>	<b>423</b>	1.48	0.03	<b>423.0</b>	<b>423</b>	1.97	0.04	423*	423
mgval-0.25-8C	16	21	41	591	545	544.9	544	1.30	0.17	544.7	544	1.61	0.32	545	<u>544</u>
mgval-0.25-9A	23	24	75	<b>371</b>	<b>371</b>	<b>371.0</b>	<b>371</b>	2.14	0.09	<b>371.0</b>	<b>371</b>	2.14	0.12	371*	371
mgval-0.25-9B	22	33	57	363	<b>358</b>	<b>358.0</b>	<b>358</b>	2.43	0.06	<b>358.0</b>	<b>358</b>	2.72	0.11	358*	358
mgval-0.25-9C	26	31	62	369	365	364.1	364	3.22	0.52	364.0	364	3.25	0.40	365	<u>364</u>
mgval-0.25-9D	24	28	69	478	429	426.7	425	2.52	0.93	426.9	425	3.57	1.33	429	<u>424</u>
mgval-0.25-10A	26	24	79	<b>492</b>	<b>492</b>	<b>492.0</b>	<b>492</b>	3.26	0.12	<b>492.0</b>	<b>492</b>	3.58	0.29	492*	492
mgval-0.25-10B	24	24	75	<b>528</b>	<b>528</b>	<b>528.0</b>	<b>528</b>	3.18	0.12	<b>528.0</b>	<b>528</b>	3.47	0.23	528*	528
mgval-0.25-10C	23	27	75	501	<b>483</b>	<b>483.0</b>	<b>483</b>	3.06	0.34	<b>483.0</b>	<b>483</b>	3.49	0.29	483	483
mgval-0.25-10D	23	31	65	616	567	566.7	566	2.24	0.95	567.4	566	3.20	0.92	567*	<u>566</u>
mgval-0.30-1A	15	14	24	<b>170</b>	<b>170</b>	<b>170.0</b>	<b>170</b>	0.55	0.01	<b>170.0</b>	<b>170</b>	0.52	0.01	170*	170
mgval-0.30-1B	12	9	26	<b>194</b>	<b>194</b>	<b>194.0</b>	<b>194</b>	0.80	0.02	<b>194.0</b>	<b>194</b>	0.89	0.02	194*	194
mgval-0.30-1C	12	11	25	280	<b>270</b>	<b>270.0</b>	<b>270</b>	0.61	0.09	<b>270.0</b>	<b>270</b>	0.61	0.03	270	270
mgval-0.30-2A	12	11	19	<b>233</b>	<b>233</b>	<b>233.0</b>	<b>233</b>	0.75	0.00	<b>233.0</b>	<b>233</b>	0.98	0.00	233*	233
mgval-0.30-2B	13	8	28	<b>347</b>	<b>347</b>	<b>347.0</b>	<b>347</b>	0.93	0.06	<b>347.0</b>	<b>347</b>	1.06	0.11	347*	347
mgval-0.30-2C	12	9	24	542	<b>495</b>	<b>495.0</b>	<b>495</b>	0.60	0.10	<b>495.0</b>	<b>495</b>	0.61	0.07	495	495
mgval-0.30-3A	13	10	23	<b>105</b>	<b>105</b>	<b>105.0</b>	<b>105</b>	0.60	0.00	<b>105.0</b>	<b>105</b>	0.61	0.00	105*	105
mgval-0.30-3B	10	11	20	<b>115</b>	<b>115</b>	<b>115.0</b>	<b>115</b>	0.53	0.00	<b>115.0</b>	<b>115</b>	0.72	0.00	115*	115
mgval-0.30-3C	12	12	17	156	<b>153</b>	<b>153.0</b>	<b>153</b>	0.44	0.01	<b>153.0</b>	<b>153</b>	0.55	0.04	153	153
mgval-0.30-4A	21	18	48	<b>477</b>	<b>477</b>	<b>477.0</b>	<b>477</b>	2.06	0.10	<b>477.0</b>	<b>477</b>	2.29	0.06	477*	477
mgval-0.30-4B	27	13	58	537	<b>533</b>	<b>533.0</b>	<b>533</b>	2.29	0.29	<b>533.0</b>	<b>533</b>	2.71	0.37	533	533
mgval-0.30-4C	27	14	57	513	500	499.2	498	1.88	0.26	499.2	498	2.50	0.32	500	<u>498</u>
mgval-0.30-4D	22	14	58	718	<b>653</b>	<b>653.0</b>	<b>653</b>	1.43	0.44	<b>653.0</b>	<b>653</b>	1.85	0.26	653	653
mgval-0.30-5A	20	15	51	<b>445</b>	<b>445</b>	<b>445.0</b>	<b>445</b>	1.64	0.02	<b>445.0</b>	<b>445</b>	1.83	0.05	445*	445
mgval-0.30-5B	20	24	39	492	<b>490</b>	<b>490.0</b>	<b>490</b>	2.04	0.19	<b>490.0</b>	<b>490</b>	2.13	0.24	490	490
mgval-0.30-5C	20	11	56	568	553	553.0	553	1.50	0.03	553.0	553	1.88	0.07	553	<u>551</u>
mgval-0.30-5D	22	20	44	675	621	618.3	618	1.65	0.75	616.9	<b>616</b>	2.55	1.09	621	<u>616</u>
mgval-0.30-6A	19	15	32	<b>252</b>	<b>252</b>	<b>252.0</b>	<b>252</b>	1.06	0.01	<b>252.0</b>	<b>252</b>	1.04	0.01	252*	252
mgval-0.30-6B	19	15	30	268	<b>262</b>	<b>262.0</b>	<b>262</b>	1.30	0.24	<b>262.0</b>	<b>262</b>	1.18	0.05	262*	262
mgval-0.30-6C	17	16	31	339	320	317.6	317	0.95	0.28	317.0	317	1.07	0.20	320	<u>317</u>
mgval-0.30-7A	17	25	35	<b>324</b>	<b>324</b>	<b>324.0</b>	<b>324</b>	1.20	0.03	<b>324.0</b>	<b>324</b>	1.14	0.04	324*	324
mgval-0.30-7B	19	17	46	<b>344</b>	<b>344</b>	<b>344.0</b>	<b>344</b>	1.64	0.03	<b>344.0</b>	<b>344</b>	1.73	0.05	344*	344
mgval-0.30-7C	23	19	43	380	<b>354</b>	<b>354.0</b>	<b>354</b>	1.09	0.02	<b>354.0</b>	<b>354</b>	1.24	0.04	354	354
mgval-0.30-8A	21	14	53	<b>431</b>	<b>431</b>	<b>431.0</b>	<b>431</b>	1.81	0.03	<b>431.0</b>	<b>431</b>	2.00	0.05	431*	431
mgval-0.30-8B	21	18	44	408	<b>400</b>	<b>400.0</b>	<b>400</b>	1.81	0.05	<b>400.0</b>	<b>400</b>	1.78	0.07	400*	400
mgval-0.30-8C	18	19	38	570	<b>522</b>	<b>522.0</b>	<b>522</b>	1.13	0.22	<b>522.0</b>	<b>522</b>	1.26	0.11	522	522
mgval-0.30-9A	26	22	70	<b>357</b>	<b>357</b>	<b>357.0</b>	<b>357</b>	2.26	0.10	<b>357.0</b>	<b>357</b>	2.16	0.11	357*	357
mgval-0.30-9B	27	30	53	356	<b>348</b>	<b>348.0</b>	<b>348</b>	2.30	0.06	<b>348.0</b>	<b>348</b>	2.46	0.10	348*	348
mgval-0.30-9C	25	29	58	347	<b>335</b>	<b>335.0</b>	<b>335</b>	2.26	0.06	<b>335.0</b>	<b>335</b>	2.68	0.13	335*	335
mgval-0.30-9D	31	26	65	475	430	428.0	428	2.58	0.78	429.7	428	3.92	1.72	430	<u>428</u>
mgval-0.30-10A	31	22	74	<b>484</b>	<b>484</b>	<b>484.0</b>	<b>484</b>	3.71	1.07	484.2	484	4.57	1.50	484*	484
mgval-0.30-10B	30	23	70	<b>441</b>	<b>441</b>	<b>441.0</b>	<b>441</b>	2.66	0.06	<b>441.0</b>	<b>441</b>	3.35	0.17	441*	441
mgval-0.30-10C	30	25	70	483	<b>475</b>	<b>475.0</b>	<b>475</b>	3.49	0.15	<b>475.0</b>	<b>475</b>	3.82	0.70	475*	475
mgval-0.30-10D	32	29	60	575	539	538.7	537	2.58	0.59	539.0	539	3.05	0.79	539	537

**Table EC.11** Results for the MCGRP – MGVAL instances (continued)

Inst	$ N_R $	$ E_R $	$ A_R $	BLMV14	DHDI14	ILS				UHGS				BKS	
				Single	Single	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Old	New
mgval-0.35-1A	12	13	22	<b>158</b>	<b>158</b>	<b>158.0</b>	<b>158</b>	0.47	0.00	<b>158.0</b>	<b>158</b>	0.46	0.00	158*	158
mgval-0.35-1B	16	8	24	<b>192</b>	<b>192</b>	<b>192.0</b>	<b>192</b>	0.81	0.02	<b>192.0</b>	<b>192</b>	0.70	0.02	192*	192
mgval-0.35-1C	14	11	23	284	290	283.5	<b>283</b>	0.66	0.30	<b>283.0</b>	<b>283</b>	0.80	0.20	284	<u>283</u>
mgval-0.35-2A	12	10	18	<b>286</b>	<b>286</b>	<b>286.0</b>	<b>286</b>	0.52	0.01	<b>286.0</b>	<b>286</b>	0.49	0.00	286*	286
mgval-0.35-2B	13	7	26	<b>326</b>	<b>326</b>	<b>326.0</b>	<b>326</b>	0.74	0.01	<b>326.0</b>	<b>326</b>	0.88	0.01	326*	326
mgval-0.35-2C	14	9	22	523	<b>485</b>	<b>485.0</b>	<b>485</b>	0.52	0.06	<b>485.0</b>	<b>485</b>	0.55	0.03	485	485
mgval-0.35-3A	13	9	21	<b>84</b>	<b>84</b>	<b>84.0</b>	<b>84</b>	0.59	0.00	<b>84.0</b>	<b>84</b>	0.61	0.01	84*	84
mgval-0.35-3B	13	10	18	<b>113</b>	<b>113</b>	<b>113.0</b>	<b>113</b>	0.53	0.00	<b>113.0</b>	<b>113</b>	0.74	0.00	113*	113
mgval-0.35-3C	13	11	16	159	<b>150</b>	<b>150.0</b>	<b>150</b>	0.37	0.00	<b>150.0</b>	<b>150</b>	0.47	0.01	150*	150
mgval-0.35-4A	24	16	44	<b>430</b>	<b>430</b>	<b>430.0</b>	<b>430</b>	1.66	0.02	<b>430.0</b>	<b>430</b>	2.18	0.03	430*	430
mgval-0.35-4B	25	12	53	<b>531</b>	<b>531</b>	<b>531.0</b>	<b>531</b>	1.59	0.02	<b>531.0</b>	<b>531</b>	2.11	0.06	531	531
mgval-0.35-4C	27	13	53	553	<b>516</b>	<b>516.0</b>	<b>516</b>	2.24	0.17	<b>516.0</b>	<b>516</b>	2.34	0.32	516*	516
mgval-0.35-4D	30	13	53	661	<b>643</b>	<b>643.0</b>	<b>643</b>	1.54	0.43	<b>643.0</b>	<b>643</b>	1.96	0.31	643	643
mgval-0.35-5A	20	14	48	<b>454</b>	<b>454</b>	454.4	<b>454</b>	1.78	0.49	454.3	<b>454</b>	2.77	0.71	454*	454
mgval-0.35-5B	23	22	36	468	<b>467</b>	467.4	<b>467</b>	2.09	0.51	467.7	<b>467</b>	2.39	0.57	467	467
mgval-0.35-5C	19	11	52	595	<b>586</b>	<b>586.0</b>	<b>586</b>	1.79	0.60	586.1	<b>586</b>	2.31	0.62	586	586
mgval-0.35-5D	22	18	40	648	578	576.5	<b>576</b>	1.66	0.69	<b>576.0</b>	<b>576</b>	1.93	0.67	578	<u>576</u>
mgval-0.35-6A	20	14	30	<b>248</b>	<b>248</b>	<b>248.0</b>	<b>248</b>	1.08	0.01	<b>248.0</b>	<b>248</b>	1.22	0.02	248*	248
mgval-0.35-6B	20	14	28	<b>250</b>	<b>250</b>	<b>250.0</b>	<b>250</b>	0.95	0.01	<b>250.0</b>	<b>250</b>	1.10	0.02	250*	250
mgval-0.35-6C	17	14	29	326	<b>312</b>	<b>312.0</b>	<b>312</b>	0.76	0.07	<b>312.0</b>	<b>312</b>	0.76	0.05	312	312
mgval-0.35-7A	23	23	32	<b>264</b>	<b>264</b>	<b>264.0</b>	<b>264</b>	1.02	0.03	<b>264.0</b>	<b>264</b>	1.01	0.03	264*	264
mgval-0.35-7B	21	16	42	<b>325</b>	<b>325</b>	<b>325.0</b>	<b>325</b>	1.44	0.02	<b>325.0</b>	<b>325</b>	1.61	0.03	325*	325
mgval-0.35-7C	24	18	40	351	<b>336</b>	<b>336.0</b>	<b>336</b>	1.26	0.10	<b>336.0</b>	<b>336</b>	1.31	0.10	336*	336
mgval-0.35-8A	22	13	49	<b>415</b>	<b>415</b>	<b>415.0</b>	<b>415</b>	1.64	0.11	<b>415.0</b>	<b>415</b>	1.51	0.06	415*	415
mgval-0.35-8B	20	17	41	<b>385</b>	<b>385</b>	<b>385.0</b>	<b>385</b>	1.44	0.03	<b>385.0</b>	<b>385</b>	1.77	0.03	385*	385
mgval-0.35-8C	22	18	35	547	494	490.7	<b>489</b>	1.18	0.45	<b>489.0</b>	<b>489</b>	1.34	0.25	494	<u>489</u>
mgval-0.35-9A	31	20	65	<b>324</b>	<b>324</b>	<b>324.0</b>	<b>324</b>	2.29	0.08	<b>324.0</b>	<b>324</b>	2.21	0.11	324*	324
mgval-0.35-9B	29	28	49	332	<b>331</b>	<b>331.0</b>	<b>331</b>	2.93	0.27	<b>331.0</b>	<b>331</b>	2.63	0.17	331*	331
mgval-0.35-9C	35	27	53	338	<b>328</b>	<b>328.0</b>	<b>328</b>	2.38	0.36	<b>328.0</b>	<b>328</b>	3.17	0.55	328*	328
mgval-0.35-9D	31	24	60	473	430	427.2	<b>426</b>	2.46	0.82	429.0	<b>426</b>	3.22	1.14	430	<u>426</u>
mgval-0.35-10A	34	20	68	<b>475</b>	<b>475</b>	<b>475.0</b>	<b>475</b>	3.43	0.12	<b>475.0</b>	<b>475</b>	3.47	0.15	475*	475
mgval-0.35-10B	32	21	65	463	<b>461</b>	<b>461.0</b>	<b>461</b>	2.54	0.07	<b>461.0</b>	<b>461</b>	3.06	0.11	461*	461
mgval-0.35-10C	34	23	65	448	<b>430</b>	<b>430.0</b>	<b>430</b>	3.01	0.38	<b>430.0</b>	<b>430</b>	3.59	0.72	430	430
mgval-0.35-10D	31	27	56	566	<b>523</b>	<b>523.0</b>	<b>523</b>	2.13	0.24	523.4	<b>523</b>	2.79	0.70	523	523
mgval-0.40-1A	15	12	21	<b>165</b>	<b>165</b>	<b>165.0</b>	<b>165</b>	0.53	0.01	<b>165.0</b>	<b>165</b>	0.48	0.01	165*	165
mgval-0.40-1B	14	7	22	<b>196</b>	<b>196</b>	<b>196.0</b>	<b>196</b>	0.68	0.01	<b>196.0</b>	<b>196</b>	0.73	0.01	196*	196
mgval-0.40-1C	15	10	21	272	<b>263</b>	<b>263.0</b>	<b>263</b>	0.67	0.15	263.2	<b>263</b>	0.78	0.23	263	263
mgval-0.40-2A	13	9	16	<b>222</b>	<b>222</b>	<b>222.0</b>	<b>222</b>	0.56	0.00	<b>222.0</b>	<b>222</b>	0.79	0.00	222*	222
mgval-0.40-2B	18	7	24	<b>311</b>	<b>311</b>	<b>311.0</b>	<b>311</b>	0.67	0.00	<b>311.0</b>	<b>311</b>	0.89	0.00	311*	311
mgval-0.40-2C	14	8	21	485	<b>469</b>	<b>469.0</b>	<b>469</b>	0.44	0.01	<b>469.0</b>	<b>469</b>	0.54	0.01	469	469
mgval-0.40-3A	13	9	19	<b>86</b>	<b>86</b>	<b>86.0</b>	<b>86</b>	0.61	0.00	<b>86.0</b>	<b>86</b>	0.56	0.00	86*	86
mgval-0.40-3B	14	9	17	<b>110</b>	<b>110</b>	<b>110.0</b>	<b>110</b>	0.51	0.01	<b>110.0</b>	<b>110</b>	0.64	0.01	110*	110
mgval-0.40-3C	13	10	15	157	<b>148</b>	<b>148.0</b>	<b>148</b>	0.35	0.01	<b>148.0</b>	<b>148</b>	0.43	0.01	148	148
mgval-0.40-4A	26	15	41	<b>400</b>	<b>400</b>	<b>400.0</b>	<b>400</b>	1.80	0.05	<b>400.0</b>	<b>400</b>	2.18	0.07	400*	400
mgval-0.40-4B	29	11	49	<b>423</b>	<b>423</b>	<b>423.0</b>	<b>423</b>	1.51	0.03	<b>423.0</b>	<b>423</b>	2.02	0.05	423	423
mgval-0.40-4C	28	12	49	487	<b>462</b>	<b>462.0</b>	<b>462</b>	1.62	0.05	<b>462.0</b>	<b>462</b>	1.94	0.07	462	462
mgval-0.40-4D	27	12	49	669	622	<b>620.0</b>	<b>620</b>	1.79	0.75	622.4	<b>620</b>	2.61	1.17	622	<u>620</u>
mgval-0.40-5A	25	13	44	<b>426</b>	<b>426</b>	<b>426.0</b>	<b>426</b>	1.45	0.02	<b>426.0</b>	<b>426</b>	1.71	0.04	426*	426
mgval-0.40-5B	23	21	33	428	<b>424</b>	<b>424.0</b>	<b>424</b>	1.43	0.03	<b>424.0</b>	<b>424</b>	1.75	0.05	424	424
mgval-0.40-5C	26	10	48	539	527	524.3	<b>524</b>	1.61	0.52	525.9	<b>524</b>	2.25	0.53	527	<u>524</u>
mgval-0.40-5D	25	17	37	665	608	605.3	604	1.43	0.73	603.8	<b>602</b>	2.68	1.41	608	<u>602</u>
mgval-0.40-6A	20	13	28	<b>224</b>	<b>224</b>	<b>224.0</b>	<b>224</b>	1.11	0.02	<b>224.0</b>	<b>224</b>	0.94	0.03	224*	224
mgval-0.40-6B	19	13	26	<b>211</b>	<b>211</b>	<b>211.0</b>	<b>211</b>	0.82	0.02	<b>211.0</b>	<b>211</b>	0.98	0.03	211*	211
mgval-0.40-6C	22	13	27	316	<b>312</b>	<b>312.0</b>	<b>312</b>	0.68	0.01	<b>312.0</b>	<b>312</b>	0.81	0.01	312	312
mgval-0.40-7A	25	21	30	<b>271</b>	<b>271</b>	<b>271.0</b>	<b>271</b>	0.95	0.02	<b>271.0</b>	<b>271</b>	0.94	0.02	271*	271
mgval-0.40-7B	23	15	39	<b>270</b>	<b>270</b>	<b>270.0</b>	<b>270</b>	1.34	0.03	<b>270.0</b>	<b>270</b>	1.47	0.03	270*	270
mgval-0.40-7C	27	16	37	336	332	<b>330.0</b>	<b>330</b>	1.34	0.37	330.4	<b>330</b>	1.73	0.54	332	<u>330</u>
mgval-0.40-8A	23	12	45	<b>393</b>	<b>393</b>	<b>393.0</b>	<b>393</b>	1.70	0.04	<b>393.0</b>	<b>393</b>	1.83	0.07	393*	393
mgval-0.40-8B	23	16	38	372	<b>371</b>	<b>371.0</b>	<b>371</b>	1.67	0.06	<b>371.0</b>	<b>371</b>	1.75	0.06	371	371
mgval-0.40-8C	23	16	33	573	<b>517</b>	<b>517.0</b>	<b>517</b>	1.09	0.15	<b>517.0</b>	<b>517</b>	1.27	0.15	517	517
mgval-0.40-9A	35	19	60	<b>341</b>	<b>341</b>	<b>341.0</b>	<b>341</b>	1.96	0.06	<b>341.0</b>	<b>341</b>	1.96	0.11	341	341
mgval-0.40-9B	34	26	45	331	<b>327</b>	<b>327.0</b>	<b>327</b>	2.17	0.05	<b>327.0</b>	<b>327</b>	2.56	0.08	327	327
mgval-0.40-9C	30	25	49	301	<b>295</b>	<b>295.0</b>	<b>295</b>	2.34	0.17	<b>295.0</b>	<b>295</b>	2.46	0.29	295	295
mgval-0.40-9D	39	22	55	414	<b>382</b>	<b>382.0</b>	<b>382</b>	2.12	0.64	382.4	<b>382</b>	3.09	1.10	382	382
mgval-0.40-10A	36	19	63	<b>406</b>	<b>406</b>	<b>406.0</b>	<b>406</b>	2.70	0.06	<b>406.0</b>	<b>406</b>	3.17	0.11	406*	406
mgval-0.40-10B	34	19	60	439	<b>433</b>	<b>433.0</b>	<b>433</b>	2.56	0.09	<b>433.0</b>	<b>433</b>	3.04	0.24	433	433
mgval-0.40-10C	33	21	60	435	433	432.2	<b>432</b>	2.77	0.67	432.3	<b>432</b>	3.57	0.91	433	<u>432</u>
mgval-0.40-10D	35	25	52	521	<b>482</b>	<b>482.0</b>	<b>482</b>	1.91	0.12	<b>482.0</b>	<b>482</b>	2.80	0.76	482	482

Table EC.12 Results for the MCGRP – MGVAL instances (end)

Inst	$ N_R $	$ E_R $	$ A_R $	BLMV14		DHD14		ILS				UHGS				BKS	
				Single	Single	Single	Single	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Old	New
mgval-0.45-1A	17	11	19	<b>168</b>	<b>168</b>	<b>168.0</b>	<b>168</b>	0.52	0.01	<b>168.0</b>	<b>168</b>	0.47	0.01	168*	168		
mgval-0.45-1B	14	7	20	<b>166</b>	<b>166</b>	<b>166.0</b>	<b>166</b>	0.65	0.02	<b>166.0</b>	<b>166</b>	0.73	0.02	166*	166		
mgval-0.45-1C	16	9	19	313	<b>258</b>	<b>258.0</b>	<b>258</b>	0.54	0.04	<b>258.0</b>	<b>258</b>	0.52	0.02	258	258		
mgval-0.45-2A	15	8	15	<b>251</b>	<b>251</b>	<b>251.0</b>	<b>251</b>	0.56	0.00	<b>251.0</b>	<b>251</b>	0.58	0.00	251*	251		
mgval-0.45-2B	18	6	22	<b>314</b>	<b>314</b>	<b>314.0</b>	<b>314</b>	0.83	0.03	<b>314.0</b>	<b>314</b>	0.91	0.05	314*	314		
mgval-0.45-2C	18	7	19	496	<b>462</b>	<b>462.0</b>	<b>462</b>	0.50	0.05	<b>462.0</b>	<b>462</b>	0.51	0.03	462	462		
mgval-0.45-3A	15	8	18	<b>82</b>	<b>82</b>	<b>82.0</b>	<b>82</b>	0.60	0.00	<b>82.0</b>	<b>82</b>	0.82	0.00	82*	82		
mgval-0.45-3B	16	8	15	<b>91</b>	<b>91</b>	<b>91.0</b>	<b>91</b>	0.50	0.01	<b>91.0</b>	<b>91</b>	0.64	0.01	91*	91		
mgval-0.45-3C	16	9	13	<b>143</b>	<b>143</b>	<b>143.0</b>	<b>143</b>	0.33	0.00	<b>143.0</b>	<b>143</b>	0.44	0.00	143	143		
mgval-0.45-4A	29	14	37	<b>381</b>	<b>381</b>	<b>381.0</b>	<b>381</b>	2.06	0.10	<b>381.0</b>	<b>381</b>	2.16	0.09	381*	381		
mgval-0.45-4B	36	10	45	481	<b>471</b>	<b>471.0</b>	<b>471</b>	1.69	0.04	<b>471.0</b>	<b>471</b>	2.12	0.06	471	471		
mgval-0.45-4C	29	11	45	508	<b>480</b>	<b>480.0</b>	<b>480</b>	1.68	0.07	<b>480.0</b>	<b>480</b>	1.89	0.16	480	480		
mgval-0.45-4D	31	11	45	626	<b>577</b>	<b>577.0</b>	<b>577</b>	1.46	0.10	<b>577.0</b>	<b>577</b>	1.47	0.08	577	577		
mgval-0.45-5A	26	12	40	<b>391</b>	392	391.7	<b>391</b>	1.70	0.32	391.4	<b>391</b>	2.45	0.58	391	391		
mgval-0.45-5B	26	19	30	416	416	415.8	<b>414</b>	1.43	0.14	415.8	<b>414</b>	1.62	0.05	416	<u>414</u>		
mgval-0.45-5C	26	9	44	502	<b>492</b>	<b>492.0</b>	<b>492</b>	1.63	0.07	<b>492.0</b>	<b>492</b>	1.77	0.19	492	492		
mgval-0.45-5D	27	15	34	612	<b>552</b>	<b>552.0</b>	<b>552</b>	1.28	0.29	<b>552.0</b>	<b>552</b>	1.34	0.12	552	552		
mgval-0.45-6A	23	12	25	<b>213</b>	<b>213</b>	<b>213.0</b>	<b>213</b>	0.79	0.01	<b>213.0</b>	<b>213</b>	0.94	0.01	213*	213		
mgval-0.45-6B	22	12	24	<b>210</b>	<b>210</b>	<b>210.0</b>	<b>210</b>	0.80	0.01	<b>210.0</b>	<b>210</b>	1.00	0.01	210*	210		
mgval-0.45-6C	22	12	24	308	<b>296</b>	<b>296.0</b>	<b>296</b>	0.66	0.01	<b>296.0</b>	<b>296</b>	0.70	0.02	296	296		
mgval-0.45-7A	27	19	27	<b>261</b>	<b>261</b>	<b>261.0</b>	<b>261</b>	1.19	0.04	<b>261.0</b>	<b>261</b>	1.08	0.03	261*	261		
mgval-0.45-7B	28	13	36	<b>294</b>	<b>294</b>	<b>294.0</b>	<b>294</b>	1.40	0.05	<b>294.0</b>	<b>294</b>	1.62	0.05	294	294		
mgval-0.45-7C	27	15	34	347	<b>336</b>	<b>336.0</b>	<b>336</b>	1.15	0.10	<b>336.0</b>	<b>336</b>	1.20	0.15	336	336		
mgval-0.45-8A	24	11	41	<b>370</b>	<b>370</b>	<b>370.0</b>	<b>370</b>	1.22	0.02	<b>370.0</b>	<b>370</b>	1.38	0.05	370	370		
mgval-0.45-8B	23	14	35	376	<b>360</b>	<b>360.0</b>	<b>360</b>	1.63	0.10	<b>360.0</b>	<b>360</b>	1.64	0.08	360	360		
mgval-0.45-8C	20	15	30	535	498	498.4	<b>496</b>	1.05	0.59	497.3	<b>496</b>	1.45	0.53	498	<u>496</u>		
mgval-0.45-9A	37	17	55	<b>306</b>	<b>306</b>	<b>306.0</b>	<b>306</b>	2.53	0.07	<b>306.0</b>	<b>306</b>	2.45	0.10	306	306		
mgval-0.45-9B	35	24	41	<b>323</b>	<b>323</b>	<b>323.0</b>	<b>323</b>	2.28	0.11	<b>323.0</b>	<b>323</b>	2.43	0.11	323	323		
mgval-0.45-9C	34	23	45	306	<b>291</b>	<b>291.0</b>	<b>291</b>	2.08	0.17	<b>291.0</b>	<b>291</b>	2.24	0.13	291	291		
mgval-0.45-9D	37	20	51	403	387	386.8	<b>386</b>	1.86	0.40	386.9	<b>386</b>	2.15	0.30	387	<u>386</u>		
mgval-0.45-10A	40	17	58	<b>388</b>	<b>388</b>	<b>388.0</b>	<b>388</b>	2.55	0.05	<b>388.0</b>	<b>388</b>	3.09	0.11	388	388		
mgval-0.45-10B	35	18	55	<b>399</b>	<b>399</b>	<b>399.0</b>	<b>399</b>	2.10	0.04	<b>399.0</b>	<b>399</b>	2.63	0.08	399	399		
mgval-0.45-10C	37	19	55	418	403	402.4	<b>401</b>	2.33	0.41	402.8	<b>401</b>	2.79	0.23	403	<u>401</u>		
mgval-0.45-10D	35	23	47	504	487	487.5	<b>486</b>	1.89	0.92	487.6	487	2.70	0.89	487	<u>486</u>		
mgval-0.50-1A	16	10	17	<b>145</b>	<b>145</b>	<b>145.0</b>	<b>145</b>	0.60	0.00	<b>145.0</b>	<b>145</b>	0.65	0.00	145*	145		
mgval-0.50-1B	17	6	19	<b>170</b>	<b>170</b>	<b>170.0</b>	<b>170</b>	0.70	0.01	<b>170.0</b>	<b>170</b>	0.81	0.01	170*	170		
mgval-0.50-1C	14	8	18	<b>261</b>	255	265.0	265	0.56	0.10	265.0	265	0.52	0.05	261	261		
mgval-0.50-2A	16	8	14	<b>248</b>	<b>248</b>	<b>248.0</b>	<b>248</b>	0.63	0.00	<b>248.0</b>	<b>248</b>	0.84	0.00	248*	248		
mgval-0.50-2B	18	6	20	<b>284</b>	<b>284</b>	<b>284.0</b>	<b>284</b>	0.66	0.01	<b>284.0</b>	<b>284</b>	0.80	0.00	284*	284		
mgval-0.50-2C	16	7	17	488	<b>464</b>	<b>464.0</b>	<b>464</b>	0.41	0.01	<b>464.0</b>	<b>464</b>	0.45	0.02	464	464		
mgval-0.50-3A	17	7	16	<b>75</b>	<b>75</b>	<b>75.0</b>	<b>75</b>	0.64	0.00	<b>75.0</b>	<b>75</b>	0.79	0.01	75*	75		
mgval-0.50-3B	15	8	14	<b>107</b>	<b>107</b>	<b>107.0</b>	<b>107</b>	0.39	0.00	<b>107.0</b>	<b>107</b>	0.57	0.00	107*	107		
mgval-0.50-3C	15	9	12	139	<b>137</b>	<b>137.0</b>	<b>137</b>	0.31	0.00	<b>137.0</b>	<b>137</b>	0.39	0.00	137*	137		
mgval-0.50-4A	31	13	34	<b>350</b>	<b>350</b>	<b>350.0</b>	<b>350</b>	1.62	0.03	<b>350.0</b>	<b>350</b>	1.85	0.04	350*	350		
mgval-0.50-4B	32	9	41	419	<b>413</b>	<b>413.0</b>	<b>413</b>	1.83	0.10	<b>413.0</b>	<b>413</b>	1.85	0.07	413	413		
mgval-0.50-4C	32	10	41	512	488	487.7	<b>487</b>	1.46	0.32	487.7	<b>487</b>	2.04	0.33	488	<u>487</u>		
mgval-0.50-4D	32	10	41	613	<b>580</b>	<b>582.8</b>	<b>580</b>	1.45	0.65	<b>580.0</b>	<b>580</b>	2.26	0.94	580	580		
mgval-0.50-5A	27	11	37	<b>367</b>	<b>367</b>	<b>367.0</b>	<b>367</b>	1.48	0.03	<b>367.0</b>	<b>367</b>	1.69	0.05	367*	367		
mgval-0.50-5B	28	17	28	<b>378</b>	<b>378</b>	<b>378.0</b>	<b>378</b>	1.36	0.03	<b>378.0</b>	<b>378</b>	1.60	0.05	378	378		
mgval-0.50-5C	26	8	40	476	<b>457</b>	<b>457.0</b>	<b>457</b>	1.59	0.13	457.6	<b>457</b>	1.83	0.38	457	457		
mgval-0.50-5D	27	14	31	570	541	541.8	<b>539</b>	1.13	0.39	539.8	<b>539</b>	1.80	0.68	541	<u>539</u>		
mgval-0.50-6A	19	11	23	<b>210</b>	<b>210</b>	<b>210.0</b>	<b>210</b>	0.82	0.01	<b>210.0</b>	<b>210</b>	0.98	0.01	210*	210		
mgval-0.50-6B	23	11	22	<b>210</b>	<b>210</b>	<b>210.0</b>	<b>210</b>	0.86	0.02	<b>210.0</b>	<b>210</b>	0.91	0.01	210*	210		
mgval-0.50-6C	22	11	22	306	<b>293</b>	<b>293.0</b>	<b>293</b>	0.64	0.03	<b>293.0</b>	<b>293</b>	0.69	0.03	293	293		
mgval-0.50-7A	31	18	25	<b>248</b>	<b>248</b>	<b>248.0</b>	<b>248</b>	1.40	0.03	<b>248.0</b>	<b>248</b>	1.21	0.04	248*	248		
mgval-0.50-7B	26	12	33	<b>276</b>	<b>276</b>	<b>276.0</b>	<b>276</b>	1.05	0.02	<b>276.0</b>	<b>276</b>	1.19	0.04	276*	276		
mgval-0.50-7C	29	14	31	333	<b>320</b>	<b>320.0</b>	<b>320</b>	1.03	0.06	<b>320.0</b>	<b>320</b>	1.07	0.05	320	320		
mgval-0.50-8A	26	10	38	<b>388</b>	<b>388</b>	<b>388.0</b>	<b>388</b>	1.66	0.05	<b>388.0</b>	<b>388</b>	1.72	0.06	388	388		
mgval-0.50-8B	23	13	32	356	<b>350</b>	<b>350.0</b>	<b>350</b>	1.32	0.04	<b>350.0</b>	<b>350</b>	1.46	0.04	350	350		
mgval-0.50-8C	22	14	27	535	<b>501</b>	<b>501.0</b>	<b>501</b>	0.79	0.03	<b>501.0</b>	<b>501</b>	0.98	0.04	501	501		
mgval-0.50-9A	39	16	50	<b>306</b>	<b>306</b>	<b>306.0</b>	<b>306</b>	2.30	0.08	<b>306.0</b>	<b>306</b>	2.03	0.12	306*	306		
mgval-0.50-9B	37	22	38	<b>278</b>	<b>278</b>	<b>278.0</b>	<b>278</b>	2.04	0.07	<b>278.0</b>	<b>278</b>	2.23	0.08	278	278		
mgval-0.50-9C	38	21	41	303	<b>292</b>	<b>292.0</b>	<b>292</b>	1.96	0.27	<b>292.0</b>	<b>292</b>	2.23	0.14	292	292		
mgval-0.50-9D	38	19	46	408	<b>358</b>	<b>358.0</b>	<b>358</b>	1.87	0.54	<b>358.0</b>	<b>358</b>	2.26	0.51	358	358		
mgval-0.50-10A	40	16	53	<b>385</b>	<b>385</b>	<b>385.0</b>	<b>385</b>	2.85	0.36	<b>385.0</b>	<b>385</b>	2.87	0.20	385	385		
mgval-0.50-10B	44	16	50	371	<b>369</b>	<b>369.0</b>	<b>369</b>	2.57	0.09	<b>369.0</b>	<b>369</b>	2.88	0.16	369	369		
mgval-0.50-10C	40	18	50	416	<b>406</b>	<b>406.0</b>	<b>406</b>	2.24	0.06	<b>406.0</b>	<b>406</b>	2.61	0.12	406	406		
mgval-0.50-10D	46	21	43	492	457	456.9	<b>456</b>	2.17	1.07	457.8	<b>456</b>	3.44	1.45	457	<u>456</u>		
Gap (%)				2.621%	0.072%	0.047%	0.015%					0.047%	0.013%				
T(min)				16.74	60.00			1.41				1.65					
T*(min)				—	3.69			0.15				0.20					
CPU				Xe 3.0G	CPU 3G			Xe 3.07G				Xe 3.07G					

Table EC.13 Results for the MCGRP – CBMix instances

Inst	$ N_R $	$ E_R $	$ A_R $	HKSG12	BLMV14	DHDI14	ILS				UHGS				BKS	
				Best-2	Single	Single	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Old	New
CBMix1	11	0	37	2589	2587	2585	2585.4	2585	0.72	0.06	2570.6	<b>2569</b>	1.05	0.34	2585	<u>2569</u>
CBMix2	36	0	149	12222	12241	11749	11708.5	11643	8.76	4.74	11663.7	<b>11610</b>	16.17	12.19	11749	<u>11610</u>
CBMix3	16	8	55	3643	3643	3614	3612.6	3604	1.40	0.40	3614.7	<b>3612</b>	1.99	0.69	3614	<u>3590</u>
CBMix4	10	75	13	7802	7583	7483	7459.5	7431	2.65	1.15	7437.0	<b>7429</b>	3.80	2.15	7483	<u>7429</u>
CBMix5	23	4	38	4531	4531	4459	4459.0	4459	0.91	0.20	4454.4	<b>4436</b>	1.27	0.26	4459	<u>4436</u>
CBMix6	40	4	64	7087	6968	6969	6846.9	<b>6813</b>	2.01	1.01	6816.9	<b>6813</b>	3.46	1.38	6968	<u>6813</u>
CBMix7	54	8	106	9607	9859	9428	9396.6	9341	5.39	2.79	9400.2	<b>9303</b>	8.66	4.88	9428	<u>9260</u>
CBMix8	63	6	108	10669	10658	10338	10292.7	10257	6.81	3.24	10270.1	<b>10162</b>	15.42	11.62	10338	<u>10148</u>
CBMix9	6	39	5	4130	4060	3991	3981.4	<b>3965</b>	0.80	0.24	3973.0	<b>3965</b>	1.35	0.62	3991	<u>3965</u>
CBMix10	4	94	9	7794	7755	7525	7486.7	7463	2.96	2.06	<b>7462.0</b>	<b>7462</b>	3.33	1.46	7525	<u>7462</u>
CBMix11	65	6	11	4525	4561	4484	4480.5	4466	1.47	0.36	4470.9	<b>4456</b>	2.61	1.22	4484	<u>4456</u>
CBMix12	1	0	52	3235	<b>3138</b>	<b>3138</b>	<b>3138.0</b>	<b>3138</b>	0.71	0.02	<b>3138.0</b>	<b>3138</b>	0.88	0.03	3138	3138
CBMix13	79	2	60	9135	9110	9037	8985.1	8947	4.20	1.94	8970.3	<b>8937</b>	6.76	3.97	9037	<u>8934</u>
CBMix14	93	0	0	8579	8671	8473	8442.4	8420	2.25	1.12	8448.2	<b>8438</b>	2.64	1.08	8473	<u>8413</u>
CBMix15	0	91	0	8371	8359	8221	8206.3	8179	1.70	0.71	8177.6	<b>8164</b>	2.46	1.06	8221	<u>8164</u>
CBMix16	36	0	133	9022	8933	8742	8627.7	<b>8605</b>	4.68	2.23	8607.3	<b>8605</b>	7.46	3.98	8742	<u>8605</u>
CBMix17	16	16	31	4097	4037	<b>4034</b>	<b>4034.0</b>	<b>4034</b>	0.68	0.03	<b>4034.0</b>	<b>4034</b>	0.81	0.02	4034	4034
CBMix18	39	0	88	7133	7254	7052	7011.6	<b>6979</b>	2.94	1.73	6997.8	<b>6979</b>	4.21	1.89	7052	<u>6962</u>
CBMix19	61	9	142	16692	16554	16155	16112.1	16072	12.01	6.93	15982.7	<b>15895</b>	17.58	12.88	16155	<u>15895</u>
CBMix20	38	2	33	4859	4885	<b>4738</b>	4758.8	<b>4738</b>	1.35	0.70	4752.8	<b>4738</b>	2.10	0.93	4738	<u>4738</u>
CBMix21	55	68	57	18809	18509	17875	17815.7	17739	6.89	3.67	17749.0	<b>17638</b>	12.07	8.50	17875	<u>17597</u>
CBMix22	7	10	25	<b>1941</b>	<b>1941</b>	<b>1941</b>	<b>1941.0</b>	<b>1941</b>	0.53	0.00	<b>1941.0</b>	<b>1941</b>	0.70	0.01	1941	1941
CBMix23	3	2	15	<b>780</b>	<b>780</b>	<b>780</b>	<b>780.0</b>	<b>780</b>	0.12	0.00	<b>780.0</b>	<b>780</b>	0.18	0.00	780	780
Gap (%)				3.083%	2.705%	0.891%	0.574%	0.287%			0.361%	0.088%				
T(min)				120.00	44.72	60.00			3.13				5.09			
T*(min)				56.92	—	19.60							3.09			
CPU				CPU 3G	Xe 3.0G	CPU 3G			Xe 3.07G				Xe 3.07G			

Table EC.14 Results for the MCGRP – BHW instances

Inst	$ N_R $	$ E_R $	$ A_R $	HKSG12	DHDI14	ILS				UHGS				BKS	
				Best-2	Single	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Old	New
BHW1	7	11	11	<b>337</b>	<b>337</b>	<b>337.0</b>	<b>337</b>	0.20	0.00	<b>337.0</b>	<b>337</b>	0.27	0.00	337*	337
BHW2	4	0	25	<b>470</b>	<b>470</b>	<b>470.0</b>	<b>470</b>	0.20	0.00	<b>470.0</b>	<b>470</b>	0.19	0.00	470*	470
BHW3	5	8	7	<b>415</b>	<b>415</b>	<b>415.0</b>	<b>415</b>	0.12	0.00	<b>415.0</b>	<b>415</b>	0.16	0.00	415*	415
BHW4	6	0	44	<b>240</b>	<b>240</b>	<b>240.0</b>	<b>240</b>	0.46	0.00	<b>240.0</b>	<b>240</b>	0.50	0.00	240*	240
BHW5	30	0	132	506	<b>502</b>	<b>502.0</b>	<b>502</b>	2.86	0.07	<b>502.0</b>	<b>502</b>	3.65	0.10	502*	502
BHW6	15	37	58	<b>388</b>	<b>388</b>	<b>388.0</b>	<b>388</b>	1.26	0.02	<b>388.0</b>	<b>388</b>	1.35	0.02	388*	388
BHW7	35	0	194	1094	1070	1066.0	1062	8.34	4.25	1065.0	<b>1060</b>	12.15	6.99	1070	<u>1054</u>
BHW8	20	0	97	672	<b>668</b>	668.4	<b>668</b>	1.99	0.57	668.4	<b>668</b>	3.35	0.98	668	668
BHW9	10	26	142	920	875	871.0	867	4.50	1.82	866.0	<b>865</b>	7.06	3.36	875	<u>863</u>
BHW10	40	0	102	8596	8524	8506.4	<b>8482</b>	3.21	1.54	8483.0	<b>8482</b>	3.54	0.88	8524	<u>8480</u>
BHW11	20	0	51	5023	4914	<b>4883.0</b>	<b>4883</b>	1.80	0.69	<b>4883.0</b>	<b>4883</b>	1.37	0.16	4914	<u>4883</u>
BHW12	40	0	75	11042	10887	10871.3	<b>10849</b>	2.53	1.16	10858.9	<b>10849</b>	3.72	1.75	10887	<u>10849</u>
BHW13	25	0	150	14510	14346	14326.9	14310	5.81	2.80	14318.0	<b>14301</b>	7.92	4.59	14346	<u>14301</u>
BHW14	25	0	196	25194	24833	24802.8	24730	13.10	7.66	24739.4	<b>24692</b>	18.58	13.98	24833	<u>24668</u>
BHW15	30	0	98	15509	15354	15361.9	15341	4.01	1.80	15351.6	<b>15331</b>	6.90	4.65	15354	<u>15331</u>
BHW16	30	0	380	44527	43948	43412.5	43147	43.08	31.30	43293.3	<b>43121</b>	56.82	51.39	43948	<u>43058</u>
BHW17	50	0	190	26768	26235	26001.9	25898	18.03	13.03	25965.3	<b>25876</b>	24.65	19.23	26235	<u>25850</u>
BHW18	20	0	174	15833	15170	15203.4	15118	6.92	3.49	15163.0	<b>15098</b>	11.72	7.77	15170	<u>15060</u>
BHW19	20	0	87	9480	9388	9393.6	<b>9379</b>	2.64	1.50	9382.9	<b>9379</b>	5.37	3.45	9388	<u>9379</u>
BHW20	50	51	192	16625	16291	16127.2	<b>16038</b>	17.15	8.24	16142.0	16060	29.69	23.16	16291	<u>16019</u>
Gap (%)				1.976%	0.581%	0.338%	0.126%			0.230%	0.084%				
T(min)				120.00	60.00										
T*(min)				60.07	21.44										
CPU				CPU 3G	CPU 3G	Xe 3.07G				Xe 3.07G					



Table EC.15 Results for the MCGRP – DI-NEARP instances

Inst	$ N_R $	$ E_R $	$ A_R $	HKSG12		ILS				UHGS				BKS	
				Best-2	Single	Avg-10	Best-10	T	T*	Avg-10	Best-10	T	T*	Old	New
n240-Q2k	120	120	0	24371	23807	23834.1	<b>23782</b>	20.98	15.04	23799.3	<b>23782</b>	35.45	24.53	23807	<u>23782</u>
n240-Q4k	120	120	0	18352	18197	<b>18181.0</b>	<b>18181</b>	22.11	10.73	18184.0	<b>18181</b>	22.89	8.84	18197	<u>18181</u>
n240-Q8k	120	120	0	15937	15884	<b>15865.0</b>	<b>15865</b>	27.10	2.25	<b>15865.0</b>	<b>15865</b>	34.36	14.02	15884	<u>15865</u>
n240-Q16k	120	120	0	14953	<b>14717</b>	<b>14717.0</b>	<b>14717</b>	24.79	1.28	<b>14717.0</b>	<b>14717</b>	12.35	0.89	14717	14717
n422-Q2k	302	120	0	18990	18943	18905.9	<b>18902</b>	51.41	27.62	18903.0	<b>18902</b>	45.82	26.89	18943	<u>18902</u>
n422-Q4k	302	120	0	15987	15869	15851.6	<b>15849</b>	60.11	32.62	<b>15849.0</b>	<b>15849</b>	47.28	17.42	15869	<u>15849</u>
n422-Q8k	302	120	0	14627	<b>14442</b>	<b>14442.0</b>	<b>14442</b>	59.23	4.29	<b>14442.0</b>	<b>14442</b>	37.22	3.43	14442	14442
n422-Q16k	302	120	0	14357	<b>14339</b>	<b>14339.0</b>	<b>14339</b>	60.17	7.65	<b>14339.0</b>	<b>14339</b>	40.53	3.58	14339	14339
n442-Q2k	294	148	0	51656	51052	50442.3	50402	56.26	24.80	50416.2	<b>50386</b>	52.90	32.17	51052	<u>50382</u>
n442-Q4k	294	148	0	45605	44952	44931.6	44927	49.60	20.02	44931.2	<b>44926</b>	46.74	22.03	44952	<u>44926</u>
n442-Q8k	294	148	0	44652	43264	<b>43247.0</b>	<b>43247</b>	60.16	18.97	<b>43247.0</b>	<b>43247</b>	49.70	15.10	43264	<u>43247</u>
n442-Q16k	294	148	0	42797	42683	<b>42681.0</b>	<b>42681</b>	60.19	3.80	<b>42681.0</b>	<b>42681</b>	55.07	3.21	42683	<u>42681</u>
n477-Q2k	203	274	0	23124	22896	22879.0	<b>22868</b>	47.97	28.36	22876.4	<b>22868</b>	46.54	22.46	22896	<u>22868</u>
n477-Q4k	203	274	0	20198	20035	<b>19950.0</b>	<b>19950</b>	56.66	7.42	19950.2	<b>19950</b>	46.40	15.05	20035	<u>19950</u>
n477-Q8k	203	274	0	18561	<b>18490</b>	<b>18490.0</b>	<b>18490</b>	60.16	5.03	<b>18490.0</b>	<b>18490</b>	60.28	33.68	18490	18490
n477-Q16k	203	274	0	18105	18040	<b>17930.0</b>	<b>17930</b>	60.22	4.39	<b>17930.0</b>	<b>17930</b>	56.81	5.23	18040	<u>17930</u>
n699-Q2k	335	364	0	59817	58948	58668.1	58595	60.21	34.97	58708.6	<b>58422</b>	60.28	56.13	58948	<u>58422</u>
n699-Q4k	335	364	0	40473	40124	39723.5	39656	60.21	32.58	39717.0	<b>39637</b>	60.31	47.58	40124	<u>39608</u>
n699-Q8k	335	364	0	30992	30799	30532.6	<b>30531</b>	60.24	41.47	30536.1	<b>30531</b>	60.40	40.87	30799	<u>30531</u>
n699-Q16k	335	364	0	27028	26999	26740.5	<b>26718</b>	60.29	32.78	26725.4	<b>26718</b>	60.53	46.47	26999	<u>26703</u>
n833-Q2k	347	486	0	56877	56102	55818.4	<b>55507</b>	60.30	36.99	55674.1	55556	60.44	57.36	56102	<u>55335</u>
n833-Q4k	347	486	0	42407	41192	40893.7	40762	60.34	35.30	40747.4	<b>40688</b>	60.53	56.02	41192	<u>40620</u>
n833-Q8k	347	486	0	35267	34812	34467.5	34374	60.41	45.29	34392.8	<b>34336</b>	60.69	54.92	34812	<u>34275</u>
n833-Q16k	347	486	0	33013	32567	32348.2	32310	60.47	28.09	32277.2	<b>32242</b>	61.05	55.46	32567	<u>32213</u>
Gap (%)				1.640%	0.537%	0.159%	0.074%			0.106%	0.041%				
T(min)				120.00	60.00			58.23				53.48			
T*(min)				92.95	36.32				23.62				30.75		
CPU				CPU 3G	CPU 3G		Xe 3.07G				Xe 3.07G				

**Table EC.16** Results for the CARP with turn penalties – CMMS instances

Inst	$ N_R $	$ E_R $	$ A_R $	ILS			UHGS (10,000 it)			UHGS (20,000 it)			BKS
				Avg-10	Best-10	T	Avg-10	Best-10	T	Avg-10	Best-10	T	
p01-8	0	80	80	10710.9	<b>10687</b>	9.62	10722.3	10697	19.53	10721.6	10696	34.50	10687
p01-16	0	80	80	12504.6	12452	6.38	12461.4	12434	15.58	12439.4	<b>12414</b>	33.71	12414
p02-11	0	80	140	13631.4	<b>13596</b>	21.27	13636.0	13622	36.86	13639.4	13608	65.69	13596
p02-22	0	80	140	14936.1	<b>14891</b>	18.48	14912.4	14893	27.31	14906.2	14893	61.76	14891
p03-14	0	80	200	17749.9	<b>17725</b>	32.01	17771.7	17743	48.64	17760.2	17733	90.39	17725
p03-28	0	80	200	20236.5	20175	20.99	20183.1	20155	49.60	20145.4	<b>20108</b>	106.49	20108
p04-9	0	100	88	11425.2	11404	9.95	11426.6	11405	18.14	11418.0	<b>11401</b>	40.63	11401
p04-17	0	100	88	12461.3	<b>12438</b>	10.26	12472.6	12442	17.00	12462.4	12446	29.34	12438
p05-12	0	100	154	14689.7	<b>14672</b>	22.34	14708.8	14678	44.52	14692.2	14674	71.84	14672
p05-24	0	100	154	16671.6	16599	18.76	16574.9	16532	37.12	16554.5	<b>16528</b>	66.97	16528
p06-15	0	100	220	18707.3	<b>18674</b>	39.32	18742.1	18723	56.21	18728.9	18688	102.82	18674
p06-30	0	100	220	21597.2	21469	28.14	21551.9	21462	56.21	21499.3	<b>21458</b>	102.08	21458
p07-10	0	120	104	13915.9	13870	10.85	13886.9	<b>13853</b>	27.75	13874.7	<b>13853</b>	36.17	13853
p07-19	0	120	104	15350.1	15255	9.13	15222.3	15166	20.89	15208.5	<b>15181</b>	39.44	15181
p08-14	0	120	182	19086.6	19051	32.11	19101.8	19057	48.34	19093.9	<b>19049</b>	100.92	19049
p08-27	0	120	182	21598.8	21533	32.91	21590.4	21496	45.09	21542.8	<b>21488</b>	93.41	21488
p09-17	0	120	260	22398.8	<b>22339</b>	55.08	22519.3	22374	58.81	22495.5	22373	111.66	22339
p09-34	0	120	260	24841.5	24681	49.19	24810.5	24725	54.95	24720.9	<b>24615</b>	114.01	24615
p10-11	0	140	132	15957.7	<b>15888</b>	21.66	15971.8	15900	38.76	15970.0	15935	75.80	15888
p10-22	0	140	132	17995.8	17943	18.22	17950.8	17912	29.67	17925.7	<b>17753</b>	60.24	17753
p11-16	0	140	231	23153.8	<b>23091</b>	48.22	23223.5	23103	57.88	23193.2	23102	112.58	23091
p11-32	0	140	231	27943.6	27873	26.91	27867.2	27791	57.84	27810.4	<b>27757</b>	110.01	27757
p12-21	0	140	330	27931.7	27860	60.03	28081.2	27961	58.72	27939.3	<b>27813</b>	115.29	27813
p12-41	0	140	330	32521.1	32430	53.05	32520.6	32368	60.10	32413.2	<b>32318</b>	120.00	32318
p13-12	0	160	148	18538.5	<b>18489</b>	20.26	18567.5	18501	36.98	18528.6	18505	60.09	18489
p13-23	0	160	148	20881.8	20806	21.64	20835.8	20623	37.28	20704.3	<b>20563</b>	83.53	20563
p14-17	0	160	259	25875.2	<b>25833</b>	52.82	25954.9	25846	60.13	25922.2	25850	119.52	25833
p14-34	0	160	259	29594.5	29449	45.56	29549.4	29431	56.34	29527.9	<b>29446</b>	117.21	29446
p15-23	0	160	370	32827.5	32720	60.05	32857.0	32786	60.11	32779.9	<b>32626</b>	120.00	32626
p15-45	0	160	370	38039.0	37915	60.03	38134.1	37963	60.04	37897.1	<b>37749</b>	118.40	37749
p16-12	0	180	160	19748.4	19705	26.76	19739.4	19689	51.94	19717.4	<b>19668</b>	98.09	19668
p16-24	0	180	160	23118.4	22999	20.79	22980.6	22841	42.16	22944.6	<b>22825</b>	83.79	22825
p17-18	0	180	280	26295.1	<b>26170</b>	57.19	26405.2	26278	60.00	26377.8	26280	120.02	26170
p17-36	0	180	280	31417.0	31335	58.08	31445.1	31257	54.71	31368.5	<b>31201</b>	118.28	31201
p18-24	0	180	400	31532.1	31221	59.41	31504.9	31295	60.18	31319.7	<b>31175</b>	120.00	31175
p18-48	0	180	400	36709.3	36587	60.05	36772.5	36543	59.55	36528.0	<b>36424</b>	120.00	36424
p19-13	0	200	168	19542.0	<b>19499</b>	26.63	19526.2	19460	42.15	19521.0	19503	66.71	19499
p19-25	0	200	168	22046.6	21963	26.89	21994.9	21724	42.15	21896.7	<b>21717</b>	84.41	21717
p20-19	0	200	294	27896.1	27670	35.59	27710.1	27551	60.08	27608.6	<b>27480</b>	119.60	27480
p20-38	0	200	294	32987.6	32903	59.44	32992.3	32783	60.00	32863.0	<b>32668</b>	112.23	32668
p21-25	0	200	420	34754.4	34590	60.05	34669.1	34451	60.04	34475.4	<b>34278</b>	120.09	34278
p21-50	0	200	420	41919.9	41726	60.06	41947.6	41605	60.00	41683.4	<b>41504</b>	120.07	41504
Gap (%)				0.645%	0.289%		0.600%	0.184%		0.370%	0.035%		
T(min)						34.91			46.41			90.42	
CPU				Xe 3.07G			Xe 3.07G			Xe 3.07G			

**Table EC.17** Results for the MCGRP with turn penalties – DI-TP instances

Inst	$ N_R $	$ E_R $	$ A_R $	ILS			UHGS (10,000 it)			UHGS (20,000 it)			BKS
				Avg-10	Best-10	T	Avg-10	Best-10	T	Avg-10	Best-10	T	
n80-Q2k	40	40	0	<b>6941.0</b>	<b>6941</b>	2.69	<b>6941.0</b>	<b>6941</b>	2.40	<b>6941.0</b>	<b>6941</b>	4.80	6941
n80-Q4k	40	40	0	<b>6296.0</b>	<b>6296</b>	3.21	<b>6296.0</b>	<b>6296</b>	2.85	<b>6296.0</b>	<b>6296</b>	5.87	6296
n160-Q2k	80	80	0	16734.4	16688	13.54	16706.0	16688	15.55	16693.3	<b>16682</b>	30.12	16682
n160-Q4k	80	80	0	14289.3	<b>14229</b>	16.93	14265.6	14247	20.71	14248.4	<b>14229</b>	37.24	14229
n240-Q2k	120	120	0	30258.8	30181	41.80	30189.4	<b>30141</b>	50.64	30171.3	<b>30141</b>	83.03	30141
n240-Q4k	120	120	0	23658.9	<b>23643</b>	45.26	23657.4	<b>23643</b>	50.64	23652.1	<b>23643</b>	83.59	23643
n240-Q8k	120	120	0	21061.9	<b>21035</b>	58.80	21064.7	<b>21035</b>	60.39	21064.2	<b>21035</b>	118.87	21035
n240-Q16k	120	120	0	20085.8	19942	60.10	19966.9	<b>19851</b>	53.63	20033.5	19877	112.48	19851
n422-Q2k	302	120	0	25847.8	25776	60.09	25760.4	<b>25718</b>	60.34	25760.8	25719	120.13	25718
n422-Q4k	302	120	0	22239.4	22156	60.12	22133.5	22125	60.03	22120.0	<b>22098</b>	120.10	22098
n422-Q8k	302	120	0	20614.5	20560	60.20	20554.8	20526	60.01	20530.5	<b>20523</b>	120.14	20523
n422-Q16k	302	120	0	20516.6	20432	60.70	20413.4	<b>20376</b>	60.93	20402.1	<b>20376</b>	120.01	20376
n442-Q2k	294	148	0	72638.4	72477	60.07	72483.1	72455	58.56	72446.9	<b>72367</b>	117.42	72367
n442-Q4k	294	148	0	65322.5	65243	60.11	<b>65175.0</b>	<b>65175</b>	59.17	<b>65175.0</b>	<b>65175</b>	118.38	65175
n442-Q8k	294	148	0	63005.6	<b>62902</b>	60.22	62975.9	<b>62902</b>	60.42	62983.3	<b>62902</b>	120.02	62902
n442-Q16k	294	148	0	62177.4	<b>62125</b>	60.61	62188.7	<b>62125</b>	63.25	62200.3	<b>62125</b>	120.01	62125
n477-Q2k	203	274	0	31002.0	30924	60.09	30950.1	30876	60.01	30943.2	<b>30835</b>	120.00	30835
n477-Q4k	203	274	0	26980.1	26903	60.14	26885.7	<b>26820</b>	60.17	26852.8	<b>26820</b>	121.02	26820
n477-Q8k	203	274	0	24990.2	24838	60.21	24983.5	24855	60.01	24867.1	<b>24802</b>	121.31	24802
n477-Q16k	203	274	0	24272.9	24104	60.27	24170.4	24086	65.39	24153.6	<b>24053</b>	120.02	24053
n699-Q2k	335	364	0	77302.9	77155	60.16	77206.4	77120	60.01	77053.0	<b>76804</b>	120.20	76804
n699-Q4k	335	364	0	52944.6	52748	60.22	52746.1	<b>52611</b>	60.01	52750.0	52655	120.29	52611
n699-Q8k	335	364	0	41020.1	40902	60.20	41029.0	40953	60.35	40946.5	<b>40883</b>	120.01	40883
n699-Q16k	335	364	0	36253.8	35991	60.32	36275.0	36045	64.91	36100.1	<b>35955</b>	120.01	35955
n833-Q2k	347	486	0	73094.2	72835	60.34	73241.8	72693	60.27	72700.4	<b>72310</b>	120.02	72310
n833-Q4k	347	486	0	53514.9	53179	60.31	53590.9	53147	60.02	53288.1	<b>52906</b>	120.05	52906
n833-Q8k	347	486	0	45403.7	45035	60.34	45291.2	45051	60.02	44879.5	<b>44640</b>	124.05	44640
n833-Q16k	347	486	0	42958.7	42756	60.44	42716.6	<b>42167</b>	60.02	42479.5	42282	120.02	42167
Gap (%)				0.597%	0.256%		0.414%	0.130%		0.251%	0.018%		
T(min)						51.70			52.53			102.83	
CPU					Xe 3.07G			Xe 3.07G			Xe 3.07G		







**Table EC.21** Results for the MDCARP – VAL instances

Inst	$ E $	$d$	UHGS				BKS
			Avg-10	Best-10	T	T*	New
1A	39	3	<b>173.0</b>	<b>173</b>	0.51	0.00	173
1B	39	3	<b>173.0</b>	<b>173</b>	0.48	0.00	173
1C	39	3	<b>192.0</b>	<b>192</b>	0.46	0.02	192
2A	34	3	<b>217.0</b>	<b>217</b>	0.40	0.00	217
2B	34	3	<b>217.0</b>	<b>217</b>	0.38	0.00	217
2C	34	3	<b>289.0</b>	<b>289</b>	0.39	0.01	289
3A	35	3	<b>77.0</b>	<b>77</b>	0.41	0.00	77
3B	35	3	<b>77.0</b>	<b>77</b>	0.40	0.00	77
3C	35	3	<b>85.0</b>	<b>85</b>	0.37	0.00	85
4A	69	3	<b>388.0</b>	<b>388</b>	1.17	0.02	388
4B	69	3	<b>388.0</b>	<b>388</b>	1.12	0.02	388
4C	69	3	<b>388.0</b>	<b>388</b>	1.20	0.05	388
4D	69	3	403.0	<b>402</b>	1.29	0.23	402
5A	65	3	<b>415.0</b>	<b>415</b>	1.05	0.01	415
5B	65	3	<b>415.0</b>	<b>415</b>	1.07	0.02	415
5C	65	3	<b>415.0</b>	<b>415</b>	1.11	0.02	415
5D	65	3	<b>431.0</b>	<b>431</b>	1.56	0.45	431
6A	50	3	<b>221.0</b>	<b>221</b>	0.87	0.09	221
6B	50	3	<b>229.0</b>	<b>229</b>	0.80	0.03	229
6C	50	3	303.2	<b>303</b>	0.81	0.11	303
7A	66	3	<b>279.0</b>	<b>279</b>	0.99	0.01	279
7B	66	3	<b>279.0</b>	<b>279</b>	1.01	0.01	279
7C	66	3	<b>308.0</b>	<b>308</b>	1.36	0.26	308
8A	63	3	<b>385.0</b>	<b>385</b>	1.03	0.02	385
8B	63	3	<b>385.0</b>	<b>385</b>	1.03	0.02	385
8C	63	3	404.3	<b>404</b>	1.28	0.27	404
9A	92	3	<b>323.0</b>	<b>323</b>	1.65	0.07	323
9B	92	3	<b>323.0</b>	<b>323</b>	1.73	0.06	323
9C	92	3	<b>323.0</b>	<b>323</b>	1.74	0.06	323
9D	92	3	339.9	<b>339</b>	2.09	0.33	338
10A	97	3	<b>424.0</b>	<b>424</b>	2.10	0.09	424
10B	97	3	<b>424.0</b>	<b>424</b>	2.17	0.13	424
10C	97	3	<b>424.0</b>	<b>424</b>	2.12	0.11	424
10D	97	3	443.9	<b>443</b>	2.40	0.29	442
Gap (%)			0.041%	0.015%			
T(min)					1.13		
T*(min)						0.08	
CPU					Xe 3.07G		







**Table EC.24** Results for the MM-kWRPP – 3 vehicles

Inst	E	BCS10 Single	UHGS				BKS	
			Avg-10	Best-10	T	T*	Old	New
C20 <sub>110</sub>	63	<b>125</b>	<b>125.0</b>	<b>125</b>	0.30	0.03	125*	125
C20 <sub>15</sub>	63	<b>90</b>	<b>90.0</b>	<b>90</b>	0.36	0.05	90*	90
C20 <sub>18</sub>	63	<b>112</b>	<b>112.0</b>	<b>112</b>	0.34	0.04	112*	112
C20 <sub>2100</sub>	63	1292	1281.3	<b>1281</b>	0.50	0.16	1287	<u>1281</u>
C20 <sub>2200</sub>	63	2402	<b>2360.0</b>	<b>2360</b>	0.52	0.23	2373	<u>2360</u>
C20 <sub>2500</sub>	63	5733	<b>5673.0</b>	<b>5673</b>	0.57	0.28	5673*	5673
C21 <sub>110</sub>	67	124	<b>122.0</b>	<b>122</b>	0.45	0.09	122*	122
C21 <sub>15</sub>	67	68	<b>67.0</b>	<b>67</b>	0.41	0.09	67*	67
C21 <sub>18</sub>	67	<b>112</b>	<b>112.0</b>	<b>112</b>	0.52	0.10	112*	112
C21 <sub>2100</sub>	67	1164	<b>1144.0</b>	<b>1144</b>	0.71	0.37	1146	<u>1144</u>
C21 <sub>2200</sub>	67	2238	2211.5	<b>2209</b>	0.84	0.50	2210	<u>2209</u>
C21 <sub>2500</sub>	67	5556	<b>5487.0</b>	<b>5487</b>	0.58	0.23	5501	<u>5487</u>
C22 <sub>110</sub>	74	205	<b>201.0</b>	<b>201</b>	0.60	0.21	201*	201
C22 <sub>15</sub>	74	<b>136</b>	<b>136.0</b>	<b>136</b>	0.56	0.11	136*	136
C22 <sub>18</sub>	74	166	<b>163.0</b>	<b>163</b>	0.61	0.17	163*	163
C22 <sub>2100</sub>	74	1058	1040.9	<b>1040</b>	1.32	0.83	1049	<u>1040</u>
C22 <sub>2200</sub>	74	1943	1908.1	<b>1907</b>	1.32	0.83	1909	<u>1907</u>
C22 <sub>2500</sub>	74	6054	6016.0	<b>6013</b>	1.18	0.70	6058	<u>6013</u>
C23 <sub>110</sub>	78	149	146.3	<b>146</b>	0.94	0.42	147	<u>146</u>
C23 <sub>15</sub>	78	90	<b>89.0</b>	<b>89</b>	0.66	0.16	89*	89
C23 <sub>18</sub>	78	132	<b>131.0</b>	<b>131</b>	0.77	0.15	131*	131
C23 <sub>2100</sub>	78	1231	1220.9	<b>1219</b>	1.44	0.90	1221	<u>1219</u>
C23 <sub>2200</sub>	78	2371	<b>2286.0</b>	<b>2286</b>	1.10	0.64	2298	<u>2286</u>
C23 <sub>2500</sub>	78	5507	5453.1	<b>5435</b>	1.32	0.80	5445	<u>5435</u>
C24 <sub>110</sub>	55	<b>118</b>	<b>118.0</b>	<b>118</b>	0.28	0.05	118*	118
C24 <sub>15</sub>	55	<b>87</b>	<b>87.0</b>	<b>87</b>	0.28	0.03	87*	87
C24 <sub>18</sub>	55	120	<b>119.0</b>	<b>119</b>	0.39	0.05	119*	119
C24 <sub>2100</sub>	55	964	951.3	<b>951</b>	0.66	0.38	957	<u>951</u>
C24 <sub>2200</sub>	55	1940	1928.7	<b>1928</b>	0.52	0.24	1928	1928
C24 <sub>2500</sub>	55	4578	4555.6	<b>4555</b>	0.52	0.22	4555*	4555
Gap (%)		0.230%	0.008%	0.000%				
T(min)		0.41			0.18			
T*(min)						0.07		
CPU		I2 2.4G		Xe 3.07G				

**Table EC.25** Results for the MM-kWRPP – 4 vehicles

Inst	E	BCS10 Single	UHGS				BKS	
			Avg-10	Best-10	T	T*	Old	New
C20 <sub>110</sub>	63	<b>108</b>	<b>108.0</b>	<b>108</b>	0.33	0.05	108*	108
C20 <sub>15</sub>	63	76	<b>75.0</b>	<b>75</b>	0.37	0.07	75*	75
C20 <sub>18</sub>	63	<b>94</b>	<b>94.0</b>	<b>94</b>	0.35	0.06	94*	94
C20 <sub>2100</sub>	63	1035	<b>1029.0</b>	<b>1029</b>	0.44	0.15	1029*	1029
C20 <sub>2200</sub>	63	1911	<b>1891.0</b>	<b>1891</b>	0.42	0.14	1925	<u>1891</u>
C20 <sub>2500</sub>	63	4548	<b>4520.0</b>	<b>4520</b>	0.37	0.09	4533	<u>4520</u>
C21 <sub>110</sub>	67	100	<b>99.0</b>	<b>99</b>	0.55	0.10	99*	99
C21 <sub>15</sub>	67	<b>56</b>	<b>56.0</b>	<b>56</b>	0.41	0.09	56*	56
C21 <sub>18</sub>	67	92	<b>91.0</b>	<b>91</b>	0.48	0.15	91*	91
C21 <sub>2100</sub>	67	900	<b>888.0</b>	<b>888</b>	0.55	0.16	888	888
C21 <sub>2200</sub>	67	1740	1709.7	<b>1709</b>	0.77	0.40	1737	<u>1709</u>
C21 <sub>2500</sub>	67	4334	4293.9	<b>4289</b>	0.62	0.28	4348	<u>4289</u>
C22 <sub>110</sub>	74	167	<b>161.0</b>	<b>161</b>	0.62	0.23	167	<u>161</u>
C22 <sub>15</sub>	74	113	<b>110.0</b>	<b>110</b>	0.80	0.24	111	<u>110</u>
C22 <sub>18</sub>	74	133	131.4	<b>131</b>	0.73	0.30	132	<u>131</u>
C22 <sub>2100</sub>	74	840	824.8	<b>824</b>	1.12	0.65	832	<u>824</u>
C22 <sub>2200</sub>	74	1517	1502.5	<b>1501</b>	1.05	0.58	1515	<u>1501</u>
C22 <sub>2500</sub>	74	4812	4761.0	<b>4752</b>	0.92	0.45	4797	<u>4752</u>
C23 <sub>110</sub>	78	123	<b>120.0</b>	<b>120</b>	0.77	0.28	121	<u>120</u>
C23 <sub>15</sub>	78	74	<b>73.0</b>	<b>73</b>	0.62	0.14	73	73
C23 <sub>18</sub>	78	109	<b>107.0</b>	<b>107</b>	0.66	0.17	109	<u>107</u>
C23 <sub>2100</sub>	78	957	939.5	<b>937</b>	1.16	0.69	950	<u>937</u>
C23 <sub>2200</sub>	78	1808	1740.8	<b>1733</b>	1.26	0.82	1752	<u>1733</u>
C23 <sub>2500</sub>	78	4308	4214.1	<b>4208</b>	1.44	0.95	4264	<u>4208</u>
C24 <sub>110</sub>	55	99	<b>97.0</b>	<b>97</b>	0.32	0.08	97*	97
C24 <sub>15</sub>	55	<b>72</b>	<b>72.0</b>	<b>72</b>	0.31	0.06	72*	72
C24 <sub>18</sub>	55	<b>101</b>	<b>101.0</b>	<b>101</b>	0.35	0.08	101*	101
C24 <sub>2100</sub>	55	770	<b>767.0</b>	<b>767</b>	0.41	0.11	770	<u>767</u>
C24 <sub>2200</sub>	55	1598	1585.6	<b>1584</b>	0.56	0.27	1593	<u>1584</u>
C24 <sub>2500</sub>	55	3663	3619.8	<b>3618</b>	0.59	0.31	3634	<u>3618</u>
Gap (%)		0.303%	0.014%	0.000%				
T(min)		0.29			0.18			
T*(min)						0.06		
CPU		I2 2.4G		Xe 3.07G				

**Table EC.26** Results for the MM-kWRPP – 5 vehicles

Inst	E	BCS10 Single	UHGS				BKS	
			Avg-10	Best-10	T	T*	Old	New
C20 <sub>110</sub>	63	<b>95</b>	<b>95.0</b>	<b>95</b>	0.33	0.05	95	95
C20 <sub>15</sub>	63	<b>68</b>	<b>68.0</b>	<b>68</b>	0.42	0.07	68*	68
C20 <sub>18</sub>	63	88	87.5	<b>87</b>	0.43	0.13	87*	87
C20 <sub>2100</sub>	63	<b>905</b>	<b>905.0</b>	<b>905</b>	0.40	0.10	905	905
C20 <sub>2200</sub>	63	1676	<b>1649.0</b>	<b>1649</b>	0.57	0.28	1663	<u>1649</u>
C20 <sub>2500</sub>	63	3901	3796.2	<b>3793</b>	0.54	0.27	3879	<u>3793</u>
C21 <sub>110</sub>	67	87	<b>86.0</b>	<b>86</b>	0.42	0.08	86*	86
C21 <sub>15</sub>	67	51	<b>50.0</b>	<b>50</b>	0.43	0.10	50*	50
C21 <sub>18</sub>	67	80	<b>79.0</b>	<b>79</b>	0.43	0.11	79*	79
C21 <sub>2100</sub>	67	755	747.3	<b>747</b>	0.73	0.35	753	<u>747</u>
C21 <sub>2200</sub>	67	1525	1469.9	<b>1468</b>	0.64	0.31	1519	<u>1468</u>
C21 <sub>2500</sub>	67	3662	3581.3	<b>3577</b>	0.76	0.43	3637	<u>3577</u>
C22 <sub>110</sub>	74	145	139.1	<b>139</b>	1.05	0.53	141	<u>139</u>
C22 <sub>15</sub>	74	97	<b>95.0</b>	<b>95</b>	0.67	0.19	96	<u>95</u>
C22 <sub>18</sub>	74	117	113.3	<b>113</b>	0.83	0.34	114	<u>113</u>
C22 <sub>2100</sub>	74	707	692.5	<b>692</b>	0.94	0.43	704	<u>692</u>
C22 <sub>2200</sub>	74	1282	1257.9	<b>1256</b>	1.25	0.74	1280	<u>1256</u>
C22 <sub>2500</sub>	74	4069	4005.9	<b>4001</b>	1.09	0.63	4046	<u>4001</u>
C23 <sub>110</sub>	78	107	<b>104.0</b>	<b>104</b>	0.83	0.34	105	<u>104</u>
C23 <sub>15</sub>	78	65	<b>64.0</b>	<b>64</b>	0.69	0.15	65	<u>64</u>
C23 <sub>18</sub>	78	96	<b>93.0</b>	<b>93</b>	0.69	0.25	95	<u>93</u>
C23 <sub>2100</sub>	78	788	774.2	<b>771</b>	1.33	0.87	787	<u>771</u>
C23 <sub>2200</sub>	78	1459	1409.2	<b>1405</b>	1.45	1.02	1455	<u>1405</u>
C23 <sub>2500</sub>	78	3623	3512.5	<b>3505</b>	1.67	1.11	3545	<u>3505</u>
C24 <sub>110</sub>	55	86	<b>85.0</b>	<b>85</b>	0.32	0.09	85*	85
C24 <sub>15</sub>	55	<b>63</b>	<b>63.0</b>	<b>63</b>	0.31	0.07	63*	63
C24 <sub>18</sub>	55	90	<b>89.0</b>	<b>89</b>	0.34	0.09	89*	89
C24 <sub>2100</sub>	55	673	666.2	<b>666</b>	0.54	0.23	666*	666
C24 <sub>2200</sub>	55	1379	1375.6	<b>1374</b>	0.53	0.26	1374*	1374
C24 <sub>2500</sub>	55	<b>3091</b>	<b>3091.0</b>	<b>3091</b>	0.39	0.12	3091	3091
Gap (%)		0.392%	0.021%	0.000%				
T(min)		0.24			0.19			
T*(min)						0.07		
CPU		I2 2.4G		Xe 3.07G				

**Table EC.27** Results for the MM-kWRPP – 6 vehicles

Inst	E	UHGS				BKS	
		Avg-10	Best-10	T	T*	Old	New
C16 <sub>110</sub>	34	<b>57.0</b>	<b>57</b>	0.15	0.03	57*	57
C16 <sub>15</sub>	34	<b>36.0</b>	<b>36</b>	0.14	0.02	36*	36
C16 <sub>18</sub>	34	<b>54.0</b>	<b>54</b>	0.10	0.01	54*	54
C16 <sub>2100</sub>	34	365.2	<b>362</b>	0.17	0.05	362*	362
C16 <sub>2200</sub>	34	<b>735.0</b>	<b>735</b>	0.17	0.05	735*	735
C16 <sub>2500</sub>	34	<b>1584.0</b>	<b>1584</b>	0.24	0.11	1593	<u>1584</u>
C17 <sub>110</sub>	17	<b>41.0</b>	<b>41</b>	0.04	0.00	41*	41
C17 <sub>15</sub>	17	<b>28.0</b>	<b>28</b>	0.04	0.00	28*	28
C17 <sub>18</sub>	17	<b>38.0</b>	<b>38</b>	0.03	0.00	38*	38
C17 <sub>2100</sub>	17	<b>404.0</b>	<b>404</b>	0.03	0.00	404*	404
C17 <sub>2200</sub>	17	<b>523.0</b>	<b>523</b>	0.05	0.01	523*	523
C17 <sub>2500</sub>	17	<b>1497.0</b>	<b>1497</b>	0.04	0.01	1497*	1497
C18 <sub>110</sub>	16	<b>71.0</b>	<b>71</b>	0.04	0.00	71*	71
C18 <sub>15</sub>	16	<b>48.0</b>	<b>48</b>	0.03	0.00	48*	48
C18 <sub>18</sub>	16	<b>58.0</b>	<b>58</b>	0.03	0.00	58*	58
C18 <sub>2100</sub>	16	<b>577.0</b>	<b>577</b>	0.04	0.01	577*	577
C18 <sub>2200</sub>	16	<b>1027.0</b>	<b>1027</b>	0.04	0.00	1027*	1027
C18 <sub>2500</sub>	16	<b>2584.0</b>	<b>2584</b>	0.04	0.01	2584*	2584
C19 <sub>110</sub>	29	<b>99.0</b>	<b>99</b>	0.10	0.02	99*	99
C19 <sub>15</sub>	29	<b>67.0</b>	<b>67</b>	0.09	0.01	67*	67
C19 <sub>18</sub>	29	<b>91.0</b>	<b>91</b>	0.09	0.01	91*	91
C19 <sub>2100</sub>	29	<b>604.0</b>	<b>604</b>	0.10	0.02	604*	604
C19 <sub>2200</sub>	29	<b>946.0</b>	<b>946</b>	0.12	0.03	975	<u>946</u>
C19 <sub>2500</sub>	29	<b>2506.0</b>	<b>2506</b>	0.13	0.04	2506*	2506
C20 <sub>110</sub>	63	<b>89.0</b>	<b>89</b>	0.31	0.06	89*	89
C20 <sub>15</sub>	63	<b>65.0</b>	<b>65</b>	0.37	0.08	65	65
C20 <sub>18</sub>	63	84.3	<b>84</b>	0.48	0.19	84*	84
C20 <sub>2100</sub>	63	<b>820.0</b>	<b>820</b>	0.50	0.20	829	<u>820</u>
C20 <sub>2200</sub>	63	<b>1476.0</b>	<b>1476</b>	0.51	0.21	1487	<u>1476</u>
C20 <sub>2500</sub>	63	3324.4	<b>3312</b>	0.58	0.30	3431	<u>3312</u>
C21 <sub>110</sub>	67	<b>78.0</b>	<b>78</b>	0.50	0.17	79	<u>78</u>
C21 <sub>15</sub>	67	<b>46.0</b>	<b>46</b>	0.42	0.10	47	<u>46</u>
C21 <sub>18</sub>	67	<b>74.0</b>	<b>74</b>	0.45	0.11	74*	74
C21 <sub>2100</sub>	67	660.3	<b>660</b>	0.78	0.44	660*	660
C21 <sub>2200</sub>	67	1280.3	<b>1274</b>	0.78	0.44	1299	<u>1274</u>
C21 <sub>2500</sub>	67	3181.8	<b>3174</b>	0.78	0.44	3230	<u>3174</u>
C22 <sub>110</sub>	74	<b>125.0</b>	<b>125</b>	0.66	0.22	128	<u>125</u>
C22 <sub>15</sub>	74	<b>85.0</b>	<b>85</b>	0.65	0.21	86	<u>85</u>
C22 <sub>18</sub>	74	101.2	<b>101</b>	0.84	0.42	103	<u>101</u>
C22 <sub>2100</sub>	74	606.1	<b>606</b>	1.25	0.71	619	<u>606</u>
C22 <sub>2200</sub>	74	1104.4	<b>1104</b>	1.01	0.53	1118	<u>1104</u>
C22 <sub>2500</sub>	74	3507.7	<b>3500</b>	1.34	0.83	3548	<u>3500</u>
C23 <sub>110</sub>	78	<b>95.0</b>	<b>95</b>	0.86	0.28	96	<u>95</u>
C23 <sub>15</sub>	78	<b>59.0</b>	<b>59</b>	0.62	0.15	60	<u>59</u>
C23 <sub>18</sub>	78	<b>85.0</b>	<b>85</b>	0.65	0.17	86	<u>85</u>
C23 <sub>2100</sub>	78	660.4	<b>658</b>	1.57	1.05	675	<u>658</u>
C23 <sub>2200</sub>	78	1184.2	<b>1179</b>	1.28	0.85	1238	<u>1179</u>
C23 <sub>2500</sub>	78	3023.7	<b>3002</b>	1.48	1.00	3071	<u>3002</u>
C24 <sub>110</sub>	55	<b>79.0</b>	<b>79</b>	0.29	0.06	79*	79
C24 <sub>15</sub>	55	<b>58.0</b>	<b>58</b>	0.31	0.06	58*	58
C24 <sub>18</sub>	55	<b>84.0</b>	<b>84</b>	0.36	0.09	84*	84
C24 <sub>2100</sub>	55	613.7	<b>613</b>	0.66	0.35	613*	613
C24 <sub>2200</sub>	55	<b>1233.0</b>	<b>1233</b>	0.54	0.27	1233*	1233
C24 <sub>2500</sub>	55	2737.3	<b>2737</b>	0.55	0.27	2743	<u>2737</u>
Gap (%)	0.031%	0.000%					
T(min)				0.19			
T*(min)						0.08	
CPU	Xe 3.07G						